

# Gates for the Kane quantum computer in the presence of dephasing

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In this paper we investigate the effect of dephasing on proposed quantum gates for the solid-state Kane quantum computing architecture. Using a simple model of the decoherence, we find that the typical error in a controlled-NOT gate is  $8.3 \times 10^{-5}$ . We also compute the fidelities of Z, X, swap, and controlled Z operations under a variety of dephasing rates. We show that these numerical results are comparable with the error threshold required for fault tolerant quantum computation.

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## I. INTRODUCTION

One of the most exciting advances in physics has been the development of quantum algorithms [1,2] which outperform their best known classical counterparts. These algorithms are described in the absence of noise, and decoherence. In experiment this will certainly not be the case. In this paper we investigate numerically how a simple model of decoherence affect gates on the Kane quantum computer [3]. The Kane quantum computer is one of a number of promising silicon based quantum computer proposals [4–8].

Dephasing in systems similar to the Kane architecture has been investigated since the introduction of the spin echo technique [9]. The nuclear spin and electronic spin decoherence times of P donors in Si is relatively long [10–17]. For example, recently a time of 60 ms was measured for the electronic dephasing time  $T_2^e$  [18]. Although dephasing times are comparatively long, if left unchecked the accumulated errors introduced by dephasing will destroy coherence in the computation.

Using quantum error correction protocols it may be possible to correct the errors caused by decoherence [19–23]. To successfully reduce the overall error in the system, we must correct errors faster than they accumulate, in a fault-tolerant manner [24]. This consideration leads to an error threshold [25–28]. Typically such a threshold requires the probability of introducing an error in each gate to be below  $1 \times 10^{-4}$  to as low as  $1 \times 10^{-6}$ . For the Kane architecture the exact threshold is still under investigation [29]. In this paper we ask if it is possible for the Kane architecture to achieve this error threshold. To do this we must know how much error is introduced by each of the gates used in the Kane quantum computing architecture.

This paper shows the results of simulations of gates on the Kane architecture in the presence of dephasing. Single qubit gates presented here are similar to those used in nuclear mag-

netic resonance [30,31] for rotations of individual qubits. Voltage fluctuations on the “A” gate and stochastic modeling of the system has been investigated analytically in Refs. [32,33]. The two qubit gates presented here use non-adiabatic pulse schemes [34].

The analysis given here is a direct analogue to that of Fowler *et al.* [35] for the adiabatic controlled-NOT (CNOT) gate. The gates analyzed in this paper are simpler, faster, and potentially higher fidelity than the adiabatic gate [34]. These gates do not rely on complicated pulse shapes, but simply turning on or off the voltages applied. In contrast to adiabatic gates, the timing of these gates could easily be run off a digital clock cycle. In addition, whereas the adiabatic CNOT operation is required to be applied up to three times to create an arbitrary two qubit gate, using nonadiabatic schemes it is possible to create an arbitrary two qubit gate directly. For example, the swap gate analyzed here would require three adiabatic CNOT gates to construct. Using nonadiabatic schemes we are able to construct it in a single pulse sequence [34], which is much faster, simpler, and higher fidelity than the corresponding adiabatic scheme.

We compare each of the gates analyzed to the error threshold for fault tolerant quantum computing. We simulate the master equation for typical values of spin dephasing expected in the Kane architecture. We find that the error in the gates analyzed is less than or comparable to that required for fault tolerant quantum computation.

This paper is organized as follows. In Sec. II we describe the simple model of decoherence which we use and present the master equation for the system. Sec. III presents the results for one qubit gates, including free evolution in Sec. III A, Z rotations in Sec. III B and X rotations in Sec. III C. Two qubit nonadiabatic gates are shown in Sec. IV. These include the CNOT gate in Sec. IV A, and the swap gate and controlled Z gates in Sec. IV B. Finally, conclusions are drawn in Sec. V.

## II. THE MASTER EQUATION

A brief introduction to the Kane quantum computing architecture [3] is given here. The Kane architecture consists of P donor atoms embedded in Si. The orientation of the nuclear spin of each P donor represents one qubit.

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TABLE I. Typical parameters used for numerical calculations.

Description	Term	Value
Unperturbed hyperfine interaction	$A$	$0.1211 \times 10^{-3}$ meV
Hyperfine interaction during $Z$ rotation	$A_z$	$0.0606 \times 10^{-3}$ meV
Hyperfine interaction during $X$ rotation	$A_x$	$0.0606 \times 10^{-3}$ meV
Constant magnetic field strength	$B$	2.000 T
Rotating magnetic field strength	$B_{ac}$	0.0025 T
Hyperfine interaction during interaction	$A_U$	$0.1197 \times 10^{-3}$ meV
Exchange interaction during interaction	$J_U$	0.0423 meV

When placed in a magnetic field applied in the  $z$  direction, Zeeman splitting occurs. This is given by the Hamiltonian

$$H_B = -g_n \mu_n B Z_n + \mu_B B Z_e, \quad (1)$$

where  $Z$  is the Pauli  $Z$  matrix, and the subscripts  $e(n)$  indicate electronic (nuclear) spin. The magnetic field  $B$  may be controlled externally. The application of a resonant rotating magnetic field adds the following terms to the spin Hamiltonian:

$$H_{ac} = -g_n \mu_n B_{ac} [X_n \cos(\omega_{ac} t) + Y_n \sin(\omega_{ac} t)] + \mu_B B_{ac} [X_e \cos(\omega_{ac} t) + Y_e \sin(\omega_{ac} t)]. \quad (2)$$

The electronic spins couple to their corresponding nuclear spins via the hyperfine interaction

$$H_A = A \boldsymbol{\sigma}_e \cdot \boldsymbol{\sigma}_n, \quad (3)$$

where  $A$  is the strength of the interaction. The design of the Kane quantum computer calls for control of the strength of the hyperfine interaction externally by applying appropriate voltages to “ $A$ ” gates. The electronic spins couple to adjacent electrons via the exchange interaction

$$H_J = J \boldsymbol{\sigma}_{e_1} \cdot \boldsymbol{\sigma}_{e_2}, \quad (4)$$

where  $e_1$  and  $e_2$  are two adjacent electrons, and  $J$  is the strength of the exchange interaction which may be controlled externally through the application of voltages to the “ $J$ ” gates.

