# Estimation of Sparse MIMO Channels with Common Support 

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#### Abstract

We consider the problem of estimating sparse communication channels in the MIMO context. In small to medium bandwidth communications, as in the current standards for OFDM and CDMA communication systems (with bandwidth up to 20 MHz , such channels are individually sparse and at the same time share a common support set. Since the underlying physical channels are inherently continuous-time, we propose a parametric sparse estimation technique based on finite rate of innovation (FRI) principles. Parametric estimation is especially relevant to MIMO communications as it allows for a robust estimation and concise description of the channels.

The core of the algorithm is a generalization of conventional spectral estimation methods to multiple input signals with common support. We show the application of our technique for channel estimation in OFDM (uniformly/contiguous DFT pilots) and CDMA downlink (Walsh-Hadamard coded schemes). In the presence of additive white Gaussian noise, theoretical lower bounds on the estimation of sparse common support (SCS) channel parameters in Rayleigh fading conditions are derived. Finally, an analytical spatial channel model is derived, and simulations on this model in the OFDM setting show the symbol error rate (SER) is reduced by a factor 2 ( 0 dB of SNR) to 5 (high SNR) compared to standard non-parametric methods e.g. lowpass interpolation.


Index Terms—Channel estimation, MIMO, OFDM, CDMA, finite rate of innovation.

## I. INTRODUCTION

MULTIPLE input multiple output (MIMO) antenna wireless systems enable significant gains in both throughput and reliability [1]-[4] and are now incorporated in several commercial wireless standards [5], [6]. However, critical to realizing the full potential of MIMO systems is the need for accurate channel estimates at the receiver, and, for certain schemes, at the transmitter as well. As the number of transmit antennas is increased, the receiver must estimate proportionally more channels, which in turn increases the pilot overhead and tends to reduce the overall MIMO throughput gains [7].

[^0]To reduce this channel estimation overhead, the key insight of this paper is that most MIMO channels have an approximately sparse common support (SCS). That is, the channel in each transmit-receive (TX-RX) antenna pair can be modeled as a discrete multipath channel, with the relative time delays being common across different TX-RX pairs. The commonality across the different antenna pairs reduces the overall number of degrees of freedom to estimate, which can in turn be used to reduce the pilot overhead or improve the channel estimate. Also, in communication systems that depend on channel state feedback from the RX to the TX, the SCS model may enable a more compressed representation.

To exploit the SCS property of MIMO channels, we proposed [8], [9] a variant along [10] of the finite rate of innovation (FRI) framework, originally developed in [11]. The present manuscript is a recapitulation and augmented version of [8], [9].

The method, which we call SCS-FRI, uses classical spectral estimation techniques such as Prony's method, ESPRIT and Cadzow denoising to recover the delay positions in frequency domain. The method is computationally simple, and our simulations demonstrate excellent performance in practical scenarios. The proposed SCS-FRI algorithm applies immediately to channel estimation in multi-output OFDM communication with contiguous or uniformly scattered DFT pilots. Interestingly enough it can be used on other modulation schemes provided a suitable pilot layout. The Walsh-Hadamard transform (WHT), used in CDMA downlink channel among others, qualifies if one controls the pilots layout in the WHT domain.

We also derive a simple scalar formula for the CramérRao bound on the estimation of separable ToAs, and also point to a more general result by Yau and Bresler [12]. Both bounds are extended to Rayleigh fading SCS channels to lower bound the expected estimation error in fading conditions. Our simulations indicate the proposed SCS-FRI method is close to this bound at high SNRs.

## A. SCS MIMO Models

Due to the physical properties of outdoor electromagnetic propagation, wireless channels are often modeled as having a channel impulse response (CIR) that is sparse in the sense that they contain few significant paths [13]. With multiple antennas, the CIR measured at different antennas share a common support, i.e. the times of arrival (ToA) at different antennas are similar while the paths amplitudes and phases are distinct. An example of sparse common support (SCS) channels is shown in Fig. 1. The SCS channel model is usually


Fig. 1. (a) Transmission over a bandlimited medium with two scatterers and $P$ receiving antennas. (b) The $P$ channels contain two paths arriving at the same time up to $\pm \varepsilon$, and are thus no exact SCS channels for $\varepsilon>0$.

TABLE I
Channel Bandwidth in Popular Wireless Systems

| System | Code | Bandwidth $B$ | Resolvable <br> distancec/B |
| :--- | :--- | ---: | ---: |
| DVB-T [44] | DFT | $5-8 \mathrm{MHz}$ | $38-60 \mathrm{~m}$ |
| IS-95 [54] | WHT | 1.25 MHz | 240 m |
| 3GPP LTE [6] | DFT | $1.4-20 \mathrm{MHz}$ | $15-215 \mathrm{~m}$ |
| UWB | - | $>500 \mathrm{MHz}$ | $<60 \mathrm{~cm}$ |

assumed in the literature, though its physical motivations are not always put forth.

It is important to note that the sparsity and common support assumptions only hold with respect to the channel bandwidth $B$ and the SNR of the channel. Indeed, in the presence of noise, the spatial resolution of channel paths is limited by the inverse bandwidth $1 / B$, even if one knows exactly which parametric model the signal obeys. At medium to low SNR, $1 / 10^{\text {th }}$ of the inverse bandwidth is a reasonable resolution to shoot for. The limited resolution has the effect of clustering paths from a single scatterer into a single path (promoting sparsity), and the small shift in the ToA due to the distance between antennas becomes negligible (promoting common support). Table I gives the channel bandwidth of several modern standards and $c / B$ which is the distance travelled by an electromagnetic wave in a time lapse equal to the inverse bandwidth.

In the presence of highly diffusive multipaths, this simple sparsity model does not hold as outlined by Berger et al in [14]. Therefore, care shall be taken when applying this model.

## B. Related Work

In OFDM systems, the majority of commercial channel estimators often simply perform some form of linear filtering or interpolation of the pilot symbols [15], [16]. Such nonparametric techniques are computationally very simple, but fundamentally cannot exploit the common sparsity in MIMO channel models. Since the phases and magnitudes are generally independent on the paths on different antenna pairs, the frequency response of sparse common support (SCS) channels are not correlated in any simple manner that can be exploited by basic linear interpolation of pilots.

TABLE II
CHANNEL ESTIMATION METHODS ARE NATURALLY CLASSIFIED IN TERMS OF THE CHANNEL PROPERTIES THEY EXPLOIT.

|  | Algorithm | Exploited channel properties |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  |  | Short delay-spread | Sparsity | Common support |
|  | Lowpass | $\checkmark$ |  |  |
| - ${ }^{0}$ | FRI |  | $\checkmark$ |  |
| ${ }_{3}^{G} 0$ | SCS-FRI |  | $\checkmark$ | $\checkmark$ |
| - | Lowpass | $\checkmark$ |  |  |
| 它 | FRI | $\checkmark$ | $\checkmark$ |  |
| - | SCS-FRI | $\checkmark$ | $\checkmark$ | $\checkmark$ |

A different line of work has proposed compressed sensing based methods for sparse channel estimation [17]-[20]. In the compressed sensing context, the SCS property is equivalent to joint or group sparsity for which there are several methods including group LASSO [21], [22], group OMP [23], simultaneous OMP [24], belief propagation [25], or more recently [26]. Techniques for mixes of joint and individual sparsity are considered in [27], [28]. All of these compressed sensing methods, however, require that the delay locations are discretized and exact sparsity is achieved only when the true path locations fall exactly on one of the discrete points. With continuous value path locations, each path components will require a number of terms to approximate well, or demand a larger number of dictionary elements to offer a finer discretization.

Approximate joint-sparsity in the CS framework is treated in [14], [29] and the effect of leakage is thoroughly studied in [30]. Since the sparsity models assumed by FRI and CS do not coincide exactly, a fair comparison can only be made on measured CIR, as simulations with the model supposed by a particular method is bound to favors the later. Test on field measurements raise the question of the sparsity level estimation which is not known a priori, also complexity is important from an applicative point of view. These considerations and comparisons on field measurements will be addressed in
another paper under preparation [31].
Another joint estimation problem with FRI signals is studied in [32], where the authors study channels which are jointly sparse over time instead of space. They implicitly use an FRI sampling model, which leads to a similar solution.

