

# polyLARVA: Technology Agnostic Runtime Verification

Christian Colombo    Adrian Francalanza    Ruth Mizzi    Gordon J. Pace

Department of Computer Science  
University of Malta

{christian.colombo|adrian.francalanza|rmiz0015|gordon.pace}@um.edu.mt

With numerous specialised technologies available to industry, it is become increasingly frequent for computer systems to be composed of heterogeneous components, built over, and using different technologies and languages. While this enables developers to use the appropriate technologies for specific contexts, it becomes more challenging to ensure the correctness of the overall system. In this paper we propose a framework to enable extensible technology agnostic runtime verification and we present an extension of polyLARVA, a runtime-verification tool able to handle the monitoring of heterogeneous-component system. The approach is then applied to a case study with C and Java components.

## 1 Introduction

Component-based approaches to software and service design is becoming more widespread, allowing for heterogeneous systems to be compartmentalised into components; these components encapsulate their internal behaviour while revealing uniform interfaces through which other components may interact. This approach to system arrangement has facilitated the construction of large complex systems, where each component is allowed to internally employ different technologies from operating systems and hardware to programming languages.

However, the sheer complexity of the systems constructed, together with the decentralised nature of how these heterogenous systems are developed, creates new potential points of failure. This, in turn, increases the need for some form of correctness verification. Runtime monitoring [12] has been shown to be a viable solution to the verification of large complex systems — by limiting the analysis to the actual current runtime path, the approach is tractable (for a reasonable choice of specification logic), while still guaranteeing detection (albeit at runtime) of property violations. Especially in the context of open systems, where correctness is also partly dependent on interaction with the environment, this approach has proved to be viable and scalable to real-life industrial systems e.g. [3, 1].

An important research question in the field of runtime verification has been how to extend the approach to handle non-monolithic systems. Much of the work so far has been limited to a high-level view of such systems, treating components as *black boxes* and focussing instead on the verification of the component interactions and on strategies for engineering monitoring in such distributed settings. To date, there has been little work that attempts to push verification *inside* the components so as to verify their inner workings as part of the wider system. This poses additional challenges to runtime verification, both practical and theoretic, such as the need for a standardised framework for generating events to be monitored (irrespective of the underlying technology used by the monitored components), or the adaptation of the monitoring logic (specifying correctness properties) to distinct underlying technologies used by different components, which may invariably lead to different semantic interpretation of said logics.

A typical example would be an online betting system, which is inherently component based (e.g., the web-portal subsystem, the billing subsystem, the fraud detection subsystem etc), and where each

component may use different technologies and resources. Typical operations in this system, such as an online-betting transaction, may go through different components ranging from the company’s transaction database, to an internal certified logging component kept for legal purposes, to an external bank system. Whereas existing verification technologies would typically only be able to monitor external units as black boxes, thus only referring to the interactions taking place, consistency and correctness properties of such betting transactions may depend on the inner workings of internal components. For instance, one may want to ensure that the value written in the certified logging component matches that written in the company’s transaction database. Such component-spanning properties have to be instrumented on different components, necessitating direct interaction with the underlying technologies of each component. Despite this, the present lack of tools supporting such technology heterogeneity is a major stumbling block towards a full adoption of runtime verification techniques for component-based systems.

In this paper we present a novel runtime verification framework which supports the monitoring of component-based systems, possibly using different technologies. The approach has been embodied as an extension to the runtime verification tool polyLARVA [2], to support the generation and instrumentation of separate monitoring code for multiple components in a system, and from a single property specification. Furthermore, the property specification language has been designed to be technology and programming-language agnostic, and hence allows reusability across different technologies. To evaluate the proposed component-based monitoring framework, polyLARVA has been applied to OpenEmm, an open source web-based e-mail marketing tool.

In the rest of the paper we start by discussing the various design options when it comes to runtime verifying component-based systems in Section 2. In Section 3, these issues are then discussed in the context of multi-technology component-based systems where the monitor needs to access the components’ internal states. Subsequently, Sections 4 and 5 describe how we support our design decisions through an extensible monitoring framework polyLARVA. In Section 6 we give account of a case-study and discuss related work and conclusions in the last section.

## 2 Component-Based System Monitoring

From a monitoring perspective, component-based systems pose particular challenges which go beyond those presented by monolithic systems. While the well-behaviour of individual components can be dealt with locally, properties and specifications which span across components and expressing the correctness of the system as a whole raise various issues, both pragmatic and conceptual. The cross-cutting nature of monitoring, which may require access to the internal state of different components, raises issues regarding the architecture of the monitors with respect to the system components and practical challenges in systems with components built using different technologies.