Altogether the spin Hamiltonian of a two donor system is given by

$$H_s = \sum_{i=1}^2 H_{B_i} + H_{A_i} + H_J + H_{ac_i}. \quad (5)$$

The times and fidelities of the gates naturally depend on exactly which parameters are used to calculate them. For many of the gates in this paper the typical parameters shown in Table I were used. These parameters are similar to the parameters used for the pure state calculations in Ref. [34].

A simple model of decoherence was used for these calculations. There are many different decoherence mechanisms, but our model only considers pure dephasing (without energy relaxation). Whereas dephasing is certainly not the only source of decoherence, it is likely to be the dominant effect on a time scale shorter than the energy relaxation (dissipa-

tion) time,  $T_1$ . For example, Feher and Gere [12] measured  $T_{1_n} > 10$  h for nuclear spin at a temperature of  $T=1.25$  K,  $B=3.2$  T, and  $T_{1_e} \approx 30$  h under similar conditions. In contrast, experimentally measured times for  $T_2$  have been much shorter. Gordon and Bowers [11] measured  $T_{2_e} = 520 \mu\text{s}$  for P:Si at  $T=1.4$  K in isotopically enriched  $^{28}\text{Si}$ . Chiba and Harai [16] have also measured the electronic decoherence times of P:Si, finding a rate of  $T_{2_e} = 100 \mu\text{s}$ . For the nucleus, recent results for the nuclear spin of a  $^{29}\text{Si}$  nucleus show a maximum value of  $T_{2_n} = 25$  s [36].

Recently Tyryshkin *et al.* [18] obtained an experimental measurement of  $T_{2_e} = 14.2$  ms at  $T=8.1$  K and  $T_{2_e} = 62$  ms at  $T=6.9$  K for a donor concentration of  $0.87 \times 10^{15} \text{ cm}^{-3}$  in isotopically pure Si. At milliKelvin temperatures the decoherence time is likely to be even longer. Additionally these measurements were carried out in a bulk doped sample, and in our case we will be considering a specifically engineered sample. Interactions such as exchange and dipole-dipole interactions contribute to the coupling between electrons. In some experiments, such as in Ref. [18], these potentially beneficial coupling have been treated as sources of decoherence, but in the operation of a quantum computer these interactions can either be decoupled or used to generate entanglement, useful for quantum computation. Hence, it is expected that  $T_{2_e}$  may even be longer than those reported in Ref. [18]. Nevertheless, we use the value of 60 ms as a conservative estimate for electronic dephasing time. We expect the nuclear dephasing times to be several orders of magnitude bigger than electronic dephasing times. We choose the following parameters to be typical of P:Si the systems we are considering

$$T_{2_e} = 60 \text{ ms}, \quad (6)$$

$$T_{2_n} = 1 \text{ s}. \quad (7)$$

The typical errors presented in the tables contained in the next three sections are evaluated at these typical dephasing times.

The simple decoherence model we consider corresponds to the master equation

$$\dot{\rho} = -\frac{i}{\hbar} [H_s, \rho] - \mathcal{L}[\rho], \quad (8)$$

where the dephasing terms are given by

$$\mathcal{L}[\rho] = \sum_{i=1}^2 \Gamma_e [Z_{e_i}, [Z_{e_i}, \rho]] + \Gamma_n [Z_{n_i}, [Z_{n_i}, \rho]]. \quad (9)$$

Characteristic dephasing rates,  $\Gamma_{2_e}$  and  $\Gamma_{2_n}$ , are related to the dephasing rates by the equations

$$T_{2_e} = \frac{1}{4\Gamma_e}, \quad (10)$$

$$T_{2_n} = \frac{1}{4\Gamma_n}. \quad (11)$$

We define fidelity (and therefore error) in terms of the actual state after applying an operation  $\rho$  and the intended state after that operation,  $\rho'$ . Due to systematic errors and decoherence these states will not necessarily be the same. When comparing against a pure state  $\rho'$ , the fidelity  $F$  of an operation is defined as

$$F(\rho, \rho') = \text{Tr}(\rho\rho'). \quad (12)$$

Error is defined in terms of fidelity

$$E(\rho, \rho') = 1 - F(\rho, \rho'). \quad (13)$$

Typically we would like to know the greatest error possible for any input state. This is a computationally difficult problem. In the results that follow, the approach taken is to calculate the fidelity for each of the computational basis states, and each of the input states which would ideally generate a Bell state. This has two main benefits. The first is that a high fidelity indicates that the gate is successfully creating or preserving entanglement. The second is that Bell states are superposition states, which are susceptible to dephasing. We also calculated the effect of each gate on the four Bell input states for the CNOT gate. For typical parameters, Bell input states give similar fidelities to those shown in this paper.