## C. Contributions

The contributions of this work are four-folds:

- Extension of classical FRI sampling and estimation to multiple SCS channels (Section II)
- Derivation of simple scalar formulas for the CRB of SCS channels (Section III)
- Application to OFDM and Walsh-Hadamard coded (e.g. CDMA downlink) communications with contiguous or uniformly scattered DFT pilots (Section IV)
- Characterization of a precise spatial analytical model for SCS channels (Section V)
The proposed SCS-FRI algorithm stands out compared to FRI or lowpass interpolation as it exploits more channel properties, as indicated in Table II. Lowpass based techniques are a sensible non-parametric approach as they exploit the short delay-spread property. In general, any estimation technique based on uniformly scattered DFT pilots uses this property, as it is a necessary condition to the unicity of the solution.

We conclude our study with numerical simulations showing the efficiency of the SCS-FRI algorithm in a Rayleigh fading scenario, and compare its equalization gain to a standard nonparametric approach, i.e. lowpass interpolation in the DFT domain.

## II. Sparse Common Support FRI: Theory and AlGORITHMS

## A. Problem Formulation

We consider the physical setup described in Fig. 1.(a). A periodic signal of limited bandwidth is transmitted over a multipath channel and uniformly sampled by a receiver with $P$ antennas. This leads to $P$ parallel multipath channels as shown in Fig. 1.(b). The channels either share a common support exactly, in which case they are called exact SCS channels, or approximately, in which case they are called SCS channels (e.g. Fig. 1.(b)).

Consider $P$ exact SCS channels shaped by a kernel $\varphi$, with complex baseband equivalent model:

$$
\begin{equation*}
h_{p}(t)=\sum_{k=1}^{K} c_{k, p} \varphi\left(t-t_{k}\right), \quad c_{k, p} \in \mathbb{C}, t_{k} \in[0 \tau[ \tag{1}
\end{equation*}
$$

where $\varphi(t)$ is the $\tau$-periodic sinc function or Dirichlet kernel:

$$
\begin{equation*}
\varphi(t)=\sum_{k \in \mathbb{Z}} \operatorname{sinc}(B(t-k \tau))=\frac{\sin (\pi B t)}{B \tau \sin \left(\frac{\pi t}{\tau}\right)} \tag{2}
\end{equation*}
$$

The kernel $\varphi$ is considered periodic as the filtering of a periodically padded signal by a linear shift invariant filter. Therefore, linear convolution of the CIR with the shaping kernel becomes circular.

We assume that the bandwidth parameter $B$ satisfies $B=$ $(2 M+1) / \tau$ for $M \geq K$. The paths coefficients $c_{k, p}$ are treated as complex random variables. $N$ measurements $y_{p}[n]$
are acquired at a rate $1 / T=N / \tau$ (with $\tau$ the signal period and $N \geq B \tau=2 M+1$ ) and corrupted by AWGN

$$
\begin{equation*}
y_{p}[n]=h_{p}[n]+q_{p}[n] \quad n \in\{0, \ldots, N-1\} \tag{3}
\end{equation*}
$$

where $\boldsymbol{q}_{p} \sim \mathcal{N}_{\mathbb{C}}\left(\mathbf{0}, \sigma^{2} \mathbb{I}\right)$ if the channel is complex-valued or $\boldsymbol{q}_{p} \sim \mathcal{N}\left(\mathbf{0}, \sigma^{2} \mathbb{I}\right)$ if real-valued. In the DFT domain, the received signal is:

$$
\begin{equation*}
\widehat{y}_{p}[m]=\widehat{\varphi}[m] \cdot \sum_{k=1}^{K} c_{k, p} W^{m t_{k}}+\widehat{q}_{p}[m] . \tag{4}
\end{equation*}
$$

where $W=e^{-2 \pi j / \tau}$ and $\widehat{\varphi}[m]=1 /(2 M+1)$ for $|m| \leq M$ and is zero otherwise. The goal is to estimate the support $\left\{t_{k}\right\}_{k=1 \ldots K}$ and the paths amplitudes $\left\{c_{k, p}\right\}_{k=1 \ldots K, p=1 \ldots P}$ from the $N P$ samples collected in (3). Once the support is known, estimation of the path amplitudes is simple linear algebra as seen in (4).

## B. Support Recovery from Baseband DFT Coefficients

We start from (4). The DFT samples $\widehat{y}_{p}[m]$ in the baseband ( $|m| \leq M$ ) are the DFT coefficients of the channel corrupted by some Gaussian noise.

The noiseless DFT coefficients of a $K$-multipath channel have the well-known and interesting property to form a linear recurrent sequence of order $K+1$, i.e., any coefficient $\widehat{h}_{p}[m]$ ( $m \geq-M+K$ ) can be expressed as a unique linear combination of the $K$ previous DFT coefficients common to all indices $m$ :

Lemma 1. Given $\widehat{h}_{p}[m]=\sum_{k=1}^{K} c_{k, p} W^{m t_{k}}$ for $m=-M+$ $K, \ldots, M$ and $t_{i} \neq t_{j}, \forall i \neq j$, there exists a unique set of coefficients $\left\{f_{k}\right\}_{k=1, \ldots, K}$ such that:
$\widehat{h}_{p}[m]=f_{1} \widehat{h}_{p}[m-1]+f_{2} \widehat{h}_{p}[m-2]+\cdots+f_{K} \widehat{h}_{p}[m-K]$ where $x^{K}-f_{1} x^{K-1}-\cdots-f_{K-1} x-f_{K}$ is the polynomial with roots $\left\{W^{t_{k}}\right\}_{k=1, \ldots, K}$.

Proof: A linear recursion of degree $K$ can be written as:

$$
\begin{equation*}
x_{n}=f_{1} x_{n-1}+\cdots+f_{K} x_{n-K}, f_{K} \neq 0 \tag{5}
\end{equation*}
$$

Its characteristic equation is:

$$
\begin{equation*}
x^{K}-f_{1} x^{K-1}-\cdots-f_{K-1} x-f_{K}=0 \tag{6}
\end{equation*}
$$

If $\lambda_{x}$ is a solution of (6) then multiplying both sides of the equation by $\lambda_{x}^{n-K}\left(\neq 0\right.$ since $\left.f_{K} \neq 0\right)$ shows that $\lambda_{x}^{n}$ is a solution of (5). Moreover by linearity, any linear combination of solutions of (5) is still a solution, and if (6) has $K$ distinct solutions, $\left\{f_{k}\right\}_{k=1, \ldots, K}$ is uniquely defined by a set of $K$ independent linear equations. Hence, for $\sum_{k=1}^{K} c_{k, p} W^{m t_{k}}$ "solution" of (5), $t_{k} \not \equiv t_{l} \bmod \tau$ for all $k \neq l$, there exists a unique set $\left\{f_{k}\right\}_{k=1, \ldots, K}$ such that $\left\{W^{t_{k}}\right\}_{k=1, \ldots, K}$ are the $K$ distinct roots of