Two important issues, which influence choices in the monitoring architecture are:

**Orchestrated vs. choreographed monitoring:** Different approaches have been proposed in the literature regarding the locality (physical or conceptual) of the monitors with respect to the different components in a system. In an orchestrated monitoring approach (e.g. see [8, 13, 11]), the monitor is a separate unit from the rest of the system, but with privileges to enable it to listen to the behaviour of the other parts of the system and pass judgement about them. On the other hand, a choreographed approach (e.g. see [16, 17, 15]) extends the different parts of the system such that each part *locally* eavesdrops on its own behaviour, ensuring correct behaviour, and communicating with other monitors (belonging to other parts of the system) whenever synchronisation or information regarding the other parts is required. Approaches to split monitors in this manner using

a static [16] or dynamic [7] manner have been proposed. Although most of this work focuses on distributed systems, the classification applies equally well to general component-based systems. Orchestration is usually easier to instrument and setup, but adds dependencies between components which may not be desirable. On the other hand, a choreographed approach respects locality, but at the cost of more complex instrumentation of monitoring code and specification slicing. Especially in the case of heterogeneous component-based systems with components using different technologies, local instrumentation can prove to be particularly challenging, since the monitoring tool has to be able to instrument code built in different languages.

**Intrusiveness of monitoring:** Another independent choice is the level of abstraction at which the monitors can eavesdrop the behaviour of the system. Much work, especially in setting of services, focuses on specifications of messages passed between components [8, 13, 11]. This black-box approach ensures that the instrumentation of monitors is relatively straightforward, since they only need to hook to the communication channels and process the behaviour appropriately. However, the approach has serious limitations when specifications refer to the components' local states or internal events, since without re-engineering to expose such states and events their monitoring would not be possible. In a component-based system setting, such properties would, for instance, be required to ensure data-consistency across components. A more intrusive approach, to enable direct access to the state of the system components poses challenges, especially when the components are built using different technologies.

Note that these two issues are largely independent of each other. For instance both black-box orchestrated monitoring [8, 13] and intrusive orchestrated monitoring [11] approaches have been investigated in the literature. In the next section we investigate how challenges in extending the existing approaches to technology-diverse component-based systems can be addressed.

### 3 Intrusive Monitoring of Component-Based Systems

Although specifications which can be checked through eavesdropping on component communication suffices for a number of applications, there are situations where more intrusive monitoring is required. In the case of systems built from components using different technologies this poses various challenges.

#### 3.1 Monitoring Architecture

Supporting technology agnosticism, using an intrusive choreographed approach limits extensibility, since for each possible technology would have to be coupled with each other one. In contrast, with an orchestrated approach, adding support for a new technology only requires adding support for instrumenting code to communicate with the central monitor.

For this reason, we have adopted a centralised monitor which receives events and processes them, thus making it a largely orchestrated approach — with the verification unit of the tool being independent of the the other components. For every supported technology, the tool can intrusively instrument event generation into the components.

Since a purely orchestrated yet intrusive approach may result in breaches of data and control encapsulation within components, we have built into the monitoring system manually instrument parts of the monitoring logic onto the system-side [2]. This allows for manual choreography, without sacrificing extensibility, since inter-component coordination still happens through the central orchestrated monitor.

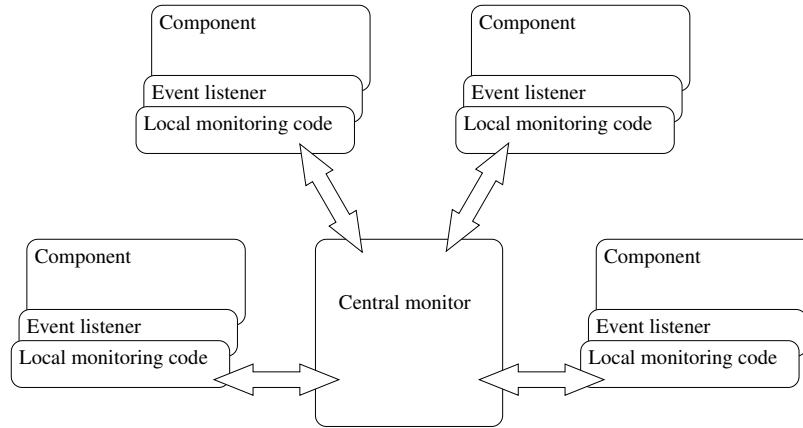


Figure 1: Architecture of instrumented monitors on components

Fig. 1 shows the general architecture of the system after the monitoring parts are instrumented onto the components.

### 3.2 Extensible Technology Support

Through the general monitoring architecture, adding support for new technologies requires the possibility of instrumenting monitors into components built using that technology. The responsibility of the monitoring code is to (i) generate events to send to the central monitor; and (ii) execute any local monitoring code specified in the properties.

Since it is desirable to support cross-component properties, it is crucial to support monitoring instrumentation from a single specification script, the approach adopted in polyLARVA. Furthermore, since extensibility to further technologies is also an essential feature, we have separated the monitoring instrumentation into different parts (i) the central verification code is generated from the global parts of the specification using a language-independent part of the runtime verification tool; and (ii) for each different technology, a separate tool is provided, which generates the automated instrumentation scripts from a subpart of the specification (as tagged by the user).