Throughout this paper we will use the states  $|0\rangle$  and  $|1\rangle$  to represent the nuclear spin up and spin down states, respectively. We will use the  $|\uparrow\rangle$  and  $|\downarrow\rangle$  to represent electronic spin up and spin down states, respectively.

### III. ONE QUBIT GATES

#### A. Free evolution

The spin of an isolated nucleus undergoing Larmor precession in the presence of a magnetic field  $B$  [37]. In this case it is easy to solve the master equation with dephasing exactly. Considering only the nuclear spin, we have

$$H_s = g_n \mu_n B Z_n. \quad (14)$$

The decoherence terms has only the single term

$$\mathcal{L}[\rho] = \Gamma_n [Z_n, [Z_n, \rho]] \quad (15)$$

$$= 2\Gamma_n (\rho - Z_n \rho Z_n). \quad (16)$$

In the rotating frame, the master equation has the solution

$$\rho(t) = \begin{bmatrix} \rho_{00}(0) & \rho_{01}(0)e^{-4\Gamma_n t} \\ \rho_{10}(0)e^{-4\Gamma_n t} & \rho_{11}(0) \end{bmatrix}. \quad (17)$$

This has the effect of exponentially decaying the off diagonal terms of the density matrix, but leaves the diagonal components unchanged. For a single isolated nuclear spin, the simple model has no effect on eigenstates of  $Z_n$  (i.e., there is no relaxation process for these states). In contrast, it has a dramatic influence on superposition states whose off diago-

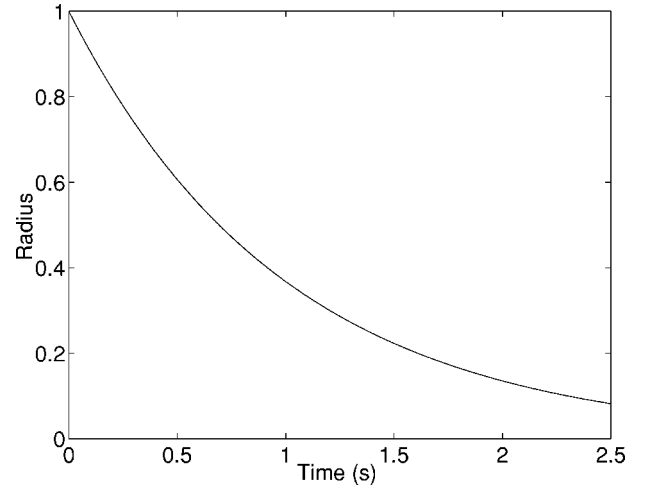


FIG. 1. Effect of dephasing on the free evolution of a single embedded P atom in the Kane architecture.

nal terms decay exponentially (i.e., dephasing). Two such states are

$$|+\rangle = \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle),$$

$$|-\rangle = \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle).$$

We can easily calculate the expectation value of the Pauli  $X$  matrix,  $\langle X \rangle$ , for the  $|+\rangle$  state

$$\langle X \rangle = \text{Tr}(X\rho), \quad (18)$$

$$= \exp(-4\Gamma_n t). \quad (19)$$

For a single nuclear spin coupled to an electronic spin via the hyperfine interaction the Hamiltonian is given by

$$H_s = H_B + H_A. \quad (20)$$

We assume that electron is initially polarized by the large magnetic field  $B$ . The evolution of this Hamiltonian was calculated for their typical values [given in Eq. (6), Eq. (7), and Table I]. The fidelity after different times is shown in Fig. 1. This figure shows the Bloch sphere radius, given by

$$r = \sqrt{\langle X \rangle^2 + \langle Y \rangle^2 + \langle Z \rangle^2}. \quad (21)$$

A pure state has a radius of one, and a radius of less than one indicates a mixed state. An initial state of  $|\downarrow+\rangle$  was used. The radius decays at a rate governed by the nuclear decoherence time, which in this case is  $T_{2n} = 1$  s. As is expected, this decay is the same as the well known solution to the Bloch equations.

#### B. Z rotations

$Z$  rotations on the Kane architecture may be performed by varying the Larmor precession frequency of a single qubit [34]. Graphs showing the error in the  $Z$  gate at different dephasing rates are shown in Figs. 2 and 3.

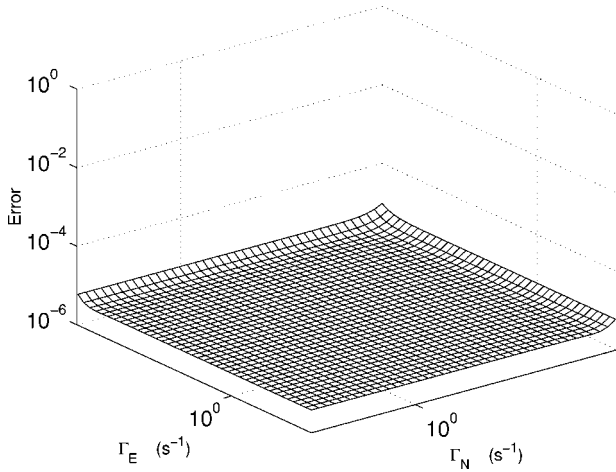


FIG. 2. Error in the Z gate for  $|\downarrow 0\rangle$  initial state at differing rates of dephasing.

Figure 2 shows the error at different rates of decoherence for both electrons  $\Gamma_e$  and nuclei  $\Gamma_n$  for a single qubit in the  $|\downarrow 0\rangle$  state. The calculated error does not significantly depend on the electronic or nuclear dephasing rate. The pure  $|\downarrow 0\rangle$  state is not affected by decoherence terms in the master equation. Therefore the only effect of dephasing occurs when the hyperfine interaction rotates this state. The effect of dephasing on this state is negligible.