$$
x^{K}-f_{1} x^{K-1}-\cdots-f_{K-1} x-f_{K}
$$

The coefficients $\widehat{y}_{p}[m]$ maybe arranged in a tall block-

```
Algorithm 1 Block-Prony_TLS
Require: An estimate on the number of effective paths \(K^{\text {est }}\),
    \(2 M+1(M \geq K)\) channel DFT coefficients \(\widehat{y}_{p}[m]=\)
    \(\sum_{k=1}^{K} c_{k, p} W_{N}^{m \bar{t}_{k}}+\widehat{q}_{p}[m]\) for \(|m| \leq M, p=1 \ldots P\).
    Build \(\boldsymbol{H}^{\left(K^{\text {est }}+1\right)}\) according to (7).
    Compute the SVD decomposition of the data matrix:
    \(\boldsymbol{H}^{\left(K^{\text {est }}+1\right)}=\boldsymbol{U} \boldsymbol{S} \boldsymbol{V}^{*}\).
    \(\phi \leftarrow \operatorname{roots}(\boldsymbol{v})\), such that \(\boldsymbol{v}\) is the right singular vector
    associated to the least singular value.
    return \(\left\{t_{k}^{\text {est }}\right\}_{k=1 \ldots K^{\text {est }}} \leftarrow-\frac{\tau}{2 \pi} \arg \phi\).
```

```
Algorithm 2 Block-ESPRIT_TLS
Require: An estimate on the number of effective paths \(K^{\text {est }}\),
    \(2 M+1(M \geq K)\) channel DFT coefficients \(\widehat{y}_{p}[m]=\)
    \(\sum_{k=1}^{K} c_{k, p} W_{N}^{m \bar{t}_{k}}+\widehat{q}_{p}[m]\) for \(|m| \leq M, p=1 \ldots P\).
    Build \(\boldsymbol{H}^{(M)}\) according to (7).
    Compute the SVD decomposition of the data matrix:
    \(\boldsymbol{H}^{(M)}=\boldsymbol{U} \boldsymbol{S} \boldsymbol{V}\).
    Extract the signal subspace basis \(\boldsymbol{\Xi}_{0}=\boldsymbol{V}_{1:(M-1), 1: K^{\text {est }}}\).
    Extract the rotated signal space basis \(\boldsymbol{\Xi}_{1}=\boldsymbol{V}_{2: M, 1: K^{\text {est }}}\).
    Solve \(\boldsymbol{\Xi}_{1}=\boldsymbol{\Xi}_{0} \boldsymbol{\Psi}\) in the TLS sense.
    return \(\left\{t_{k}^{\text {est }}\right\}_{k=1 \ldots K^{\text {est }}} \leftarrow\left\{-\frac{\tau}{2 \pi} \arg \lambda_{k}(\boldsymbol{\Psi})\right\}_{k=1 \ldots K^{\text {est }}}\).
```

Toeplitz matrix

$$
\boldsymbol{H}^{(L)}=\left[\begin{array}{c}
\boldsymbol{H}_{1}^{(L)} \\
\boldsymbol{H}_{2}^{(L)} \\
\vdots \\
\boldsymbol{H}_{P}^{(L)}
\end{array}\right]
$$

such that

$$
\boldsymbol{H}_{p}^{(L)}=\left[\begin{array}{cccc}
\widehat{y}_{p, L-M-1} & \widehat{y}_{p, L-M-2} & \cdots & \widehat{y}_{p,-M}  \tag{7}\\
\widehat{y}_{p, L-M} & \widehat{y}_{p, L-M-1} & \cdots & \widehat{y}_{p, 1-M} \\
\vdots & \vdots & \ddots & \vdots \\
\widehat{y}_{p, M} & \widehat{y}_{p, M-1} & \cdots & \widehat{y}_{p, M-L+1}
\end{array}\right],
$$

where $\widehat{y}_{i, j}=\widehat{y}_{i}[j]$. The data matrix $\boldsymbol{H}^{(L)}$ is made of $P$ Toeplitz blocks of size $(2 M+2-L) \times L$, and we assume that both block dimensions are larger or equal to $K$. It possess interesting algebraic properties which form the core of line spectra estimation techniques. We will use Lemma 1 to show three well-known spectral estimation tools which extend straightforwardly from Toeplitz data matrices to blockToeplitz ones.

The annihilating filter property for a set of exact SCS channels with $K$ distinct paths is

$$
\begin{equation*}
\boldsymbol{H}^{(K+1)} \boldsymbol{f}=\mathbf{0} \tag{8}
\end{equation*}
$$

where $\boldsymbol{f}=\left[\begin{array}{llll}1 & -f_{1} \cdots-f_{K}\end{array}\right]^{T}$ are the annihilating filter coefficients such that the polynomial $p_{\boldsymbol{f}}(z)=1-\sum_{k=1}^{K} f_{k} z^{k}$ has $K$ roots $\left\{e^{-2 \pi j t_{k} / \tau}\right\}_{k=1 \ldots K}$. The matrix $\boldsymbol{H}^{(K+1)}$ is built with blocks as in (7) (with $L=K+1$ ).

This is a direct consequence of Lemma 1. Also, the rank of the data matrix is exactly $K$,

```
Algorithm 3 Block-Cadzow denoising
Require: A block-Toeplitz matrix \(\boldsymbol{H}^{(L)}\) and a target rank \(K\).
Ensure: A block-Toeplitz matrix \(\boldsymbol{H}^{(L)}\) with rank \(\leq K\).
    repeat
        Reduce \(\boldsymbol{H}^{(L)}\) to rank \(K\) by a truncated SVD.
        Make \(\boldsymbol{H}_{p}^{(L)} p=1 \ldots P\), Toeplitz by averaging diago-
        nals.
    until convergence
```

Algorithm 4 SCS-FRI channel estimation
Require: An estimate on the number of effective paths $K^{\text {est }}$,
$2 M+1(M \geq K)$ noisy channel DFT coefficients
$\widehat{y}_{p}[m]=\sum_{k=1}^{K} c_{k, p} W_{N}^{m t_{k}}+\widehat{q}_{p}[m]$ for $|m| \leq M, p=$
$1 \ldots P$.
Ensure: Support estimate $\left\{t_{k}^{\text {est }}\right\}_{k=1 \ldots K^{\text {est }}}$
Build $\boldsymbol{H}^{(M)}$ according to (7).
$\boldsymbol{H}^{(M)} \leftarrow$ Block-Cadzow $\left(\boldsymbol{H}^{(M)}, K^{\text {est }}\right)$ [optional].
Update $\widehat{y}_{p}[m]$ with the first row and column of the
denoised block $\boldsymbol{H}_{p}^{(M)}$.
Estimate the common support with Block-Prony_TLS or
Block-ESPRIT_TLS.
Estimate $\left\{c_{k, p}\right\}$ solving $P$ linear Vandermonde systems
(4).

Proposition 1. For a set of exact SCS channels with $K$ distinct paths and in the absence of noise, $\boldsymbol{H}^{(L)}$ satisfies

$$
\operatorname{rank} \boldsymbol{H}^{(L)}=K
$$

for $K \leq L \leq 2 M+2-K$.
Proof: Let $\widetilde{\boldsymbol{H}}$ be the top-left $K \times K$ minor of $\boldsymbol{H}^{(L)}$. It can be written as the sum of $K$ rank-1 matrices:

$$
\widetilde{\boldsymbol{H}}=\sum_{k=1}^{K} c_{k,} W^{(M+1-L) t_{k}} \boldsymbol{\xi}_{k}^{*} \boldsymbol{\xi}_{k}
$$

such that $\boldsymbol{\xi}_{k}=\left[\begin{array}{llll}1 & W^{t_{k}} & W^{2 t_{k}} & \cdots\end{array} W^{(K-1) t_{k}}\right]$. If $t_{k} \not \equiv$ $t_{l} \bmod \tau$ for all $k \neq l$, then $\left\{\boldsymbol{\xi}_{1}, \ldots, \boldsymbol{\xi}_{K}\right\}$ form a set of non collinear vectors. Therefore rank $\boldsymbol{H}^{(L)} \geq K$.

Choose $K$ mutually independent rows of $\boldsymbol{H}^{(L)}$. From Lemma 1, truncating these row vectors to length $K$ preserves the linear independence. Therefore, given a row $\boldsymbol{h}=$ $[h[0], \ldots, h[L-1]]$ of $\boldsymbol{H}^{(L)}$, there exists a linear combination of these $K$ rows $\boldsymbol{h}^{\prime}$ such that

$$
h[k]=h^{\prime}[k], \quad k=0, \ldots, K-1
$$

By construction, since all rows of $\boldsymbol{H}^{(L)}$ verify the same linear recursion of degree $K, \boldsymbol{h}$ and $\boldsymbol{h}^{\prime}$ verifies this recursion. Hence $\boldsymbol{h}=\boldsymbol{h}^{\prime}$, implying rank $\boldsymbol{H}^{(L)}$ is at most $K$.
a) Block-Prony algorithm: Equation (8) is Prony's method [33], [34] for block-Toeplitz matrices. We call the corresponding algorithm "Block-Prony TLS", listed under Algorithm 1. It solves the annihilating filter equation (8) in the total least-square (TLS) sense. The crucial step is the identification of what shall be the unidimensional null space of $\boldsymbol{H}^{(K+1)}$ in a noiseless case. Solving this problem in the TLS sense yields the least right singular vector of $\boldsymbol{H}^{(K+1)}$. Prony's method is notoriously sensitive to noise,
which is to be expected as the result relies on identification of the unidimensional complement of the $K$-dimensional signal space. This sensitivity can be mitigated with prior denoising of the measurements.
b) Block-ESPRIT algorithm: Proposition 1 implies each block in the data matrix shares the same signal subspace. Hence the ESPRIT TLS algorithm outlined in [35] applies as-is to the block-Toeplitz data matrix $\boldsymbol{H}^{(L)}$. The BlockESPRIT TLS algorithm is outlined in Algorithm 2. The BlockESPRIT algorithm fulfills the same goal as the Block-Prony algorithm, but its essence is entirely different. Where Prony's method identifies a line with least energy in a $K$ dimensional space, ESPRIT finds the rotation between two $K$-dimensional subspaces in an $M$-dimensional space. What makes ESPRIT much more resilient to noise is that the two subspaces are computed from the most energetic part of the signal.
c) Block-Cadzow denoising: Proposition 1 used together with the block-Toeplitz structure property yields the "lift-andproject" denoising Algorithm 3, which we call Block-Cadzow denoising [36]. Using the same argument as in [37], the blockCadzow algorithm provably converges.
d) SCS-FRI: We have all the elements to describe the SCS-FRI algorithm. The Block-Cadzow algorithm may be used to denoise the measurements and is followed by either Block-Prony or Block-ESPRIT estimation of the common ToAs (solved in the TLS sense).