The workflow for the usage of polyLARVA is depicted in Fig. 2, with: (i) the specification script is passed through the language agnostic part of the tool to produce the central monitor (bottom arrow of the figure); and (ii) for each different component, the appropriate technology tool of polyLARVA processes the specification to produce the instrumentation instructions for event extraction and local monitoring on that particular component (top arrows in the figure).

The actual instrumentation typically takes place by processing the scripts produced by the polyLARVA language-dependent compilers, and using additional external tools (such as aspect-oriented programming compilers).

### 3.3 Replicated Monitors and Language Agnosticism

Since most systems describe multiple instances of an abstract concept (e.g. multiple users, accounts, sessions, etc.) frequently, one would desire to replicate a property for each instance. For instance, for each bank transaction, one may want to set up a monitor for a property which states that the incoming

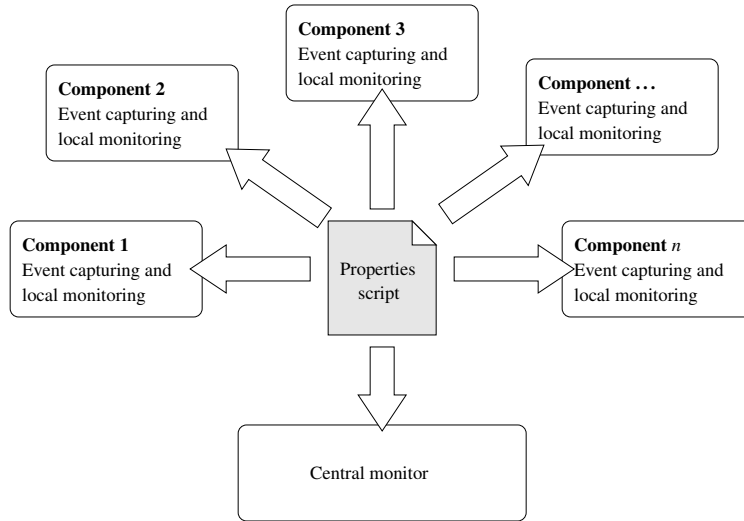


Figure 2: Splitting the specification script across components

and outgoing balances cancel out. In practice, the way such concepts are encoded depends on the technology being used. For instance, if the system is written in Java, a transaction may be encoded as an object, while if written in Erlang, it may correspond to a separate thread and data structures handling the transaction. While monitoring tools for single-technology monolithic systems associate such replication in combination with the technology, in a language agnostic system one needs to be more general. In component-based systems, this poses further challenges when the concept’s lifetime spans across different components. To support such replication, one solution is to demarcate concept instances’ lifetimes by identifying events marking their start and those marking their end. Furthermore, all events related to such an instance are tagged with an identifier which indicates to which instance they belong.

For example, in the bank transaction example, we would identify a call to *initialiseTransaction(transid)* to be the starting point, while *concludeTransaction(transid)* with the same parameter to be its end. Furthermore, any calls to *transferFunds(transid,...)* are associate with the instance of the monitor which was started with the transaction identity passed as parameter.

## 4 polyLARVA Specifications

At the lowest level, our monitoring framework, polyLARVA, uses a simple guarded-command style specification language. Properties are expressed as a list of rules of the following form:

$$event \mid condition \mapsto action$$

Whenever an *event* (possibly having parameters) is generated by the system, the list of monitor rules is scanned for rule matches relating to that event. If a match is found, the expression specified in the *condition* of the rule is evaluated and, if satisfied, the *action* is triggered.

**Example 4.1** Consider a scenario in which one desires to check that an online payment on a web-based system is carried out after the credit card used has been registered — tagging the customer as *untrusted* if this rule is violated. This may be expressed in terms of the rules enclosed within the **rules** block, ①, in Program 4.1. In these rules, *register* and *pay* are system events, parametrised by the values *customer* and *card*;  $\neg$ registeredCards[*card*] is a condition, while registeredCards[*card*] := true and se-

`tUntrusted(customer)` are actions triggered by the runtime monitor. Note that to monitor the property for each system user, we define the event `newCustomerSession` as the point triggering the replication of the rules (②). Consequently, the events declared within this scope must all define a variable `customer` which binds the event with a particular monitor instance. Finally, ③ marks the end of the context upon event `endCustomerSession`.

The `state`, `conditions` and `actions` blocks define specification-specific monitor state, conditions and actions respectively that are used in the `rules` section as macros. The `events` block highlights those specific points which, during system execution, should trigger the monitoring functionality.