The error in the Z operation for an initial state of  $|\downarrow 0\rangle$ , is primarily due to systematic error of  $3.8 \times 10^{-6}$ . This error is due to the hyperfine interaction coupling between electrons and nuclei. This allows a small probability of finding the electrons in an excited state. At typical rates of decoherence for the short duration of a Z rotation (approximately 21 ns), systematic error is the dominant effect. For the  $|\downarrow 0\rangle$  state, the typical error is  $3.8 \times 10^{-6}$ .

In the previous section we noted that superposition states  $|\downarrow +\rangle$  and  $|\downarrow -\rangle$  are affected by dephasing terms more than eigenstates of Z. This is illustrated in Fig. 3. Electronic dephasing times have little effect on the overall fidelity. As

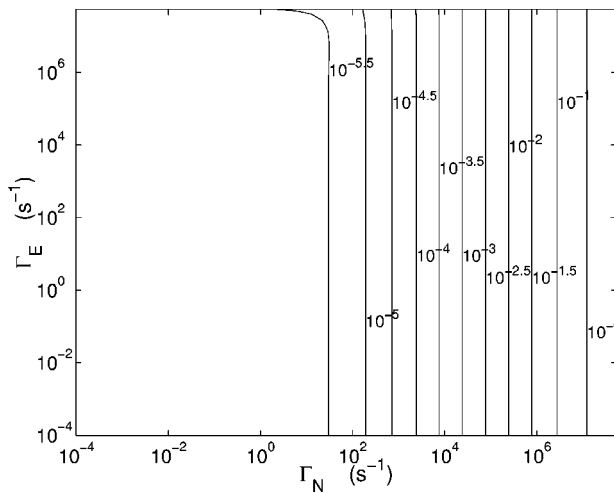


FIG. 3. Contour plot showing the error in the Z gate for  $|\downarrow +\rangle$  initial state evolving at differing rates of dephasing. Contour lines connect, and are labeled by points of equal error.

TABLE II. Summary of Z gate error.

State	Systematic error	Typical error
$ 0\rangle$	$3.8 \times 10^{-6}$	$3.8 \times 10^{-6}$
$ +\rangle$	$1.9 \times 10^{-6}$	$1.9 \times 10^{-6}$
Maximum	$3.8 \times 10^{-6}$	$3.8 \times 10^{-6}$

the nuclear dephasing rate increases, the fidelity decreases with a maximum error of approximately 0.5. This indicates all quantum coherence was lost and we are in classical mixture of the states  $|\downarrow 0\rangle$  and  $|\downarrow 1\rangle$ . For typical rates of dephasing [given in Eqs. (6) and (7)], the error is found to be  $1.9 \times 10^{-6}$  which is largely due to systematic error.

The maximum error, for typical rates of dephasing, of any of the states tested for the Z gate is  $3.8 \times 10^{-6}$ . This error is largely due to systematic effects rather than dephasing. This error suggests it is theoretically possible to do a Z rotation with an error of less than the  $1 \times 10^{-4}$  limit suggested for fault tolerant quantum computing. The results for the Z gate are summarized in Table II.

### C. X rotations

X rotations are performed using a resonant magnetic field  $B_{ac}$ . The error in the X rotation was found at different rates of dephasing. This is shown for the two basis states in Figs. 4 and 5. Evident on both of these graphs is a valley which levels out at a minimum error. This error is primarily due to systematic error in the gate. For the  $|\downarrow 0\rangle$  initial state the systematic error is found to be  $2.3 \times 10^{-6}$ , and for the  $|\downarrow 1\rangle$  initial state the systematic error is found to be  $4.9 \times 10^{-6}$ . As the nuclear dephasing rates  $\Gamma_{2n}$  increases, the fidelity of the operation decreases. The fidelity of the operation drops to 0.0 indicating that, in the limit of large dephasing rates, the rotating magnetic field does not have the desired effect of a resonant magnetic field.

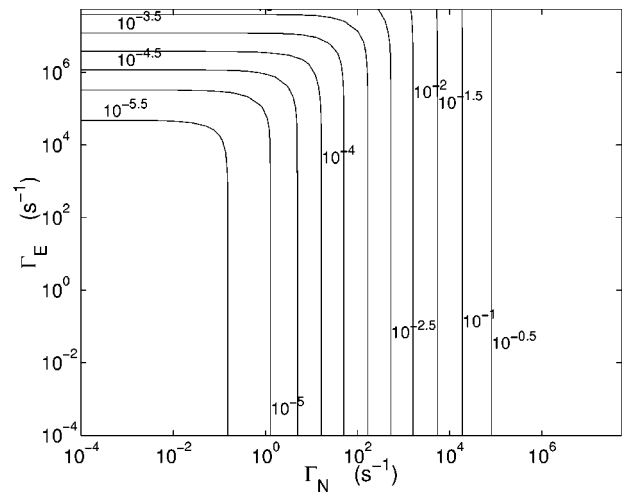


FIG. 4. Contour plot showing the error in the X gate for the  $|\downarrow 0\rangle$  initial state evolving at differing rates of dephasing. Contour lines connect, and are labeled by points of equal error.