For Cadzow denoising and ESPRIT, it is empirically found that a data matrix with square blocks works well. The last step is to estimate the path amplitudes independently for each channel. This is done by solving a linear Vandermonde system (4) [38]. The processing chain at the receiver is listed in Algorithm 4 and shown in Fig. 2. Combination of Cadzow and ESPRIT for estimation of a single OFDM channel is considered in [39]. We assumed the number of paths to be known. For information theoretic criterions to estimate the number of paths, we refer to [35]. In [31], a criterion based on the effective rank is used on synthetic and measured CIR.

## C. Deterministic Multipath Channels

In [11], the authors derive the Cramér-Rao lower bound [40], [41] for estimating the positions and weights of the Diracs in FRI signals. Considering a single path with deterministic amplitude in a single-channel real-valued scenario, the minimal relative uncertainties on the location of the Dirac, $t_{1}$, and on its amplitude, $c_{1}$, are given by

$$
\begin{aligned}
& \mathbb{E}\left[\left(\frac{\Delta t_{1}}{\tau}\right)^{2}\right] \geq \frac{3(2 M+1)}{4 \pi^{2} N M(M+1)} \mathrm{PSNR}^{-1} \\
& \mathbb{E}\left[\left(\frac{\Delta c_{1}}{c_{1}}\right)^{2}\right] \geq \frac{2 M+1}{N} \mathrm{PSNR}^{-1}
\end{aligned}
$$

where $\mathrm{PSNR}=c_{1}^{2} / \sigma^{2}$ is the input peak signal to noise ratio. When there are more than two Diracs, the Cramér-Rao formula for one Dirac still holds approximately when the locations are sufficiently far apart ${ }^{1}$.

[^1]
## D. Jointly Gaussian Multipath Channels

We derive bounds on the support estimation accuracy with measurements taken according to (3). The paths coefficients $c_{k, p}$ are assumed to be jointly Gaussian, and modeled as the product of $a_{k, p}=\mathbb{E}\left[\left|c_{k, p}\right|\right]$ by a standard normal random variable $Z_{k, p}$ having the following properties, consistent with the well-known Rayleigh-fading model:

- $Z_{k, p} \sim \mathcal{N}_{\mathbb{C}}(0, \sqrt{1 / 2} \mathbb{I})$.
- Similar expected path amplitude between antennas: $a_{k} \stackrel{\text { def }}{=}$ $a_{k, 1}=a_{k, 2}=\cdots=a_{k, P}$.
- Independence between paths: $\mathbb{E}\left[Z_{k, p} Z_{k^{\prime}, p^{\prime}}^{*}\right]=0, k \neq$ $k^{\prime}$.
- The random vector $\boldsymbol{Z}_{k}=\left[Z_{k, 1} \cdots Z_{k, P}\right]^{T}$ is defined as $\boldsymbol{Z}_{k}=\boldsymbol{L}_{k} \boldsymbol{r}$, where $\boldsymbol{L}_{k}$ is the Cholesky factor of the covariance matrix $\boldsymbol{R}_{k}=\mathbb{E}\left[\boldsymbol{Z}_{k} \boldsymbol{Z}_{k}^{*}\right]$ and $\boldsymbol{r}$ is a vector of iid standard complex Gaussian random variables.

The Rayleigh-fading case can be seen as deterministic if conditioned on the path amplitudes. Thus, the Cramér-Rao bounds for random paths coefficients are random variables for which we can compute statistics. Expectation and standard deviation will respectively give the expected accuracy of the estimator and its volatility. For a single path, and a symmetric or antisymmetric $\varphi$ (not necessarily a sinc kernel), the CrámerRao bound has a concise closed form formula.

With complex-valued measurements according to (3), $K=$ 1 , and $\boldsymbol{Z}_{1}$ be a random Gaussian vector, then

$$
\begin{equation*}
\mathbb{E}\left[\left(\frac{\Delta t_{1}}{\tau}\right)^{2}\right] \geq \frac{\mathbb{E}\left[\left(\boldsymbol{Z}_{1}^{*} \boldsymbol{Z}_{1}\right)^{-1}\right]}{2 N \cdot \mathrm{dSNR}} \tag{9}
\end{equation*}
$$

where dSNR $=\left|a_{1}\right|^{2}\left\|\varphi^{\prime}\left(\boldsymbol{n} T-t_{1}\right)\right\|^{2} /\left(N \sigma^{2}\right)$ is the differential SNR and $\lambda_{1}, \cdots, \lambda_{P}$ are the eigenvalues of the covariance matrix $\boldsymbol{R}_{1}$ and $P>1$ :

- Uncorrelated paths coefficients, $\lambda_{1}=\cdots=\lambda_{P}=1$ :

$$
\mathbb{E}\left[\left(\boldsymbol{Z}_{1}^{*} \boldsymbol{Z}_{1}\right)^{-1}\right]=(P-1)^{-1}, \quad P>1
$$

- Correlated path coefficients, such that $\lambda_{1} \neq \cdots \neq \lambda_{P}$ :

$$
\mathbb{E}\left[\left(\boldsymbol{Z}_{1}^{*} \boldsymbol{Z}_{1}\right)^{-1}\right]=\sum_{p=1}^{P}\left(-\lambda_{p}\right)^{P-1} \frac{\ln \lambda_{p}}{\lambda_{p}} \prod_{p^{\prime} \neq p}\left(\lambda_{p^{\prime}}-\lambda_{p}\right)^{-1}
$$

A proof is given in [42]. The expectation of an inverse$\chi^{2}$ distributed random variable is found in various statistical handbooks. For correlated random variables, see [43]. This expression is a suitable approximation for multipaths scenario with distant paths (separated by more than twice the inverse bandwidth. It gives an important insight on the evolution of the estimation performance when uncorrelated antennas are added to the system. Namely, the RMSE decays as $1 / \sqrt{P-1}$.

In general, multiple paths are interacting with each other and the information matrix cannot be considered diagonal. Let $\boldsymbol{\Phi}$ and $\boldsymbol{\Phi}^{\prime}$ be $N \times K$ matrices such that

$$
\begin{gathered}
\Phi_{n, k}=\varphi\left((n-1) T-t_{k}\right), \quad \Phi_{n, k}^{\prime}=\varphi \prime\left((n-1) T-t_{k}\right), \\
n \in\{1, \ldots, N\}, \quad k \in\{1, \ldots, K\}
\end{gathered}
$$



Fig. 2. The SCS-FRI sampling and reconstruction scheme in a multi-antenna channel estimation setting with $P$ receiving antennas.
and

$$
\boldsymbol{C}=\operatorname{diag}\left(a_{1}, \ldots, a_{K}\right)\left(\sum_{p=1}^{P} \boldsymbol{Z}_{p}^{\prime} \boldsymbol{Z}_{p}^{\prime *}\right) \operatorname{diag}\left(a_{1}^{*}, \ldots, a_{K}^{*}\right)
$$

Yau and Bresler have shown in [12] that the Fisher information matrix $\boldsymbol{J}$ conditioned on the path amplitudes is

$$
\begin{equation*}
\boldsymbol{J}=2 \sigma^{-2} \boldsymbol{\Phi}^{\prime *} P_{\mathrm{ker} \boldsymbol{\Phi}} \boldsymbol{\Phi}^{\prime} \odot \boldsymbol{C} . \tag{10}
\end{equation*}
$$

such that $P_{\mathrm{ker} \Phi}=\mathbb{I}-\boldsymbol{\Phi} \Phi^{\dagger}$ is the projection into the nullspace of $\Phi$ and " $\odot$ " denotes the entrywise matrix product.
The Cramér-Rao bounds for the estimation of the normalized times of arrival are on the diagonal of the expectation of $\boldsymbol{J}^{-1}$. The matrix $\boldsymbol{J}$ is a complex Wishart matrix. Computing its inverse moments analytically is not an easy task, nevertheless it can be numerically computed via Monte-Carlo simulations.