*ahna apposta hallejna barra d-definitions ta l-events? ma jidhrux ftit misleading hekk? /C u d-definition ta newCustomerSession fejn tkun normalment fl-iscrypt? /C*

---

**Program 4.1** Monitoring customers attempting payments with unregistered cards

---

```

② upon (newCustomerSession(customer)) {
  state {
    boolean[] registeredCards;
  }
  events {
    register(customer, card)
    pay(customer, card)
    endCustomerSession(customer)
  }
  conditions {
    isRegistered(card) = { registeredCards[card] }
  }
  actions {
    setUntrustedCustomer(customer) = ...
    registerCard(card) = { registeredCards[card] := true }
  }
① rules {
  register(customer, card) -> registerCard(card);
  pay(customer, card) \ !isRegistered(card) -> setUntrustedCustomer(customer);
  ③ endCustomerSession(customer)-> Done;
}
}

```

---

While events are a result of the execution path being followed by a system and therefore system-specific, the conditions and actions defined in a rule are typically specific to the runtime monitoring states and are thus independent from the system. This means that evaluation of rule actions and conditions has no pre-set requirement of running on the system being monitored and can safely be moved onto separate resources. However, this is not always the case since conditions may also query the system state, and actions may alter the system state. Unfortunately, it is not straightforward to automatically delineate system-dependent elements from purely monitoring elements. For this reason, our language enables the user to explicitly specify this separation through appropriate constructs<sup>1</sup>.

**Example 4.2** *Continuing on the previous example, consider the case where instead of keeping track within the monitor state of which cards have been registered or not, we query the system state (which keeps track of card registration anyway). In this case the condition which checks whether a credit card*

---

<sup>1</sup>A lengthy discussion of how one could use the distinction between system-side and monitor-side monitoring to optimise efficiency can be found in [2].

is registered is performed on the system state by being tagged as `systemSide`, ⑤. Note that the rest of the monitoring elements are still performed on the monitor side and are thus marked as such.

---

**Program 4.2** Monitoring activity between two components
 

---

```

upon (newCustomerSession(customer)) {
  state {
    remoteSide {int cardNo; }
  }
  events {
    register(customer, card)
    pay(customer, card)
    paymentservice receiveDetails(customer, card)
    endCustomerSession(customer)
  }
  conditions {
    remoteSide { validateCardDetails(card) = cardNo == card; }

    ⑤ systemSide { isRegistered(card) = { registeredCards[card] } }
  }
  actions {
    remoteSide { setUntrustedCustomer(customer) = ... }
    remoteSide { saveCardDetails(card) = cardNo := card.cardNo; }
    remoteSide { reportError = ... }
  }
  rules {
    pay(customer, card) \ !isRegistered(card) -> setUntrustedCustomer(customer);
    pay(customer, card) \ isRegistered(card) -> saveCardDetails(card);
    receiveDetails(customer, card) \ !validateCardDetails(card) -> reportError();
    endCustomerSession(customer)->Done;
  }
}

```

---

The tagging of monitor side and system side evaluation of monitoring logic suffices for a monolithic system. However, when a system is composed of heterogeneous components residing on different technologies, tagging has to be more comprehensive, *i.e.*, it has to also distinguish across components. There are two aspects to this: (i) as before states, conditions, and actions may reside within different components; and (ii) the events may now also arise from different components. The following example, demonstrates the use of extended tagging in the context of multiple components.

**Example 4.3** *As an example, we identify the fact that many online stores normally incorporate functionality offered by payment gateway web services to validate credit card details and accept transactions. In such a scenario we might want to ensure that the payment details input by the customer on the online system are the same details received by the payment gateway. A potential monitoring setup for such a scenario is as illustrated in Figure 3 where a system side monitor, associated with the online store system, notifies the remote runtime monitor about a payment transaction. Program 4.3 highlights how the proposed setup is facilitated through the use of the constructs supplied by the polyLARVA specification language. The payment transaction event ③ is identified as being an event that will occur on one particular system component through the use of a user-defined label that identifies the component as store. The event is also parametrised with the credit card details entered to the system, thus ensuring that the global monitoring component can maintain a copy of these values. On receipt of an authorisation*

request, the system side monitor associated with the payment gateway service, will communicate with the global monitoring component in order to trigger validation of the card details. This communication is identified through the definition of the event `receiveDetails` ④ which specifies that its source is one labelled as `paymentservice`, and therefore a different process from that which triggered other events.

---

**Program 4.3** Monitoring activity between two components

---

```

upon (newCustomerSession(customer)) {
  state {
    remoteSide {int cardNo; }
  }
  events {
    ③ event@store register(customer,card)
    ③ event@store pay(customer,card)
    ④ event@paymentservice receiveDetails(customer,card)
    ③ event@store endCustomerSession(customer)
  }
  conditions {
    remoteSide { validateCardDetails(card) = cardNo == card; }
  }
  ⑤ systemSide@store { isRegistered(card) = { registeredCards[card] } }
  }
  actions {
    remoteSide { setUntrustedCustomer(customer) = ... }
    remoteSide { saveCardDetails(card) = cardNo := card.cardNo; }
    remoteSide { reportError = ...}
  }
  rules {
    pay(customer,card) \ !isRegistered(card) -> setUntrustedCustomer(customer);
    pay(customer,card) \ isRegistered(card) -> saveCardDetails(card);
    receiveDetails(customer,card) \ !validateCardDetails(card) -> reportError();
    endCustomerSession(customer)->Done;
  }
}