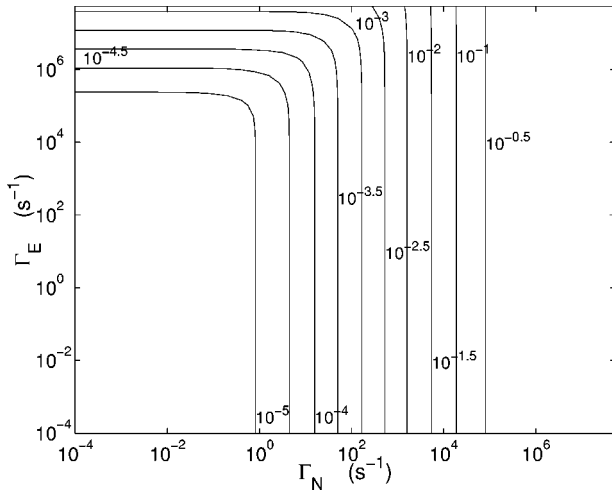


FIG. 5. Contour plot showing the error in the  $X$  gate for the  $|\downarrow\downarrow\rangle$  initial state at differing rates of dephasing. Contour lines connect and are labeled by points of equal error: (a)  $|00\rangle$  initial state, (b)  $|01\rangle$  initial state, (c)  $|10\rangle$  initial state, and (d)  $|11\rangle$  initial state.

In Figs. 4 and 5, we see that the fidelity of the  $X$  rotation depends weakly on the electronic dephasing rate  $\Gamma_e$ . The solutions to the equations under these conditions show that the principal cause of error is the electron becoming excited to a higher energy level.

Under typical conditions, the  $|\downarrow\downarrow\rangle$  initial state has an error of  $3.8 \times 10^{-6}$ . The  $|\downarrow\downarrow\rangle$  state has an error of  $6.4 \times 10^{-6}$ . In this operation dephasing has a much more important role than in  $Z$  rotations. One reason for this is that an  $X$  rotation takes longer, approximately  $6.4 \mu\text{s}$ .

The error in the  $X$  gate is summarized in Table III. The maximum error from the two basis states tested for the  $X$  gate was  $6.4 \times 10^{-6}$ . The error induced in this operation is less than a threshold of  $1 \times 10^{-4}$  required for fault tolerant quantum computing.

## IV. TWO QUBIT GATES

### A. The CNOT gate

The CNOT gate is specified by

$$\Gamma_1 X = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}, \quad (22)$$

which may be created using the steps specified in Ref. [34]. In the following discussion of two qubit gates, if not explic-

TABLE III. Summary of  $X$  gate error.

State	Systematic error	Typical error
$ 0\rangle$	$2.3 \times 10^{-6}$	$3.8 \times 10^{-6}$
$ 1\rangle$	$4.9 \times 10^{-6}$	$6.4 \times 10^{-6}$
Maximum	$4.9 \times 10^{-6}$	$6.4 \times 10^{-6}$

itly stated, all initial electron spin states are assumed to be  $|\downarrow\downarrow\rangle$ .

The error in the CNOT operation in the presence of dephasing was found by numerically solving the master equation [Eqs. (8) and (9)] for the appropriate pulse sequence [34]. For different rates of dephasing, different fidelities are obtained. Fidelities were calculated for each of the four computational basis states, and for the evolution leading to the four Bell states. Fidelities for the four computational basis states are shown in Fig. 6.

Each of the states shows a minimum error when the rates of electronic and nuclear dephasing are slow. The remaining error is due to systematic error in the gate. Some sources of systematic error for the CNOT operation include off-resonant effects, excitation into higher electronic energy levels, and imperfections in the pulse sequences (such as the breakdown of the second order approximations used to derive appropriate times for pulse sequence [34]). Systematic error for each of the states  $|00\rangle$ ,  $|01\rangle$ ,  $|10\rangle$ , and  $|11\rangle$  are  $4.0 \times 10^{-5}$ ,  $2.6 \times 10^{-5}$ ,  $1.9 \times 10^{-5}$ , and  $2.9 \times 10^{-5}$ , respectively. For evolution starting in an initial Bell state, we find the systematic error is  $3.5 \times 10^{-5}$ ,  $3.4 \times 10^{-5}$ ,  $1.9 \times 10^{-5}$ , and  $2.6 \times 10^{-5}$ . Systematic error, for states resulting in a Bell state are shown in Table IV.

In each of the four states, as the dephasing rate increases, the fidelity decreases. In the limit of large dephasing rates, the computational basis states tend to stay in their original states. This is particularly evident from the graphs of  $|00\rangle$  and  $|01\rangle$  which have higher fidelity (lower error) at high dephasing rates. In contrast, states  $|10\rangle$  and  $|11\rangle$  have lower fidelities (higher error) at high rates of dephasing. For example, the  $|00\rangle$  state stays in the  $|00\rangle$  state after the CNOT operation. At such high rates of dephasing, not even the single qubit rotations described in Sec. III C apply, which form part of the CNOT gate operation. For these states quantum coherence has been lost. When we consider, for example, the state  $|\psi\rangle = 1/\sqrt{2}(|00\rangle + |01\rangle)$  we find that at a high dephasing rate of  $\Gamma_n = \Gamma_e = 54 \times 10^6 \text{ s}^{-1}$  the error of the gate is 0.5. In this case, quantum coherence has been lost between the two states, and the qubit evolves to a completely mixed state.

