

Non-Arguesian configurations in a modular lattice

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To the memory of András Huhn

1. Introduction. In [1] we showed that if L is non-Arguesian, then there exist, in the ideal lattice of L , elements p_α , $\alpha \in \mathbf{5}^{[2]}$, and q_β , $\beta \in \mathbf{5}^{[3]}$, that are related to each other in a manner similar to the ten points and ten lines in a non-Arguesian configuration in a projective plane. In the lattice case, however, each p_α is a point in a plane P_α , and each q_β is a line in the plane Q_β , with all of these planes being intervals in the ideal lattice of L . Actually our construction yielded thirty two intervals $I_\mu = u_\mu/z_\mu$, $\mu \subseteq \mathbf{5}$, and it was shown that, with at most two exceptions, these intervals are non-degenerate projective planes. The exceptional intervals, I_\circ and I_5 , were shown to be projective geometries of dimension three or less.

Our present objective is to describe in greater detail how the various intervals I_μ fit together. The notation and terminology of [1] will be in effect. A non-Arguesian perspectivity configuration (or PC), \mathbf{d} , will be called *prime* if \mathbf{d} covers \mathbf{d}_* in $\mathbf{PC}(L)$. These PC's and their associated intervals $I_\mu = u_\mu/z_\mu$, $\mu \subseteq \mathbf{5}$, will be the primary objects of our investigation. To simplify the notation, we write I_i for $I_{\{i\}}$, I_{ij} for $I_{\{ij\}}$, $I_{\gamma i}$ for $I_{\mathbf{5} \setminus \{i\}}$, etc.

It is easy to see that if, $\emptyset \neq \mu < \nu \neq \mathbf{5}$ ($<$ means “is covered by”), then the planes I_μ and I_ν are either transposes of each other (possibly equal) or else they are connected by a two dimensional gluing (either loose or tight). Much less is known about the intervals I_\circ and I_5 . In the examples that have been constructed so far, these too are non-degenerate projective planes, but we do not know if this is always the case. We do however show that, if I_\circ is either 2 or 3 dimensional, then it is non-degenerate. By duality, the same holds for I_5 .

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There are two further technical conditions that apply to PC's. A PC, \mathbf{d} , is called *stable* if, whenever two intervals of the form I_i and I_{ij} are transposes of each other, they are equal. This supposed restriction causes no real loss of generality since we will show that, for every non-stable \mathbf{d} , there exists a stable prime \mathbf{e} with $\mathbf{e} < \mathbf{d}$. A PC, \mathbf{d} , is called *Boolean* if the two functions $\mu \rightarrow z_\mu$ and $\mu \rightarrow u_\mu$ of 2^5 into L are both lattice homomorphisms. Clearly, if \mathbf{d} is Boolean, then $L' := \cup \{I_\mu : \mu \subseteq 5\}$ is a sublattice of L of finite length. A fundamental result states that, if \mathbf{d} is both Boolean and stable, $\{I_\mu : \mu \subseteq 5\}$ consists of 2^r planes where $0 \leq r \leq 3$. In this case the length of L' is at most 9, and each simple subdirect factor of L' has length 6 or less.

Much less is known about the case when \mathbf{d} is stable but not Boolean. We do show however that in this case the twenty planes, I_μ , $2 \leq |\mu| \leq 3$, are distinct from each other and from the planes of the form I_i or $I_{\neg i}$. Hopefully this case will be broken down eventually into subcases for which reasonable descriptions can be found.

Some examples of the above cases can be found in [3].

2. The gluings. Throughout this section we work with a fixed prime PC, \mathbf{d} , in a modular lattice, L .

Lemma 2.1. *For distinct $i, j \in 5$, $z_i z_j = z_\emptyset$.*

Proof. By definition, z_\emptyset is the meet of all the entries in the matrix \mathbf{d} . Since each diagonal entry is the meet of all entries in its row (or column), it follows that z_\emptyset is the meet of the diagonal entries in \mathbf{d} . For distinct $i, j, k \in 5$, we have

$$z_i z_j = (d_{*ij} d_{*ik})(d_{*ij} d_{*jk}) \leq d_{*ik} d_{*jk} = z_k.$$

Consequently $z_i z_j = z_\emptyset$.