```

---



Figure 3: Centralised monitoring for online payments

## 5 Extending polyLARVA for New Technologies

Due to the intrusive nature of monitoring in polyLARVA, technology-specific plugins had to be implemented to support the monitoring of OpenEmm: one for Java and another for C. The two technologies differ significantly resulting in different design choices for the two plugins:

**Eliciting events** To extract monitoring information from the components we decided to opt for aspect-oriented programming (AOP) extensions for both Java and C. AOP has been heavily used for



implementing runtime verification tools [?] and both Java and C have relatively mature AOP extensions being AspectJ and ACC — the AspeCt-oriented C compiler<sup>2</sup>, respectively.

**Communication to and from monitoring component** When an event triggers at a component, this has to be communication to the global monitoring component. To this end, both Java and C plugins use standard socket connections. The same technology is also used to support communication from the monitoring component back to the other components. This is useful when monitoring conditions and actions have to access system state residing on a particular component.

I suggest adding the parts below around the text above /G

- The architecture/design within a language plugin.
- How does one add support for a new language? How much work is required?
- How similar/different are the C and Java plugins?

## 6 Case Study

polyLARVA has been used to monitor OpenEmm<sup>3</sup>, an open source, web-based tool for email marketing. The tool provides facilities to administer mailing lists, create email shots, schedule automatic email sending, track sent emails, and manage bounced emails. OpenEmm claims over hundreds of thousands of downloads to date and has a user base that includes a number of prominent large-scale companies. The tool is ideal as a case study for polyLARVA's technology-agnostic setup due to its component-based setup and the hybrid technologies it adopts: the core part is written in Java while performance-sensitive component are written in C [more details go here](#) /C. The case study focuses on monitoring the behaviour of OpenEmm to send out a customised email shot. The tool provides users with a content management interface that allows easy customisation and creation of email templates. When an email is sent out to a mailing list, each individual recipient will receive a personalised email, built by an automated process that customizes each email based on the given template. This functionality is a result of a process which flows across two components: a Java component which retrieves the email template and mailing list information from the database and collates them in an XML file, and a C component which receives this information and produces the final emails. This control flow is displayed using sequentially numbered arrows, in Figure 4.

Though unlikely, this setup introduces the possibility of a discrepancy between the two components: e.g., the XML file getting corrupted, or the database gets updated resulting in outdated emails being sent. This concern has been addressed by (intrusively) monitoring the components' states and identifying any inconsistencies which may occur. The rest of this section elaborates on how this was achieved.

### 6.1 Specifying Properties for OpenEmm

A basic property that can be monitored to ensure the XML file has not been tampered is to ensure that the total number of recipients in the mailing list is the same within the C component as it is in the Java component. This may be expressed in terms of the following rule:

$$c\_sendMails() \mid (java\_mailCount \neq c\_mailCount) \mapsto logIncorrectCount;$$

<sup>2</sup><https://sites.google.com/a/agapp.msrg.utoronto.ca/aspectc>

<sup>3</sup><http://www.openemm.org/>

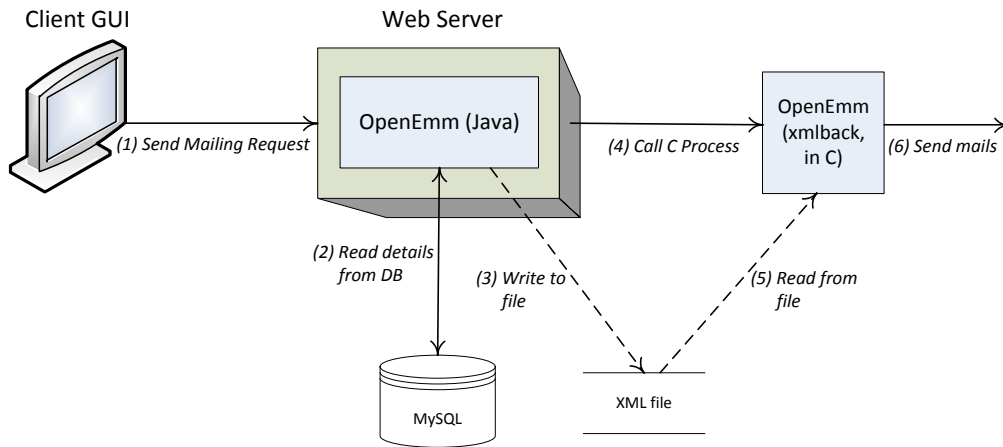


Figure 4: Partial architecture of the OpenEmm framework

where  $c\_sendMails$  is a system event occurring on the OpenEmm C process at the point when mails are about to be sent; the rule condition specifies that the total number of mails being sent by the C process  $c\_mailCount$  must be equal to the total count of mails that was specified at the Java component.  $logIncorrectCount$  is the action taken by the monitor if the values are not equal. Figure 5 “process” should change to “component” and “remote” to “global” /C shows the flow of control amongst the components involved in monitoring the specified property: the monitoring listeners instrumented within the Java and C components and the global monitoring component. When the details of a mail shot are available, the Java component notifies the global monitor with the total number of subscribers. The value is stored as the monitor state  $java\_mailCount$ . Subsequently, when the C component receives the XML file, the total number of mail recipients is communicated to the global monitor which in turn compares it to  $java\_mailCount$ . Program 6.1 shows the polyLARVA specification required to generate this setup. In ⑥ and ⑦ the specification distinguishes between the events’ component-sources while the main rule is specified in ⑧.