## III. Application to OFDM and CDMA Downlink

1) SCS-FRI with Uniformly Scattered DFT Pilots (OFDM): The theory in Section II is developed for contiguous DFT coefficients. In OFDM communications, pilots are often uniformly laid out in frequency (ETSI DVB-T [44], 3GPP LTE [6],...). The period of pilot insertion $D$ is upper-bounded by $\Delta^{-1}$, the inverse of the delay-spread of the CIR : $D<\tau / \Delta$. If not, the CIR cannot be unambiguously recovered from the pilots because of aliasing. For a fixed number of pilots, $D$ is chosen as large as possible $(D=\lfloor\tau / \Delta\rfloor)$, as interpolation of the CIR spectrum is more robust than extrapolation.
SCS-FRI can take advantage of uniformly scattered pilot layouts [42], [45]. For $\widehat{\varphi}$ flat in $\{-M D, \ldots, M D\}$, equation (4) becomes:

$$
\begin{equation*}
\widehat{y}_{p}[m D]=\sum_{k=1}^{K} c_{k, p} W_{N}^{m D t_{k}}+\widehat{q}_{p}[m D] \tag{11}
\end{equation*}
$$

which corresponds to a dilation by $D$ of the support parameters $\left\{t_{k}\right\}$. By definition $0 \leq t_{k}<\Delta$, and so the bound on $D$ prohibits aliasing of $D t_{k}$. Therefore, SCS-FRI is applicable without other modification than division of the recovered support parameters by $D$. The results of (9) can be extended to scattered pilot with minimal effort.

The minimal uncertainties on the estimation of the parameters in the SCS-FRI scenario (11) with $P$ signals are given by

$$
\begin{aligned}
& \mathbb{E}\left[\left(\frac{\Delta t_{1}}{\tau}\right)^{2}\right] \geq \frac{3 B T}{4 D^{2} \pi^{2} M(M+1)} \mathbb{E}\left[\mathrm{ESNR}^{-1}\right] \\
& \mathbb{E}\left[\left(\frac{\Delta c_{\ell}}{c_{\ell}}\right)^{2}\right] \geq B T \mathbb{E}\left[\mathrm{PSNR}_{\ell}^{-1}\right] \quad \ell=1, \ldots, P .
\end{aligned}
$$

Such that for real-valued signal and noise ESNR = $\frac{1}{\sigma^{2}} \sum_{\ell=1}^{P} c_{\ell}^{2}$ denotes the effective signal to noise ratio and $\mathrm{PSNR}_{\ell}=c_{\ell}^{2} / \sigma^{2}$. For complex-valued signal and noise $\mathrm{ESNR}=\frac{1}{2 \sigma^{2}} \sum_{\ell=1}^{P} c_{\ell}^{*} c_{\ell}$ and $\mathrm{PSNR}_{\ell}=c_{\ell}^{*} c_{\ell} /\left(2 \sigma^{2}\right)$.

It is an explicit evaluation of dSNR in (9) - assuming $\varphi$ is the Dirichlet kernel of bandwidth $B$ as defined in (2) since

$$
\mathbb{E}\left[\left(\frac{\Delta t_{1}}{\tau}\right)^{2}\right]=\mathbb{E}\left[\left(\frac{\Delta\left(D t_{1}\right)}{\tau}\right)^{2}\right] \cdot D^{-2}
$$

2) Extension to Walsh-Hadamard Coded Schemes (CDMA): Numerous applications use the $2^{n}$-WHT to code the channel into $2^{n}$ subchannels $\left(N=2^{n}\right)$. Among others, IS-95 uses a 64 -WHT to code the downlink channel. The straightforward way to insert pilots is to use one of these subchannels as a pilot itself and use correlation based channel estimation methods as the Rake-receiver for example [46]. The SCS-FRI algorithm works in the DFT domain but can nevertheless be applied in the WHT domain with pilots uniformly scattered by $D$ a power of 2 .

Proposition 2. Let $\boldsymbol{W}_{n}$ and $\boldsymbol{S}_{n}$ be respectively the $2^{n}$-points DFT and WHT matrices obtained by Sylvester's construction:

$$
\boldsymbol{S}_{1}=\frac{1}{\sqrt{2}}\left[\begin{array}{cc}
1 & 1 \\
1 & -1
\end{array}\right], \quad \boldsymbol{S}_{i+1}=\boldsymbol{S}_{1} \otimes \boldsymbol{S}_{i}
$$

Then, for any $\ell \in\{1, \ldots, n-1\}$ the set of $\boldsymbol{S}_{n}$ 's columns with indices in $\left\{2^{\ell}+i\right\}_{1,2}$ and the set of $\boldsymbol{W}_{n}$ 's columns with indices in $\left\{(i-1 / 2) \cdot 2^{n-\ell}+1\right\}_{1 \ldots 2^{\ell}}$ span the same subspace.

Proof: We partition the Walsh-Hadamard transform ma-
trix in two "left" and "right" blocks:

$$
\boldsymbol{S}_{n}=\left[\begin{array}{ll}
\boldsymbol{S}_{n}^{(l)} & \boldsymbol{S}_{n}^{(r)}
\end{array}\right], \quad \boldsymbol{S}_{n}^{(l)}=\left[\begin{array}{c}
\boldsymbol{S}_{n-1} \\
\boldsymbol{S}_{n-1}
\end{array}\right], \quad \boldsymbol{S}_{n}^{(r)}=\left[\begin{array}{c}
\boldsymbol{S}_{n-1} \\
-\boldsymbol{S}_{n-1}
\end{array}\right]
$$

Given $\boldsymbol{w}_{2^{n}}^{k}=\left[W_{2^{n}}^{0 k} \cdots W_{2^{n}}^{\left(2^{n}-1\right) k}\right]$ the $k^{t h}$ vector of the DFT basis and $\boldsymbol{s}^{(r)} \in \operatorname{span} \boldsymbol{S}_{n}^{(r)}$ :

$$
\begin{aligned}
\left\langle\boldsymbol{w}_{2^{n}}^{k}, s_{n}^{(r)}\right\rangle & =\sum_{l=0}^{N-1} W_{2^{n}}^{k l} s_{n}^{(r)}[l] \\
& =\sum_{l=0}^{2^{n-1}-1} W_{2^{n}}^{k l} s_{n}^{(r)}[l]+W_{2^{n}}^{k\left(l+2^{n-1}\right)} s_{n}^{(r)}\left[l+2^{n-1}\right] \\
& =\sum_{l=0}^{2^{n-1}-1} W_{2^{n}}^{k l} s_{n}^{(r)}[l]\left(1-W_{2^{n}}^{k 2^{n-1}}\right)
\end{aligned}
$$

Hence

$$
\left\langle\boldsymbol{w}_{2^{n}}^{k}, \boldsymbol{s}_{n}^{(r)}\right\rangle=0, \quad \text { for } k \text { even }
$$

The spans of $\boldsymbol{S}_{n}^{(r)}$ and $\boldsymbol{S}_{n}^{(r)}$ partition the original $2^{n}$ dimensional space in two subspaces of dimension $2^{n-1}$. Let $\boldsymbol{W}_{n}^{(o)}=\left\{\boldsymbol{w}_{2^{n}}^{k}\right\}_{0 \leq 2 k+1<N, k \in \mathbb{N}}$ be the DFT basis vectors with odd indices and $\boldsymbol{W}_{n}^{(e)}=\left\{\boldsymbol{w}_{2^{n}}^{k}\right\}_{0 \leq 2 k<N, k \in \mathbb{N}}$ the ones with even indices. The spans of $\boldsymbol{W}_{n}^{(o)}$ and $\boldsymbol{W}_{n}^{(e)}$ partition the original space into two subspaces of dimension $2^{n-1}$. Since $\operatorname{span} \boldsymbol{S}_{n}^{(r)} \perp \operatorname{span} \boldsymbol{W}_{n}^{(e)}$ :

$$
\begin{aligned}
& \operatorname{span} \boldsymbol{W}_{n}^{(o)}=\operatorname{span} \boldsymbol{S}_{n}^{(r)} \\
& \operatorname{span} \boldsymbol{W}_{n}^{(e)}=\operatorname{span} \boldsymbol{S}_{n}^{(l)}
\end{aligned}
$$