Electronic dephasing rates play a much bigger role in two qubit gates than in single qubit gates. For the typical dephasing times,  $T_2^e$  and  $T_2^n$  [given in Eqs. (6) and (7)], we find the error for the states  $|00\rangle$ ,  $|01\rangle$ ,  $|10\rangle$ , and  $|11\rangle$  are  $8.3 \times 10^{-5}$ ,  $6.8 \times 10^{-5}$ ,  $6.2 \times 10^{-5}$ , and  $7.2 \times 10^{-5}$  respectively. The error for an initial state of one of the four Bell states is found to be  $6.0 \times 10^{-5}$ ,  $6.0 \times 10^{-5}$ ,  $4.4 \times 10^{-5}$ , and  $5.1 \times 10^{-5}$ . Typical errors for states resulting in a Bell state are shown in Table IV. This implies that under our very simple decoherence model, the maximum error in the CNOT gate in any basis state is  $8.3 \times 10^{-5}$ . This is only marginally under the threshold of  $1 \times 10^{-4}$  required for fault tolerant quantum computation.

Errors for each of the computational basis states are shown in Table IV, and the maximum error of any of the four computational basis states or states with evolution leading to a Bell state is plotted in Fig. 7.

These results are directly analogous to calculations made for the adiabatic CNOT gate [35]. We have used the same noise model as was used in their numerical simulations. In

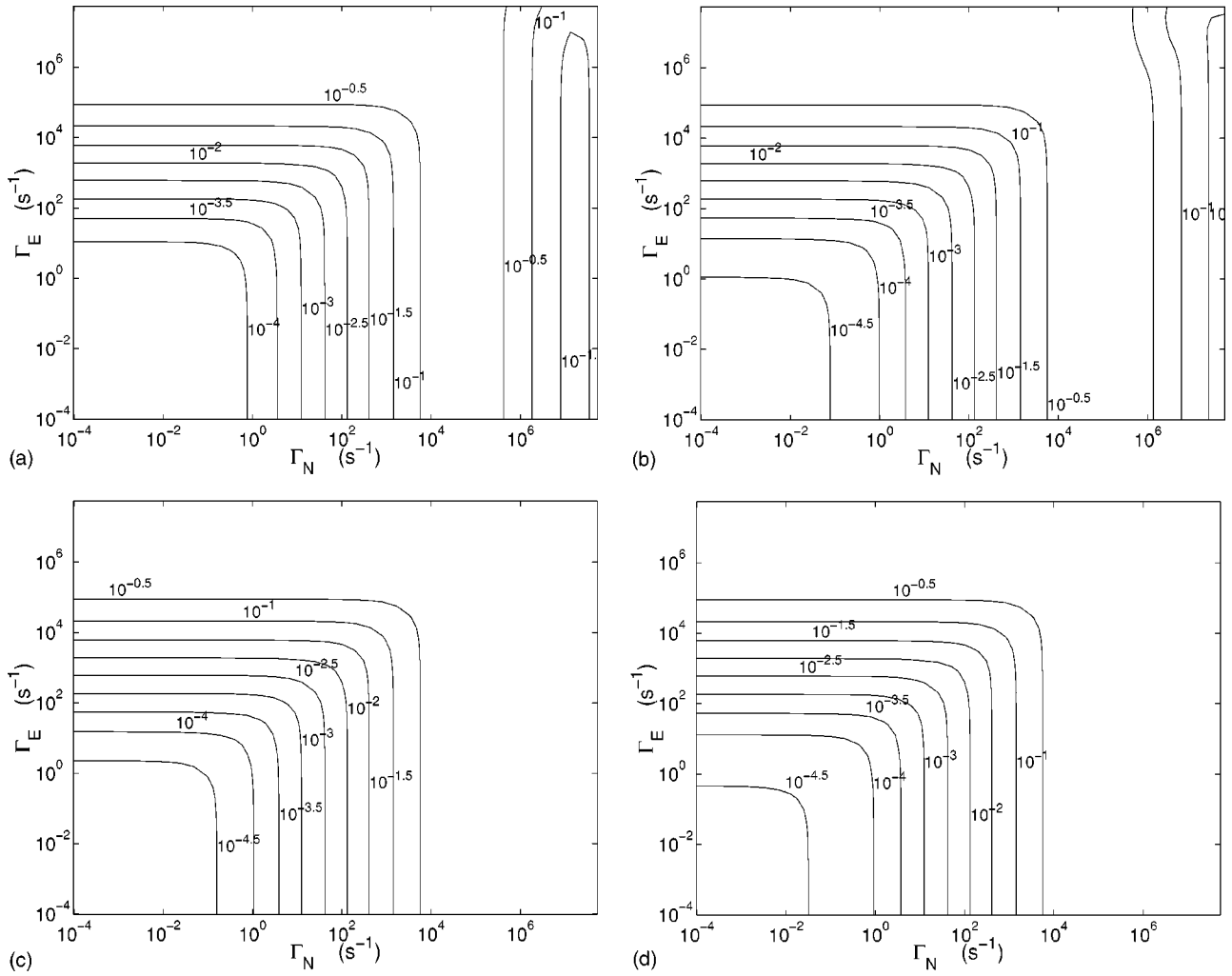


FIG. 6. Contour plots showing error in the CNOT operation for different rates of electronic and nuclear dephasing. Contour lines connect and are labeled by points of equal error.

calculations for the adiabatic CNOT gate at decoherence times, for electronic and nuclear dephasing rates of  $T_{2e} = 200 \mu s$  and  $T_{2n} = 10 s$ , respectively, the maximum error of any of the four basis states for the adiabatic CNOT gate was found to be ‘just over  $10^{-3}$ ’ [35]. In comparison, for the same conditions, we find that the maximum error in the nonadia-

TABLE IV. Summary of CNOT gate fidelities.