Another property which we monitored on OpenEmm is to ensure that users which for some reason have been “blacklisted” are never sent an email. OpenEmm adheres to this property by carrying out a filtering exercise on the mailing list recipients, leaving out any blacklisted recipients. However, if an email recipient is blacklisted while the mailing generation process is already running, there could be circumstances where the recipient is still included in the mailing list. Such an issue could be detected by setting up a polyLARVA monitor which verifies that each recipient is still non-blacklisted at the time of being sent an email. Figure 6 depicts how the global runtime monitor can be notified upon the creation of a personalised email ( $inCreateMail$ ), triggering the monitoring process to query the blacklist on the Java component ( $isEmailBlackListed?$ ). Program 6.2 shows how the property of verifying users to be non-blacklisted can be specified in polyLARVA. In particular note that the condition is specified as  $systemSide@javaComponent$  meaning that the condition is to be executed within the Java component and the result is then communicated to the global monitor.

The two properties described above have been successfully compiled by polyLARVA and applied to OpenEmm. Each specification script was processed by three compilers — (i) the standard polyLARVA com-

---

**Program 6.1** Monitoring count of mailshot recipients

---

```

upon (newMailShot(mailshotID)) {
  state {
    monitorSide {int java_mailCount; }
  }
  events {
    ⑥ event@javaComponent callMailingExecution(mailshotID, javaSubsCount)
    ⑦ event@cComponent startXMLProcessing(mailshotID, c_mailCount)
  }
  conditions {
    monitorSide { invalidMailCount(c_mailCount) =
                  java_mailCount != c_mailCount; }
  }
  actions {
    monitorSide { setJavaMailCount(javaSubsCount) =
                  java_mailCount == javaSubsCount; }
    monitorSide { logIncorrectCount = ...}
  }
  rules {
    callMailingExecution(mailshotID, javaSubsCount) \ true ->
      setJavaMailCount(javaSubsCount);
    ⑧ startXMLProcessing(mailshotID, c_mailCount) \ invalidMailCount ->
      logIncorrectCount;
  }
}

```

---



---

**Program 6.2** Monitoring blacklisted recipients

---

```

upon (newCustomer(custID)) {
  events {
    event@cComponent inCreateMail(c_custID)
  }
  conditions {
    systemSide@javaComponent{ isEmailBlacklisted(c_custID) =
                              database query to check for presence
                              of c_custID in blacklist .... }
  }
  actions {
    monitorSide { logBlacklisted = ...}
  }
  rules {
    checkEmail = inCreateMail(mailshotID, c_custID) \ isEmailBlacklisted ->
      logEmailBlacklisted;
  }
}