This property applies recursively, since for $k \in\left\{0, \ldots, 2^{n-1}-\right.$ $1\}$ :

$$
\begin{aligned}
\left\langle\boldsymbol{w}_{2^{n}}^{2 k}, \boldsymbol{s}_{n}^{(l)}\right\rangle & =\left\langle\frac{1}{\sqrt{2}}\left[\begin{array}{l}
\boldsymbol{w}_{2^{n-1}}^{k} \\
\boldsymbol{w}_{2^{n-1}}^{k}
\end{array}\right], \frac{1}{\sqrt{2}}\left[\begin{array}{c}
\boldsymbol{s}_{n-1} \\
\boldsymbol{s}_{n-1}
\end{array}\right]\right\rangle \\
& =\left\langle\boldsymbol{w}_{2^{n-1}}^{k}, \boldsymbol{s}_{n-1}\right\rangle
\end{aligned}
$$

where $\boldsymbol{s}_{n-1} \in \operatorname{span} \boldsymbol{S}_{n-1}$.
Proposition 2 states that one can choose $2^{\ell}$ contiguous Walsh-Hadamard codewords for pilots and get $2^{\ell}$ uniformly spread DFT pilots with layout gap $D=2^{n-\ell}$. The channel coding is akin CDMA, but the pilot layout matches the one used in OFDM communication. The lesson, is that the Walsh-Hadamard transform alone achieves "scrambling" of data followed by carrier mapping in the DFT domain in a fashion similar to SC-FDMA [47]. In SC-FDMA, the data are first "scrambled" by application of a shorter length DFT.

This result has a nice interpretation in the context of generalized Fourier transforms, the $2^{n}$-WHT being itself the Fourier transform on the finite group $(\mathbb{Z} / 2 \mathbb{Z})^{n}$ instead of $\mathbb{Z} / 2^{n} \mathbb{Z}$ for the classical $2^{n}$-points DFT [42], [48]. A similar result holds for DFT on any toric finite group [42].

## IV. Application: Fading Channel Estimation in Multi-Output Systems

## A. Channel Model

1) Physical Assumptions: A linear time-invariant channel is characterized by its impulse response $h$. In mobile communications, channels are transient, but we may assume the
channel to be locally invariant around a time $\tau$. This leads to the definition of a time-dependent channel impulse response $h_{\tau}$. Consider a channel impulse response made of a large number $L$ of echoes:

$$
\begin{equation*}
h_{\tau}(t)=\sum_{l=1}^{L} \alpha_{l}(\tau) \delta\left(t-t_{l}(\tau)\right) \tag{12}
\end{equation*}
$$

The number of echoes $L$, is usually far too large to warrant a finite rate of innovation approach. However individual echoes aggregate in a smaller and manageable number of clusters $K$ [49]. The rationales behind clustering are the same as for the common support assumption: a finite bandwidth combined with background noise allow only for a limited resolution. Table I lists a few examples for which clustering applies in typical operating conditions.

This simplification is at the heart of medium and narrowband wireless communications [50]. We want to estimate $h_{\tau}$ by sending probes at the input and collecting samples at the output.

Correlation of the channel with respect to time is an important feature to exploit, however we will not consider it, as scheduling in modern communication systems makes its usage uncertain. Hence we settle on a time $\tau$ and drop it from the notation.

Communication is carried over a restricted frequency band, which is achieved by pulse-shaping with a template function $\varphi(t)$ and modulation by $e^{j \omega_{c} t}$. Applying clustering to (12) the channel impulse response becomes:

$$
\begin{align*}
h(t) & =\sum_{k=1}^{K} c_{k} \varphi\left(t-t_{k}\right)  \tag{13}\\
\text { s.t. } c_{k} & =a_{k} Z_{k}=e^{j \omega_{c} t_{k}} \sum_{\left(\alpha_{l}, t_{l}\right) \in \mathcal{C}_{k}} \alpha_{l} e^{j \omega_{c}\left(t_{l}-t_{k}\right)}
\end{align*}
$$

where $Z_{k}$ has unit-variance, $a_{k}$ is the appropriate scaling parameter and $\mathcal{C}_{k}$ is the $k^{t h}$ cluster. Assuming $\left\{\alpha_{l} e^{j \omega_{c}\left(t_{l}-t_{k}\right)}\right\}_{\left(\alpha_{l}, t_{l}\right) \in \mathcal{C}_{k}}$ contains i.i.d. elements with finite first two moments (echoes of finite energy)

$$
\lim _{\# \mathcal{C}_{k} \rightarrow \infty} Z_{k} \sim \mathcal{N}_{\mathbb{C}}(0,1)
$$

This is the classical non line of sight fading scenario where the paths amplitudes $\left|Z_{k}\right|$ are independently Rayleigh distributed.
2) Multipoint Communications, One to Many: Communication through fading channels rely on spatial diversity to gain robustness. Spatial diversity is achieved with the deployment of several antennas at the receiver and/or transmitter. We describe a spatial channel model between one transmitter and several receivers, which generalizes to MIMO communications in a straightforward manner. The physical properties of the channel are the following as shown in Fig. 3:

- The distance in between antennas $m$ and $n$ is $d_{m, n}$.
- Each path is characterized by an angle of arrival (AoA) $\theta_{k}$. To simplify computations it is assumed that the AoA is the same for all antennas (far field assumption). In the near field, a scatterer surrounds the receiver and the distribution becomes almost isotropic. Hence this assumption can be made for both regimes with limited


Fig. 3. Channel model with a single scatterer. Each scatterer is characterized by its apparent width (width/distance) $\sigma_{k} / \Delta_{k}$ and its azimuth $\theta_{k}$. This model is considered valid in the near field as well, as the scatterer surrounds the receiver, thus having no intrinsic azimuth.
error.

- The direction normal to the segment between antennas $m$ and $n$ points toward azimuth $\theta_{m, n}$.

The channel model in (13) applies to the $P$ subchannels

$$
h_{m}(t)=\sum_{k=1}^{K} a_{m, k} Z_{m, k} \varphi\left(t-t_{m, k}\right), \quad m=1, \ldots, P
$$

We assume the distance in between antennas is smaller than the achievable spatial resolution, hence

$$
t_{1, k}=t_{2, k}=\cdots=t_{P, k}
$$

To fully characterize the channel, the path correlation across antennas must be known - by assumption $Z_{m, k}$ and $Z_{m^{\prime}, k^{\prime}}$ are independent for $k \neq k^{\prime}$. Following Salz and Winters [51] we derive a formula for the autocorrelation matrix of $\boldsymbol{Z}_{k}=$ [ $Z_{1, k} \cdots Z_{P, k}$ ]. However we put a Gaussian prior on the cluster shape rather than a uniform one with discontinuities at the boundaries.
3) Spatial Correlation of Paths: As in [51], a large number of reflections are assumed to be drawn from a continuous probability distribution for each scatterer.

Proposition 3. Under the spatial channel model described in Section IV-A2, the antenna cross-correlation is closely approximated by:

$$
\begin{aligned}
& \mathbb{E}\left[Z_{k, m} Z_{k, n}^{*}\right]= J_{0}\left(\frac{\omega_{c}}{c} d_{m, n}\right) \\
&+2 \sum_{l=1}^{\infty} j^{l} \frac{I_{l}\left(\kappa_{k}\right)}{I_{0}\left(\kappa_{k}\right)} J_{l}\left(\frac{\omega_{c}}{c} d_{m, n}\right) \\
& \cdot \cos \left[l\left(-\theta_{m, n}+\theta_{k}-\frac{\pi}{2}\right)\right],
\end{aligned}
$$

where $\Delta_{k}^{2} / \sigma_{k}^{2} \approx\left(1-e^{-3 \kappa_{k} / 4}\right) \kappa_{k}$, $J_{(\cdot)}$ is the Bessel function of the first kind and $I_{(\cdot)}$ is the modified Bessel function of the first kind.