Lemma 2.2. *For all $\mu, \nu \subseteq 5$,*

- (1) $z_\mu + z_\nu = z_{\mu \cup \nu}$, if $\mu \cap \nu \neq \emptyset$;
- (2) $z_\mu z_\nu = z_{\mu \cap \nu}$, if $\mu \cup \nu \neq 5$;
- (3) $u_\mu + u_\nu = u_{\mu \cup \nu}$, if $\mu \cap \nu \neq \emptyset$;
- (4) $u_\mu u_\nu = u_{\mu \cap \nu}$, if $\mu \cup \nu \neq 5$.

Proof. Statement (1) and its dual (4) are, respectively, parts (2) and (1) of [1; Lemma 5.2]. It therefore suffices to prove (2). Moreover we may assume that $|\mu| \leq |\nu|$, $\mu \cap \nu \subset \nu$, and $|\mu \cup \nu| = 4$. We consider four cases:

(A) $|\mu| = 2$; $|\nu| = 3$; $|\mu \cap \nu| = 1$. We may assume $\mu = \{i, j\}$ and $\nu = \{i, k, m\}$. Then

$$z_\mu z_\nu = d_{*ij}(d_{*ik} + d_{*im}) = d_{*ii} = z_i = z_{\mu \cap \nu}.$$

(B) $|\mu| = 3$; $|\nu| = 3$; $|\mu \cap \nu| = 2$. We may assume $\mu = \{i, j, k\}$ and $\nu = \{i, j, m\}$.

Then

$$z_\mu z_\nu = (d_{*ij} + d_{*jk})(d_{*ij} + d_{*jm}) = d_{*ij} = z_{ij} = z_\mu \cap \nu.$$

(C) $|\mu|=2; |\nu|=2; |\mu \cap \nu|=0$. We may assume $\mu = \{i, j\}$ and $\nu = \{k, m\}$. Then

$$z_\mu z_\nu = d_{*ij} d_{*km} \cong d_{*ij}(d_{*ik} + d_{*im}) d_{*km} (d_{*ik} + d_{*jk}) = d_{*ii} d_{*kk} = z_i z_k = z_\emptyset,$$

by 2.1.

(D) $|\mu|=1; |\nu|=3; |\mu \cap \nu|=0$. We may assume $\mu = \{i\}$ and $\nu = \{j, k, m\}$. Then

$$z_\mu z_\nu = z_i z_{ij} z_{jkm} = z_i z_j = z_\emptyset,$$

by (A) and 2.1.

Lemma 2.3. For $i \in 5$, the four elements, $d_{ij}u_i$ with $j \neq i$, are four points in general position in the plane I_i .

Proof. Let i, j, k, m, n be the distinct members of 5. Then, by computing with intervals,

$$d_{ij}u_i/z_i = d_{ij}u_{ikm}/d_{ij}u_{ikm}(d_{ik} + d_{im}) \cong (d_{ij}u_{ikm} + d_{ik} + d_{im})/(d_{ik} + d_{im}) =$$

(by transposition)

$$= (d_{ij} + d_{ik} + d_{im})u_{ikm}/(d_{ik} + d_{im}) = u_{ikm}/(d_{ik} + d_{im}).$$

Now $d_{ik} + d_{im}$ is a line in I_{ikm} , and is therefore covered by u_{ikm} . Thus $z_i \prec d_{ij}u_i$ for each $j \neq i$. To see that the four points are in general position, we compute

$$(d_{ij}u_i + d_{ik}u_i)d_{im}u_i \cong (d_{ij} + d_{ik})d_{im} = d_{ii} = z_i.$$