```

---

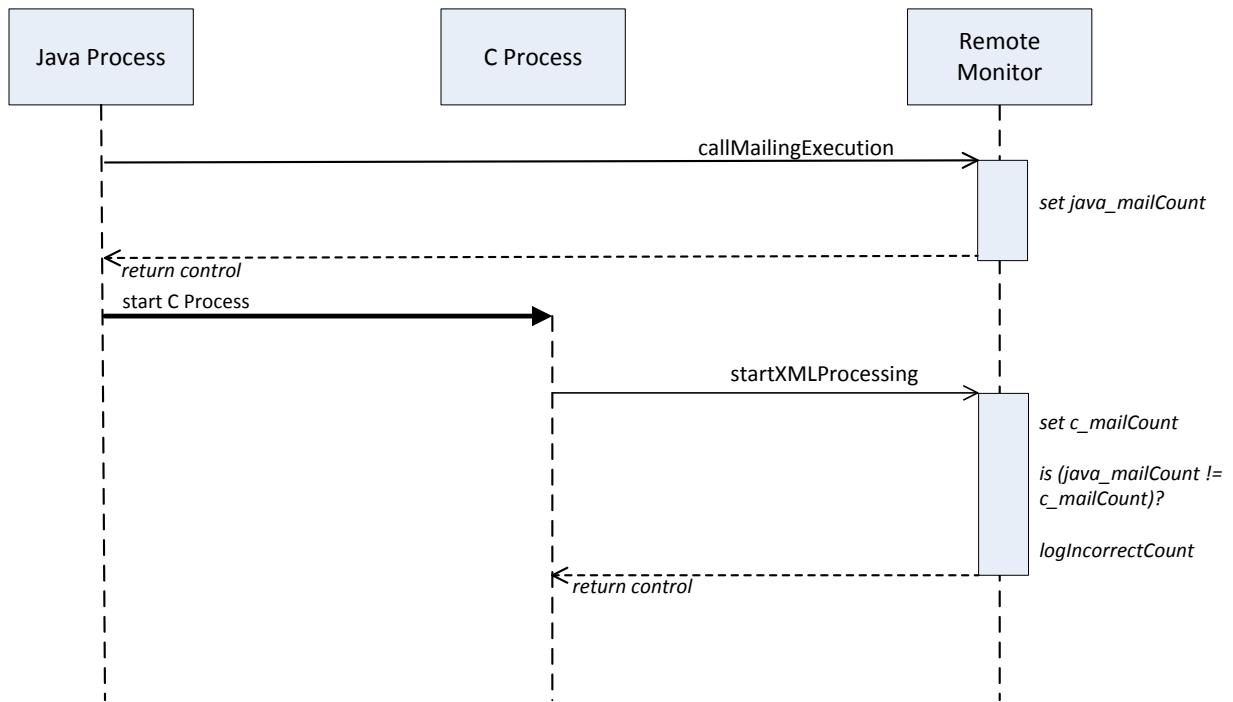


Figure 5: Monitoring email shot subscribers

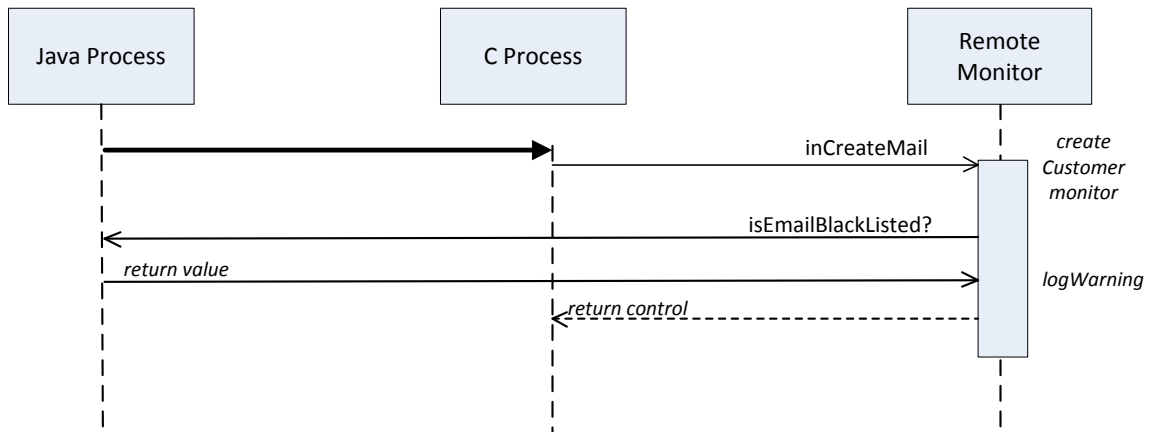


Figure 6: Monitoring blacklisted recipients

piler that creates the global monitoring component, (ii) the polyLARVA language compiler in conjunction with the Java plug-in, to create a component listener that was woven into the Java OpenEmm code, and

(iii) the polyLARVA language compiler in conjunction with the C plug-in, to create a component listener that was woven into the C OpenEmm code. OpenEmm was installed on an Ubuntu operating system while the global monitor component was executed on a separate Windows machine.

No performance tests were carried out during this case study because the aim of this work was to study the interaction between different components running different technologies and the global runtime monitoring component. The possible performance improvements that can be achieved using the polyLARVA framework are highlighted in our other work [2].

## 7 Conclusions

While a significant number of runtime verification frameworks have been proposed in the literature [4, 6, 14, 5, 10, 9], these tools are normally restricted to support one particular programming language or technology and the effort required to support new languages is prohibitive. An exception is the MOP framework [14] whose architecture makes provision for the addition of new language plug-ins that can generate a MOP runtime monitor for a particular programming language. However, MOP does not support the monitoring of a single property across a system with multiple technologies and does not have an inbuilt concept of components. On the other hand, a runtime verification approach has been proposed for the BIP component framework [6] which tackles issues specific to component-based systems. However, the work is positioned at a higher level of abstraction, focusing on the guarantees that are required to ensure sound and correct monitoring within BIP — *i.e.*, the issue of multiple technologies has not been considered in this work.

The non-monolithic nature of component-based systems means that verification techniques have to be adapted to be applicable. In this paper, we have presented an extension of an existing tool, polyLARVA [2], to handle the runtime verification of component-based systems. In particular, we have emphasised the need for the support of multiple-technologies used in such systems, with the resulting tool being easily extensible to handle new technologies. Although we have shown its applicability by deploying it on a third-party open-source, we are currently looking into its use in an industrial setting. Furthermore, we are looking into ways of combining our runtime verification approach to the unit testing of component-based systems.