Proof: See appendix VII. It is only a close approximation since the azimuthal distribution at the receiver is approximated by a Von-Mises distribution.

TABLE III
Simulation Parameters

| Parameter | Symbol | Value |
| :--- | :---: | :--- |
| Sampling step | $T$ | 50 ns |
| Bandwidth | $B$ | 20 MHz |
| Center frequency | $f_{c}$ | 2.6 GHz |
| Frame duration (without padding) | $\tau$ | $25.55 \mu \mathrm{~s}$ |
| Samples per frame | $N_{\text {frame }}$ | 511 |
| Pilots per frame | $N$ | 63 |
| Pilot gap | $D$ | 8 |
| Delay spread | $\Delta$ | $1.6 \mu \mathrm{~s}$ |

Corollary 1. For a path width $\kappa_{k}$ large enough:

$$
\begin{aligned}
\mathbb{E}\left[Z_{k, m} Z_{k, n}^{*}\right] \approx & J_{0}\left(\frac{\omega_{c}}{c} d_{m, n}\right) \\
+2 \sqrt{2 \pi \kappa_{k}} & \sum_{l=1}^{\infty} j^{l} f_{\kappa_{k}}(l) J_{l}\left(\frac{\omega_{c}}{c} d_{m, n}\right) \\
& \cdot \cos \left[l\left(-\theta_{m, n}+\theta_{k}-\frac{\pi}{2}\right)\right]
\end{aligned}
$$

such that $f_{\kappa_{k}}$ is the centered Gaussian pdf of variance $\kappa_{k}$.
The result is in its form similar to [51], however the derivation stays closer to the original physical model.

## V. Numerical Results

## A. Simulations Setup

For simulations we use the channel model developed in Section IV, and choose its parameters to loosely follow the 3GPP-LTE standard. Its characteristics are listed in Table III. We assume 63 pilots which are uniformly spaced in frequency, one every 8 . The transmitted frame is circularly padded such as to guarantee circular convolution of the transmitted signal with the CIR.

Results are derived from three different experiments. The first two experiments deliberately use a very trivial 2-taps channel model to gain insight on the behavior of the algorithm:
A The medium has two paths separated by $2 T$. The second path's expected amplitude is $1 / 10^{t h}$ of the expected amplitude of the first path. The receiver possesses 1, 2, 4 or 8 uncorrelated antennas. The channels have exact SCS ( $\varepsilon=0$ ).
B The medium has two paths separated by $T$ or $2 T$. Both paths have the same expected amplitude. The receiver has 4 uncorrelated antennas. The channels have either exact $\operatorname{SCS}(\varepsilon=0)$ or non-exact $\operatorname{SCS}(\varepsilon=T / 50=1 \mathrm{~ns})$. The discrepancy in the ToA between antennas is uniformly distributed in $[-\varepsilon \varepsilon]$. A time lapse of $2 T / 50$ corresponds to a path length difference of 60 cm .
C This experiment is more realistic from a physical standpoint. The receiver has 5 antennas equispaced on a circle of radius 10 cm (the maximum path-difference between antennas is thus 40 cm ). The propagation medium contains 10 scatterers, which position are independently


Fig. 4. (Exp. A ) For the same global input SNR, a system with more antennas estimates the ToAs more accurately and is more resilient to noise. This is a consequence of the increased receiver diversity. The second path has $1 / 10^{\text {th }}$ the amplitude of the first path and is thus quickly buried into noise as SNR decreases. The estimation reaches the Cramér-Rao bound as long as it correctly identifies the path.


Fig. 5. (Exp. A ) Part (a) shows the performances of Block ESPRIT-TLS with or without Block Cadzow denoising. In this setup, the gain obtained with the denoising is relatively small and is achieved after one iteration. Part (b) shows the performances of Block Prony-TLS with or without Block Cadzow denoising. As expected, the performance of Prony's algorithm without denoising is very poor. After 3 denoising iterations, performances of Block Prony-TLS and Block ESPRIT-TLS are indistinguishable.
drawn from a symmetrical 2D gaussian distribution centered on the receiver. It turns out that the minimum distance between paths is 0.6 times the sampling step. A random delay corresponding to a propagation distance within $\pm 50 \mathrm{~cm}$ is added to the paths of each channel. We use the spatial correlation model derived in Proposition 3.
Results were obtained on 10000 independent noise and fading realisations, and the SER is computed after demodulation and maximum ratio combining (MRC) of the channels.

## B. Simulations Result

1) Results on Exp. A : Fig. 4 shows that the SCS-FRI algorithm efficiently estimates the ToA down to a certain SNR where the recovery breaks down. This breaking point is pushed lower as spatial diversity increases, which is to be expected. Fig. 5 compares the use and combination of the various subspace identification techniques discussed earlier. The conclusion is that the performances of Block-ESPRIT TLS or Block-Prony TLS are exactly the same on a signal denoised with the Block-Cadzow algorithm. However BlockESPRIT TLS requires fewer to none Block-Cadzow iterations than Block-Prony TLS to reach the optimum. It is well-known that Prony TLS is not robust to noise [11], [34].
2) Results on Exp. B: Fig. 6 shows that the single path CRB given in (9) is a good approximation of the true bound computed via (10) for multiple paths separated by more than twice the inverse bandwidth of the channel. This experiment also verifies the usefulness of the SCS assumption when ToAs are slightly perturbed from one antenna to another:

$$
t_{k, p}=t_{k}+E_{k, p}, \quad E_{k, p} \sim \mathrm{U}([-\varepsilon \varepsilon]), \text { i.i.d. }
$$

The error caused by the random perturbation $E_{k, p}$ is of the order of the perturbation itself, and thus we may say SCS-FRI is robust on non exact SCS channels.
3) Results on Exp. C : All estimation algorithms use the fact that the delay spread is much shorter than the frame length. The difference between lowpass interpolation and other techniques is the use of the sparsity property. Using this property alone, the SER is significantly lowered as shown in Fig. 7. The addition of the SCS property proves to be valuable.

## VI. Conclusion

We outlined the SCS-FRI algorithm, studied its performances on SCS channels estimation and computed theoretical lower-bounds for comparison. A spatial channel model was proposed for simulation purposes. The algorithm takes


Fig. 6. (Exp. B ) This figure shows that the proposed algorithms behave as expected in the presence of ToA mismatches between antennas. Part (b) motivates the separability assumption to compute the CRB of paths located more than $2 T$ apart, while Part (a) shows its inadequacy for a smaller delay $T$. The "true" estimate is obtained via Monte-Carlo simulations.


Fig. 7. (Exp. C ) Using the SCS property, the SER is decreased compared to the conventional non-parametric approach. Sparsity alone provides a significant SER improvement, which shall be combined with the common support property. At very high SNR, independent channel estimation across antennas would become preferable as the channels only approximately have the common support property.
full advantage of the main properties of outdoor multipath channels and is directly applicable to most OFDM based communication standards. Simulations indicate that SCS-FRI based on the Block-ESPRIT TLS routine seems to be the most suitable since it requires only two partial SVD with size of the model order and provides optimal accuracy.

Future work is needed for estimation of the model order (sparsity level), incorporation of temporal correlation in the model and the algorithm (tracking of the model parameters). Computational complexity is also a crucial point for mobile applications, and improvements could be made with Krylov subspaces techniques as in [52].

Also, a comparison with a discrete approach such as CS, shall be made on measured CIR, since both approach use fairly different features of the signal. A manuscript with this comparison along with sparsity level estimation and speed improvements is being prepared [31].