State	Systematic error	Typical error
$ 00\rangle$	$4.0 \times 10^{-5}$	$8.3 \times 10^{-5}$
$ 01\rangle$	$2.6 \times 10^{-5}$	$6.8 \times 10^{-5}$
$ 10\rangle$	$1.9 \times 10^{-5}$	$6.2 \times 10^{-5}$
$ 11\rangle$	$2.9 \times 10^{-5}$	$7.2 \times 10^{-5}$
$ 00\rangle +  11\rangle$	$2.9 \times 10^{-5}$	$7.0 \times 10^{-5}$
$ 00\rangle -  11\rangle$	$3.2 \times 10^{-5}$	$7.3 \times 10^{-5}$
$ 01\rangle +  10\rangle$	$3.1 \times 10^{-5}$	$7.2 \times 10^{-5}$
$ 01\rangle -  10\rangle$	$2.3 \times 10^{-5}$	$6.4 \times 10^{-5}$
Maximum	$4.0 \times 10^{-5}$	$8.3 \times 10^{-5}$

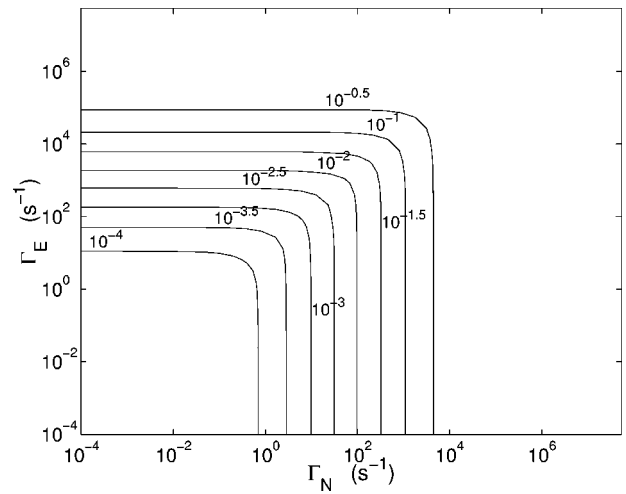


FIG. 7. Contour plot of the maximum error in basis and Bell output states of the CNOT gate shown as a function of electronic and nuclear dephasing rates. Contour lines connect and are labeled by points of equal error.

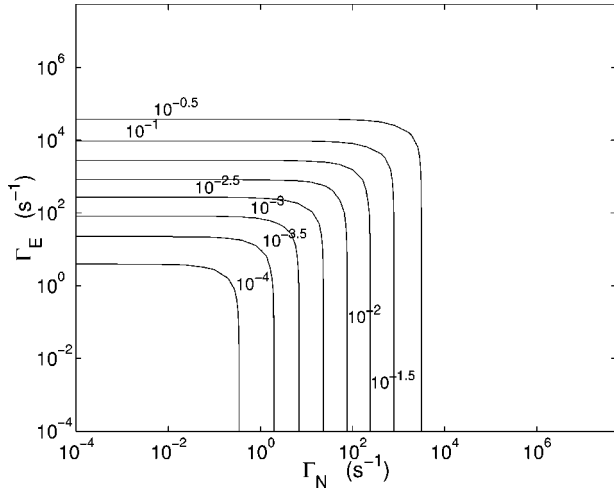


FIG. 8. Contour plot of the maximum error in basis states and Bell output states of the swap gate shown as a function of electronic and nuclear dephasing rates. Contour lines connect and are labeled by points of equal error.

batic gate is  $7 \times 10^{-3}$ . Under these conditions both nonadiabatic and adiabatic gates give similar fidelities.

Where does this error come from at the dephasing rate specified above, for the nonadiabatic CNOT gate? The error in an X rotation under these conditions is  $3 \times 10^{-6}$  and the error in the entangling operation of the gate  $U_m(\pi/4)$ , is  $4 \times 10^{-3}$  and therefore it is clear that the two qubit entangling operation is the major source of error. When electronic decoherence times are short, any electron mediated operations will be affected by this decoherence. By increasing the exchange interaction strength, the time required for the electron mediated operation may be reduced. For example, at a strength of  $J=0.0529$  meV the error decreases to  $1 \times 10^{-3}$  and the maximum error in the CNOT gate is also  $1 \times 10^{-3}$  for any of the computational basis states.

The advantage of nonadiabatic gates over adiabatic gates is that the pulse schemes required are much simpler, faster,

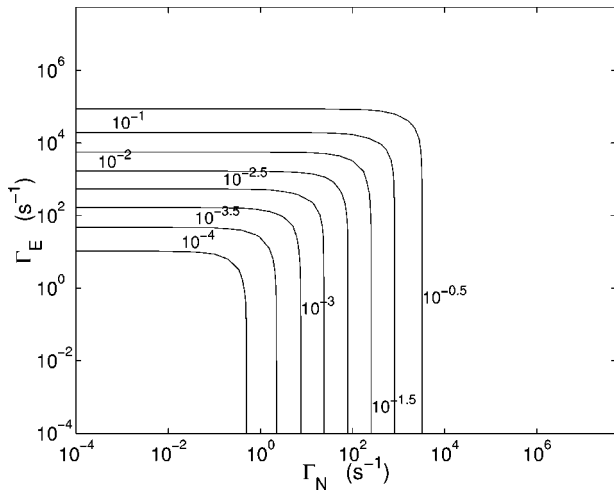


FIG. 9. Contour plot of the maximum error in basis states and Bell output states of the  $\Gamma_1 Z$  gate shown as a function of electronic and nuclear dephasing rates. Contour lines connect and are labeled by points of equal error.

TABLE V. Summary of swap gate error.