Theorem 2.4. If μ and ν are non-empty proper subsets of 5 with $\mu \prec \nu$, then either

$$z_\mu \prec u_\mu z_\nu \quad \text{and} \quad (u_\mu + z_\nu) \prec u_\nu,$$

or

$$z_\mu = u_\mu z_\nu \quad \text{and} \quad (u_\mu + z_\nu) = u_\nu.$$

Proof. The intervals I_μ and I_ν are of the same length and have comparable upper and lower endpoints. Consequently, $z_\mu \prec u_\mu z_\nu$ if and only if $(u_\mu + z_\nu) \prec u_\nu$, and $z_\mu = u_\mu z_\nu$ holds just in case $(u_\mu + z_\nu) = u_\nu$ is true. Therefore we need only show that for each $\mu \prec \nu$, at least one of the four conditions holds. By duality, we need only consider $|\mu|=1$ or 2.

Assume that $\mu = \{i\}$ and $\nu = \{i, j\}$. By 2.3, $z_i \prec u_i d_{ij}$ whence $u_i z_{ij}$ must equal one of those two elements. Thus $z_\mu \prec u_\mu z_\nu$ or $z_\mu = u_\mu z_\nu$.

Assume now that $\mu = \{i, j\}$ and $\nu = \{i, j, k\}$. By the Main Theorem of [1], the element $q = d_{ij} + d_{ik}$ is a line in the plane I_ν , and qu_μ is a line on the point d_{ij} in I_μ . Now $z_\mu \cong u_\mu z_\nu \cong qu_\mu$. This last inequality must be strict since

$$d_{ij} z_{ijk} = d_{ij}(z_{ij} + z_{ik}) = z_{ij} + d_{ij} z_{ik} = z_{ij} + z_i = z_{ij} \prec d_{ij}.$$

Therefore the length of $u_\mu z_\nu / z_\mu$ is at most 2 and one of our relations must again hold.

Lemma 2.5. *For all $i \in 5$, $u_\circ z_i$ either covers or equals z_\circ .*

Proof. For distinct i, j, k, m in 5,

$$u_\circ z_i = z_i u_i u_{jkm} = z_i u_{jkm}, \quad \text{and} \quad z_\circ = z_i z_j = z_i d_{ij}(d_{jk} + d_{jm}) = z_i(d_{jk} + d_{jm}).$$

Since $(d_{jk} + d_{jm}) < u_{jkm}$, the conclusion follows.

Lemma 2.6. *Any four of the five elements, z_i , $i \in 5$, are independent over z_\circ .*

Proof. If $i, j, k, m \in 5$ are distinct, then

$$z_i(z_j + z_k + z_m) \cong z_i z_{jkm} = z_\circ.$$

Theorem 2.7. *The following conditions are equivalent:*

- (1) *The five elements, $z_i u_\circ$, $i \in 5$, are points in general position in I_\circ ;*
- (2) *I_\circ is a non-degenerate 3-space;*
- (3) *length $(I_\circ) = 4$;*
- (4) *$z_\circ < z_i u_\circ$, for all $i \in 5$.*

Proof. Now [1; Theorem 5.4] gives us that $\text{length}(I_\circ) \cong 4$. Thus (1) \Rightarrow (2) and (2) \Rightarrow (3) are trivial. If $z_\circ = z_i u_\circ$, for some $i \in 5$, then $I_\circ \cong (z_i + u_\circ) / z_i$, a subinterval of a length 3 lattice. Therefore (3) \Rightarrow (4). Finally if (4) holds, then the $z_i u_\circ$ are five points in I_\circ by 2.5. By 2.6, any four of these points are independent. From $\text{length}(I_\circ) \cong 4$, we deduce (1).

Corollary 2.8. *If the conditions of the theorem hold, then for each $i \in 5$, I_i transposes down onto the interval $u_\circ / z_i u_\circ$ and $u_\circ = \sum (z_s u_\circ : s \neq i)$.*

Theorem 2.9. *If $\text{length}(I_\circ) = 3$, then at least two of the intervals I_i transpose down onto I_\circ . Thus in this case as well, I_\circ is a non-degenerate projective space (i.e. a plane).*

Proof. Since any four of the elements $z_i u_\circ$ are independent, at least one out of each four must be z_\circ . Therefore at least two of the five such elements must be z_\circ . But this forces, for these i , I_i to transpose down onto I_\circ since both intervals are the same length.