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## VII. Spatial Correlation Formula for Fading Channels

## A. Azimuthal Scatterers Density Distribution

The reflection density of each scatterer is normally distributed with mean $\boldsymbol{\mu}_{k}$ (its position) and covariance matrix $\sigma_{k}^{2} \mathbb{I}$ (its "girth"). The number of reflections within a scatterer is assumed to be large enough to warrant their approximation by their continuous probability density function. The azimuthal density is the integral of the scatterer's pdf over $\Gamma_{\vartheta}$ the straight path from the receiving antenna at an angle ${ }^{2} \vartheta$ :

$$
p\left(\vartheta ; \boldsymbol{\mu}_{k}, \sigma_{k}^{2}\right)=\int_{\mathbb{R}^{2}} f_{\sigma_{k}^{2}}^{(2 D)}\left(\boldsymbol{x}-\boldsymbol{\mu}_{k}\right) \mathcal{I}_{\boldsymbol{x} \in \Gamma_{\vartheta}} \boldsymbol{d} \boldsymbol{x}
$$

Reparametrization in polar coordinates yield:

$$
\begin{aligned}
p\left(\vartheta ; \boldsymbol{\mu}_{k}, \sigma_{k}^{2}\right)= & f_{\sigma_{k}^{2}}\left(\left\|\boldsymbol{\mu}_{k}\right\| \sin (\vartheta)\right) \\
& \cdot \int_{\mathbb{R}_{+}} f_{\sigma_{k}^{2}}\left(r-\left\|\boldsymbol{\mu}_{k}\right\| \cos (\vartheta)\right) J_{\boldsymbol{x}}(r, \vartheta) d r \\
= & \sigma_{k}^{-1} f\left(\sqrt{\kappa_{k}^{\prime}} \sin (\vartheta)\right) \\
& \cdot \int_{\sqrt{\kappa_{k}^{\prime}} \cos (\vartheta)}^{+\infty} \sigma_{k}^{-1} f\left(r-\left\|\boldsymbol{\mu}_{k}\right\| \cos (\vartheta)\right) \\
& \cdot\left(s+\sqrt{\kappa_{k}^{\prime}} \cos (\vartheta)\right) \sigma_{k}^{2} d s
\end{aligned}
$$

such that $\kappa_{k}^{\prime}=\left\|\boldsymbol{\mu}_{k}\right\|^{2} / \sigma_{k}^{2}$ and $J_{\boldsymbol{x}}(r, \vartheta)=r$ is the Jacobian of the cartesian to polar transformation. We performed the change of variable $s=r-\sqrt{\kappa_{k}^{\prime}} \cos (\vartheta)$. Hence, the distribution has only one degree of freedom, and after some calculus:

$$
\begin{align*}
p_{\kappa_{k}^{\prime}}(\vartheta)= & f\left(\sqrt{\kappa_{k}^{\prime}} \sin \vartheta\right)  \tag{14}\\
& \cdot\left[\sqrt{\kappa_{k}^{\prime}} \cos \vartheta \cdot F\left(\sqrt{\kappa_{k}^{\prime}} \cos \vartheta\right) f\left(\sqrt{\kappa_{k}^{\prime}} \cos \vartheta\right)\right] .
\end{align*}
$$

The circular distribution (14) is well approximated by a VonMises distribution of scale $\kappa_{k}$ :

$$
q_{\kappa_{k}}(\vartheta)=\frac{e^{\kappa_{k} \cos \vartheta}}{2 \pi I_{0}\left(\kappa_{k}\right)} .
$$

[^2]where $I_{0}$ is the $0^{t h}$ order modified Bessel function of the first kind. Asymptotically, $\kappa_{k}^{\prime} \stackrel{\kappa_{k}^{\prime} \rightarrow \infty}{=} \kappa_{k}$, and the approximation $\kappa_{k}^{\prime} \approx\left(1-e^{-3 \kappa_{k} / 4}\right) \kappa_{k}$ was found to be empirically accurate for all $\kappa_{k}$ (K-L divergence between $p_{\kappa_{k}^{\prime}}$ and $q_{\kappa_{k}}$ is less than 0.02 bits).

## B. Derivation of the Correlation Matrix Formula

Considering the setup of Fig. 3, and from [51]:

$$
R_{Z}^{(k)}[m, n]=\int_{-\pi}^{\pi} q_{\kappa_{k}}\left(\vartheta+\theta_{m, n}-\theta_{k}\right) e^{j \frac{\omega_{c}}{c} d_{m, n} \sin \vartheta} d \vartheta
$$

Then, $q_{\kappa_{k}}$ is expanded in terms of spherical harmonics via the Jacobi-Anger expansion [53](9.1):

$$
\begin{aligned}
& q_{\kappa_{k}}\left(\vartheta+\theta_{m, n}-\theta_{k}\right) \\
= & \frac{1}{2 \pi I_{0}\left(\kappa_{k}\right)}\left\{J_{0}\left(-j \kappa_{k}\right)\right. \\
& \left.+\sum_{l=1}^{\infty} j^{l} J_{l}\left(-j \kappa_{k}\right) \cos \left[l\left(\vartheta+\theta_{m, n}-\theta_{k}\right)\right]\right\} \\
= & \frac{1}{2 \pi}+\frac{1}{\pi I_{0}\left(\kappa_{k}\right)} \sum_{l=1}^{\infty} I_{l}\left(\kappa_{k}\right) \cos \left[l\left(\vartheta+\theta_{m, n}-\theta_{k}\right)\right]
\end{aligned}
$$

where the second equality is obtained with $I_{l}(j x)=j^{l} J_{l}(x)$ [53](9.6.3, 9.1.35).

We now have a series for $R_{Z}^{(k)}[m, n]$ with $l^{t h}$ term:

$$
\begin{aligned}
& \frac{I_{l}\left(\kappa_{k}\right)}{\pi I_{0}\left(\kappa_{k}\right)} \int_{-\pi}^{\pi} \cos \left[l\left(\vartheta+\theta_{m, n}-\theta_{k}\right)\right] e^{j \frac{\omega_{c}}{c} d_{m, n} \sin \vartheta} d \vartheta \\
& \stackrel{(a)}{=} \frac{I_{l}\left(\kappa_{k}\right)}{\pi I_{0}\left(\kappa_{k}\right)}\left\{\cos \left[l\left(\theta_{k}-\theta_{m, n}-\frac{\pi}{2}\right)\right]\right. \\
& \int_{-\pi}^{\pi} \cos l \vartheta e^{j \frac{\omega_{c}}{c} d_{m, n} \cos \vartheta} d \vartheta \\
& +\sin \left[l\left(\theta_{k}-\theta_{m, n}-\frac{\pi}{2}\right)\right] \\
& \left.\cdot \int_{-\pi}^{\pi} \sin l \vartheta e^{j \frac{\omega_{c}}{c} d_{m, n} \cos \vartheta} d \vartheta\right\} \\
& \stackrel{(b)}{=} \frac{2 I_{l}\left(\kappa_{k}\right)}{I_{0}\left(\kappa_{k}\right)} I_{l}\left(j \frac{\omega_{c}}{c} d_{m, n}\right) \cos \left[l\left(\theta_{k}-\theta_{m, n}-\frac{\pi}{2}\right)\right] \\
& \stackrel{(c)}{=} \frac{2 j^{l} I_{l}\left(\kappa_{k}\right)}{I_{0}\left(\kappa_{k}\right)} J_{l}\left(\frac{\omega_{c}}{c} d_{m, n}\right) \cos \left[l\left(\theta_{k}-\theta_{m, n}-\frac{\pi}{2}\right)\right]
\end{aligned}
$$

Equality (a) is obtained with some standard trigonometric identities and a shift by $-\frac{\pi}{2}$ of the variable of integration. Equality $(b)$ follows from the standard integral representation of $I_{l}$ ([53] 9.6.19). The second integrand is antisymmetric which leads the integral over the unit-circle to vanish. Finally $(c)$ is a consequence of $I_{l}(j x)=j^{l} J_{l}(x)$ again. Hence:

$$
\begin{aligned}
R_{Z}^{(k)}[m, n]= & J_{0}\left(\frac{\omega_{c}}{c} d_{m, n}\right) \\
& +\frac{2}{I_{0}\left(\kappa_{k}\right)} \sum_{l=1}^{\infty} j^{l} I_{l}\left(\kappa_{k}\right) J_{l}\left(\frac{\omega_{c}}{c} d_{m, n}\right) \\
& \cdot \cos \left[l\left(\theta_{k}-\theta_{m, n}-\frac{\pi}{2}\right)\right]
\end{aligned}
$$

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[^1]:    ${ }^{1}$ Empirically, the distance should be larger than $2 / B$.

[^2]:    ${ }^{2}$ Without loss of generality the scatterer origin is at azimuth 0 , and the antenna is located at position $\mathbf{0}$