## References

- [1] Howard Barringer, Allen Goldberg, Klaus Havelund & Koushik Sen (2004): *Rule-Based Runtime Verification*. In: *VMCAI, LNCS 2937*, pp. 44–57.
- [2] Christian Colombo, Adrian Francalanza, Ruth Mizzi & Gordon J. Pace (2012): *polyLarva: Runtime Verification with Configurable Resource-Aware Monitoring Boundaries*. In George Eleftherakis, Mike Hinchey & Mike Holcombe, editors: *SEFM, Lecture Notes in Computer Science 7504*, Springer, pp. 218–232.
- [3] Christian Colombo, Gordon J. Pace & Patrick Abela (2012): *Safer asynchronous runtime monitoring using compensations*. *Formal Methods in System Design* 41(3), pp. 269–294. Available at <http://dx.doi.org/10.1007/s10703-012-0142-8>.
- [4] Christian Colombo, Gordon J. Pace & Gerardo Schneider (2009): *Dynamic Event-Based Runtime Monitoring of Real-Time and Contextual Properties*. In: *Formal Methods for Industrial Critical Systems*, Springer-Verlag, pp. 135–149.
- [5] Marcelo d’Amorim & Klaus Havelund (2005): *Event-based Runtime Verification of Java Programs*. In: *Proceedings of the third international workshop on Dynamic analysis, WODA ’05*, ACM, New York, NY, USA, pp. 1–7.

- [6] Yliès Falcone, Mohamad Jaber, Thanh-Hung Nguyen, Marius Bozga & Saddek Bensalem (2011): *Runtime Verification of Component-Based Systems*. In: *Software Engineering and Formal Methods (SEFM), Lecture Notes in Computer Science 7041*, Springer, pp. 204–220.
- [7] Adrian Francalanza, Andrew Gaudi & Gordon J. Pace (2011): *Distributed System Contract Monitoring*. In: *Fifth Workshop on Formal Languages and Analysis of Contract-Oriented Software (FLACOS'11), Electronic Proceedings in Theoretical Computer Science (EPTCS) 68*.
- [8] Yuan Gan, Marsha Chechik, Shiva Nejati, Jon Bennett, Bill O'Farrell & Julie Waterhouse (2007): *Runtime Monitoring of Web Service Conversations*. In: *Proceedings of the 2007 conference of the center for advanced studies on Collaborative research, CASCON '07, IBM Corp., Riverton, NJ, USA*, pp. 42–57.
- [9] Dimitra Giannakopoulou & Klaus Havelund (2001): *Runtime Analysis of Linear Temporal Logic Specifications*. Technical Report.
- [10] Moonzoo Kim, Mahesh Viswanathan, Sampath Kannan, Insup Lee & Oleg Sokolsky (2004): *Java-MaC: A Run-Time Assurance Approach for Java Programs*. *Form. Methods Syst. Des.* 24(2), pp. 129–155.
- [11] Ingolf Krüger, Michael Meisinger & Massimiliano Menarini (2007): *Runtime Verification of Interactions: From MSCs to Aspects*. In Oleg Sokolsky & Serdar Tasiran, editors: *Runtime Verification, Lecture Notes in Computer Science 4839*, Springer Berlin / Heidelberg, pp. 63–74.
- [12] Martin Leucker & Christian Schallhart (2009): *A Brief Account of Runtime Verification*. *Journal of Logic and Algebraic Programming* 78(5), pp. 293–303. Available at <http://dx.doi.org/10.1016/j.jlap.2008.08.004>.
- [13] Z. Li, Y. Jin & J. Han (2006): *A Runtime Monitoring and Validation Framework for Web Service Interactions*. In: *Software Engineering Conference, 2006. Australian*, pp. 10 pp. –79, doi:10.1109/ASWEC.2006.6.
- [14] Patrick O'Neil Meredith, Dongyun Jin, Dennis Griffith, Feng Chen & Grigore Rosu (2012): *An Overview of the MOP Runtime Verification Framework*. *Int. J. Softw. Tools Technol. Transf.* 14(3), pp. 249–289.
- [15] Thomas S.Cook, Doron Drusinsky & Man-Tak Shing (2007): *Specification, Validation and Run-time Monitoring of SOA Based System-of Systems Temporal Behaviors*. In: *System of Systems Engineering (SoSE)*, IEEE Computer Society.
- [16] Koushik Sen, Abhay Vardhan, Gul Agha & Grigore Roşu (2004): *Efficient Decentralized Monitoring of Safety in Distributed Systems*. *International Conference on Software Engineering*, pp. 418–427.
- [17] Wenchao Zhou, Oleg Sokolsky, Boon Thau Loo & Insup Lee (2009): *DMaC: Distributed Monitoring and Checking*. In: *Runtime Verification 09, LNCS 5779*, Springer, pp. 184–201.