Theorem 2.10. *The duals of 2.7, 2.8, and 2.9 also hold. In particular, if I_5 is of length 3 or 4, it is a non-degenerate projective space.*

3. Boolean configurations. The definition of a Boolean configuration in Section 1 contains redundancies. We already know, for instance, that whenever $\mu \cap \nu \neq \emptyset$, $z_\mu + z_\nu = z_{\mu \cap \nu}$ holds for any PC. In this section we will reduce the number of con-

ditions needed to be checked in order to show that a PC, \mathbf{d} , is Boolean. As before, we assume that \mathbf{d} is a prime PC in a modular lattice, L .

Recall that a subset $U \subseteq L$ is called *distributive* if it generates a distributive sublattice of L . A 3-element subset $U = \{a, b, c\}$ is distributive if either $(a+b)c = ac+bc$ or $(a+c)(b+c) = ab+c$.

Lemma 3.1. *The following conditions on a prime PC are equivalent:*

- (1) $z_\mu + z_\nu = z_{\mu \cup \nu}$ for all $\mu, \nu \subseteq 5$;
- (2) $z_\mu + z_\nu = z_{\mu \cup \nu}$ for some $\mu, \nu \neq \emptyset$ with $\mu \cap \nu = \emptyset$, and $\mu \cup \nu \neq 5$;
- (3) $\{z_{ij}, z_{ik}, z_{jk}\}$ is distributive for all pairwise distinct $i, j, k \in 5$;
- (4) $\{z_{ij}, z_{ik}, z_{jk}\}$ is distributive for some pairwise distinct $i, j, k \in 5$;
- (5) $\{d_{23}, d_{24}, d_{34}\}$ is distributive.

Proof. By 2.2, (1) is equivalent to $z_\mu + z_\nu = z_{\mu \cup \nu}$ with the added condition that μ and ν are disjoint. By noting that $z_2 = d_{23}d_{24}$, $z_3 = d_{23}d_{34}$, and $z_{23} = d_{23}(d_{24} + d_{34})$, (5) is equivalent to $z_\mu + z_\nu = z_{\mu \cup \nu}$ with $\mu = \{2\}$ and $\nu = \{3\}$. By using the special automorphisms of $\text{PC}(L)$, we get that (5) is equivalent to $z_\mu + z_\nu = z_{\mu \cup \nu}$ with the added constraint that μ and ν are disjoint singletons. This last property and 2.2 however easily imply that $z_\mu = \sum \{z_i : i \in \mu\}$ for all $\mu \subseteq 5$ and this implies (1). Therefore (1) is equivalent to (5).

A priori, (1) implies (2). Conversely, assume that (2) holds with μ or ν a non-singleton. If $\mu = \{i\}$ and $\nu \supseteq \{j, k\}$, then

$$z_{ij} = z_{ij}(z_\mu + z_\nu) = z_i + z_{ij}z_\nu = z_i + z_j.$$

If $\mu = \{i, j\}$ and $\nu = \{k, m\}$, then

$$z_{ijk} = z_{ijk}(z_\mu + z_\nu) = z_{ij} + z_{ijk}z_\nu = z_{ij} + z_k.$$

Thus this case reduces to the previous one. Therefore (1) \Leftrightarrow (2) \Leftrightarrow (5).

Now for distinct $i, j, k \in 5$, $\{z_{ij}, z_{ik}, z_{jk}\}$ is distributive if and only if

$$z_{ij}(z_{ik} + z_{jk}) = z_{ij}z_{ik} + z_{ij}z_{jk}.$$

The left side of this equation is z_{ij} , and the right side is $z_i + z_j$. Thus (4) implies (2) and (1) implies (3). This completes the proof.