State	Systematic error	Typical error
$ 00\rangle$	$3.9 \times 10^{-5}$	$9.0 \times 10^{-5}$
$ 01\rangle$	$1.4 \times 10^{-5}$	$7.9 \times 10^{-5}$
$ 10\rangle$	$1.6 \times 10^{-5}$	$8.0 \times 10^{-5}$
$ 11\rangle$	$3.8 \times 10^{-5}$	$8.9 \times 10^{-5}$
$ 00\rangle +  11\rangle$	$5.3 \times 10^{-5}$	$1.4 \times 10^{-4}$
$ 00\rangle -  11\rangle$	$7.4 \times 10^{-5}$	$1.6 \times 10^{-4}$
$ 01\rangle +  10\rangle$	$1.7 \times 10^{-5}$	$1.0 \times 10^{-4}$
$ 01\rangle -  10\rangle$	$1.5 \times 10^{-5}$	$1.0 \times 10^{-4}$
Maximum	$7.4 \times 10^{-5}$	$1.6 \times 10^{-4}$

and as at the conditions considered in Ref. [35], the two schemes have approximately the same fidelity. Considering the electronic decoherence times measured in Ref. [18] it is likely that electronic dephasing rates are not as large as considered in Ref. [35]. At lower rates of dephasing, we approach the systematic error, which may be smaller for nonadiabatic gates than adiabatic gates [34]. Another distinct advantage of the nonadiabatic gates is that *any* two qubit gate may be made. This allows us to construct two qubit gates directly (such as the swap gate), which are faster and higher fidelity than expressing them as combinations of CNOT gates and single qubit rotations. If the Kane computer was being run from a digital clock cycle, nonadiabatic two qubit gates could be controlled at discrete times, and do not require the continuous and sophisticated pulse shapes required for adiabatic gates operating at this speed and fidelity.

### B. The swap gate and controlled-Z gate

Similar calculations to those calculated for the CNOT gate were carried out for the swap gate and the controlled-Z gate. The swap gate is specified by the matrix

$$U_{\text{Swap}} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}. \quad (23)$$

The controlled-Z gate is specified by the matrix

TABLE VI. Summary of controlled-Z gate error.

State	Systematic error	Typical error
$ 00\rangle$	$3.0 \times 10^{-5}$	$8.1 \times 10^{-5}$
$ 01\rangle$	$1.5 \times 10^{-5}$	$6.6 \times 10^{-5}$
$ 10\rangle$	$1.1 \times 10^{-5}$	$6.2 \times 10^{-5}$
$ 11\rangle$	$3.8 \times 10^{-5}$	$8.9 \times 10^{-4}$
$ 00\rangle +  11\rangle$	$3.6 \times 10^{-5}$	$9.4 \times 10^{-5}$
$ 00\rangle -  11\rangle$	$3.3 \times 10^{-5}$	$9.2 \times 10^{-5}$
$ 01\rangle +  10\rangle$	$1.7 \times 10^{-5}$	$7.5 \times 10^{-5}$
$ 01\rangle -  10\rangle$	$1.1 \times 10^{-5}$	$7.0 \times 10^{-5}$
Maximum	$3.8 \times 10^{-5}$	$9.4 \times 10^{-5}$

TABLE VII. Times and errors for each of the gates investigated.

Gate	Typical error	Systematic error	Time ( $\mu\text{s}$ )
Z	$3.8 \times 10^{-6}$	$3.8 \times 10^{-6}$	0.02
X	$6.4 \times 10^{-6}$	$4.9 \times 10^{-6}$	6.4
CNOT	$8.3 \times 10^{-5}$	$4.0 \times 10^{-5}$	16.0
swap	$1.6 \times 10^{-4}$	$7.4 \times 10^{-5}$	19.2
$\Gamma_1 Z$	$9.4 \times 10^{-5}$	$3.8 \times 10^{-5}$	16.1

$$\Gamma_1 Z = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -1 \end{bmatrix}. \quad (24)$$

The circuit which may be used to create the swap gate may be found in Ref. [34].

The master equation was solved numerically for each pulse sequence. The error in the swap gate was calculated for each basis state, and each state whose output state is a Bell state. Note that for these two gates the states which give Bell states as output are themselves Bell states. A separate simulation was completed for each combination of nuclear and electronic dephasing times. The maximum error of any of the basis states has been plotted in Fig. 8 for the swap gate, and Fig. 9 for the control Z gate.

Similar features that were evident for the CNOT gate are visible in these figures. The corresponding errors are shown in Table V for the swap gate, and Table VI for the control-Z gate.

## V. CONCLUSION

In conclusion, we have investigated the effect of dephasing on the Kane quantum computing architecture. We used a

simple model of decoherence and investigated how this model affected proposed gate schemes on the Kane quantum computer. For typical decoherence rates [given in Eqs. (6) and (7)], these results are summarized in Table VII.

Each of the errors, for typical rates of dephasing, found here are close to the error threshold required for fault tolerant quantum computation. If the temperature is lowered, and coupling between qubits is not considered a decoherence process as in Ref. [18], it is likely that the typical decoherence times for the Kane architecture may be further reduced, and therefore unambiguously under the threshold required for fault tolerant quantum computation.

Construction and operation of the Kane quantum computer is extremely challenging. In the actual physical system there will undoubtedly be noise and decoherence processes not considered in our simple physical model. This substantial effort would never be able to achieve its ultimate goal of a working quantum computer if there were fundamental reasons why such a computer could not operate. In this paper we investigated the one such effect on proposed gates for the Kane quantum computer. Our simulations indicate that errors due to dephasing, the dominant form of decoherence in the Kane architecture, do not rule out fault tolerant quantum computation.

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