Lemma 3.2. *For a PC, \mathbf{d} , the following are equivalent:*

- (1) $z_\mu z_\nu = z_{\mu \cap \nu}$ for all $\mu, \nu \subseteq 5$;
- (2) $z_\mu z_\nu = z_{\mu \cap \nu}$ for some $\mu, \nu \neq 5$ with $\mu \cup \nu = 5$ and $\mu \cap \nu \neq \emptyset$;
- (3) $\{z_{ijk}, z_{ijm}, z_{ijn}\}$ is distributive for all distinct $i, j, k, m \in 5$;
- (4) $\{z_{ijk}, z_{ijm}, z_{ijn}\}$ is distributive for some distinct $i, j, k, m \in 5$.

Proof. In considering (2), we may assume that $|\mu| \leq |\nu|$. The possible values for $s = |\mu|$ and $t = |\nu|$ are therefore

$$(s, t) = (4, 4), (3, 4), (3, 3), \text{ and } (2, 4).$$

For each of these four ordered pairs, (s, t) , we define:

$$(\forall_{st}) \quad z_\mu z_\nu = z_{\mu \cap \nu} \quad \text{for all } \mu, \nu \text{ with } \mu \cup \nu = 5, \quad |\mu| = s, \quad \text{and } |\nu| = t;$$

$$(\exists_{st}) \quad z_\mu z_\nu = z_{\mu \cap \nu} \quad \text{for some } \mu, \nu \text{ with } \mu \cup \nu = 5, \quad |\mu| = s, \quad \text{and } |\nu| = t.$$

We claim that (3), (4), and each of the eight statements above are equivalent to each other.

Assume that i, j, k, m, n are all distinct in 5, and consider the equation

$$(*) \quad z_{ijkm} z_{ijkn} = z_{ijk}.$$

This can be rewritten as

$$(z_{ijk} + z_{ijm})(z_{ijk} + z_{ijn}) = z_{ijk}$$

and since $z_{ijm} z_{ijn} = z_{ij} \leq z_{ijk}$, this is equivalent to

$$(**) \quad \{z_{ijk}, z_{ijm}, z_{ijn}\} \text{ is distributive.}$$

Since $(*)$ is $\{i, j, k\}$ -symmetric and $(**)$ is $\{k, m, n\}$ -symmetric, it follows that both conditions are invariant under all symmetries of the indices. Therefore (3), (4), (\forall_{44}) , and (\exists_{44}) are equivalent.

If (\forall_{44}) holds, then

$$z_{ijk} z_{ijmn} = z_{ijkm} z_{ijkn} z_{ijmn} = z_{ijm} z_{ijn} = z_{ij},$$

and thus (\forall_{34}) holds. On the other hand if (\exists_{34}) holds, say $z_{ijk} z_{ijmn} = z_{ij}$, then

$$z_{ijkm} z_{ijmn} = z_{im} + z_{ijk} z_{ijmn} = z_{im} + z_{ij} = z_{ijm}$$

and (\exists_{44}) holds. Consequently, (\forall_{44}) is equivalent to both (\exists_{34}) and (\forall_{34}) .

If (\forall_{34}) holds, then

$$z_{ijk} z_{imn} = z_{ijk} z_{ikmn} z_{ijmn} = z_{ij} z_{ik} = z_{ij},$$

and thus (\forall_{33}) holds. On the other hand if (\exists_{33}) holds, say $z_{ijk} z_{imn} = z_i$, then

$$z_{ijk} z_{ijmn} = z_{ij} + z_{ijk} z_{imn} = z_{ij} + z_i = z_{ij}$$

and (\exists_{33}) holds. Consequently, (\forall_{44}) is equivalent to both (\exists_{33}) and (\forall_{33}) .

A similar argument shows that each of the statements (\exists_{24}) and (\forall_{24}) is equivalent to (\forall_{44}) . Therefore (2), (3), and (4) are equivalent.

To obtain (2) implies (1) we need only consider complementary subsets of 5. Assuming (2), we obtain

$$z_i z_{jkmn} = z_{ij} z_{ik} z_{jkmn} = z_j z_k = z_\emptyset, \quad \text{and} \quad z_{ij} z_{kmn} = z_{ijk} z_{ijm} z_{kmn} = z_k z_m = z_\emptyset.$$

Thus (2) implies (1) and the proof is complete.

Corollary 3.3. *If, for some non-empty proper subsets, $\mu \subset \nu \subseteq 5$, $z_\mu = z_\nu$, then \mathfrak{d} is Boolean.*

Proof. The inclusion $\mu \subset \nu$ implies that $u_\mu \cong u_\nu$, and since I_μ and I_ν are both projective planes we get equality here as well. Now let $\kappa = \nu \setminus \mu$ and $\lambda = \mathfrak{S} \setminus \kappa = (\mathfrak{S} \setminus \nu) \cup \mu$. We compute

$$z_{\mu \cup \kappa} = z_\nu = z_\nu + z_\kappa = z_\mu + z_\kappa, \quad \text{and} \quad z_{\nu \cap \lambda} = z_\mu = z_\mu z_\lambda = z_\nu z_\lambda.$$

By Lemmas 3.1, 3.2 and their duals, it follows that \mathfrak{d} is Boolean.

The above argument works in general to produce:

Theorem 3.4. *A PC, \mathfrak{d} , is Boolean if and only if, for some distinct $i, j \in \mathfrak{S}$,*

$$z_i + z_j = z_{ij}, \quad \text{and} \quad z_{\neg i} + z_{\neg j} = z_{\neg ij}, \quad \text{and} \quad u_i + u_j = u_{ij}, \quad \text{and} \quad u_{\neg i} + u_{\neg j} = u_{\neg ij}.$$

4. Stable configurations. We still assume that \mathfrak{d} is a prime PC in a modular lattice, L .

Theorem 4.1. *Let \mathfrak{d} be prime and stable. Then \mathfrak{d} is either Boolean or satisfies*

(***). *For all $\mu, \nu \subseteq \mathfrak{S}$, if $\emptyset \subset \mu \prec \nu \subset \mathfrak{S}$, then $z_\mu \prec u_\mu z_\nu$ and $u_\mu + z_\nu \prec u_\nu$.*

Proof. Let \mathfrak{d} be stable, and take $\emptyset \subset \mu \prec \nu \subset \mathfrak{S}$. By 2.4 we must have I_μ transposing up to I_ν , or $z_\mu \prec z_\nu u_\mu$ and $u_\mu + z_\nu \prec u_\nu$. If the first property holds, then let $\{i\} = \nu \setminus \mu$ and take $j \in \mu$. Now

$$u_j z_{ij} = u_j z_{ij} u_\mu z_\nu = u_j z_\mu z_{ij} = z_i.$$

Since \mathfrak{d} is stable, this implies $I_{ij} = I_i$, and hence \mathfrak{d} is Boolean by 3.3.

Lemma 4.2. *Let \mathfrak{d} be a prime PC. For any $x \in d_{02}/z_0$, there exists a unique PC, \mathfrak{e} , such that*

$$e_{01} = d_{01}(x + d_{12}), \quad e_{02} = x, \quad e_{12} = d_{12}(x + d_{01}),$$

and for $\{i, j\} = \{0, 1\}$ and $k \in \{3, 4\}$,

$$e_{ik} = d_{ik}(d_{jk} + e_{01}).$$

Moreover if x is not less than or equal to z_{02} , then \mathfrak{e} is non-Arguesian.

Proof. The uniqueness of \mathfrak{e} is obvious for, by [1; Theorem 3.2], every PC in L is completely determined by the elements listed above. Thus we are left with showing the existence. This however also follows from [1; Lemma 2.4] and the quoted theorem. If \mathfrak{e} were Arguesian, then $\mathfrak{e} = \mathfrak{e}_*$, and

$$x = e_{02} = e_{*02} \cong d_{*02} = z_{02}.$$

Lemma 4.3. *If \mathfrak{d} is not stable, then there exists a prime PC, $\mathfrak{e} \prec \mathfrak{d}$ that is both stable and Boolean.*

Proof. If \mathbf{d} is a prime PC that is not stable then there exists $i, j \in 5$ such that I_i transposes up to I_{ij} but is not equal to I_{ij} . Thus for these i and j we have

$$z_{ij}u_j = z_j, \quad z_{ij} + u_i = u_j, \quad \text{and} \quad z_i < z_j.$$

By using the special automorphisms of $\mathbf{PC}(L)$, we may assume that $i=0$, and $j=2$. Using $x=d_{02}u_0$ in the previous lemma, we obtain a non-Arguesian PC, \mathbf{e} , with $\mathbf{e} < \mathbf{d}$ and $e_{02}=d_{02}u_0$. To see that \mathbf{e} is prime, we note that $z_0=z_0(\mathbf{d}) \cong z_{02}(\mathbf{e}) < e_{02}$ (and that $z_0 < e_{02}$). Therefore $z_0=z_0(\mathbf{e})=z_{02}(\mathbf{e}) < e_{02}$.

That \mathbf{e} is Boolean follows from 3.3 and the fact that $z_0(\mathbf{e})=z_{02}(\mathbf{e})$, but \mathbf{e} may not be stable. What this \mathbf{e} has done is replace the transpose $I_0(\mathbf{d})$ up to $I_{02}(\mathbf{d})$ with the equality $I_0(\mathbf{e})=I_{02}(\mathbf{e})$. But for all $i \in 5$, direct calculations show that

$$z_i(\mathbf{d}) = z_i(\mathbf{e}) \cong z_{ij}(\mathbf{e}) \cong z_{ij}(\mathbf{d}).$$

Therefore this \mathbf{e} preserves all equalities of the form $I_i(\mathbf{d})=I_{ij}(\mathbf{d})$. This means that after finitely many steps (at most 5^2) all transpositions are replaced by equalities and the resultant PC is both Boolean and stable.

Thus if L is a non-Arguesian modular lattice, we can find, in the lattice of ideals of L , a prime (non-Arguesian) PC, \mathbf{d} . If \mathbf{d} is stable, then \mathbf{d} is either Boolean or satisfies $(***)$. If \mathbf{d} is not stable, we can find a smaller PC, \mathbf{e} , that is both stable and Boolean. Therefore every non-Arguesian variety of modular lattices contains a non-Arguesian lattice with a stable (non-Arguesian) PC. The Boolean case has a nice finite solution which we present in the next section. By [3], there exists infinitely many distinct stable PC's satisfying $(***)$, and these authors at least have found no classification of them. Our only general result is the following.

Theorem 4.4. *Let \mathbf{d} be a stable non-Boolean PC. Then the twenty planes, I_μ , $2 \cong |\mu| \cong 3$, are distinct from each other, and from the planes, I_i and $I_{\neg i}$, $i \in 5$.*

Proof. Let $\mu \neq \nu \subseteq 5$ satisfy:

$$1 \cong |\mu|, |\nu| \cong 4, \quad \min \{|\mu|, |\nu|\} \cong 3, \quad \text{and} \quad \max \{|\mu|, |\nu|\} \cong 2.$$

We wish to show that the assumption, $z_\mu=z_\nu$, leads to a contradiction. We obtain this contradiction by producing a covering pair of subsets, $\kappa < \lambda$, with $z_\kappa=z_\lambda$, and invoking $(***)$.

If $\mu \cap \nu \neq \emptyset$, then $z_\mu+z_\nu=z_{\mu \cup \nu}$, and we may choose κ to be the set of smallest cardinality and λ to be any cover contained in $\mu \cup \nu$. This produces our contradiction on $(***)$. Therefore we may conclude that

$$[0] \quad \mu \cap \nu = \emptyset.$$

Therefore there exists $i \in \mu \setminus \nu$. But now we have for all $i \in \mu$

$$z_{\nu \cup \{i\}} = z_\nu \cup \{i\} + z_\nu = z_\nu \cup \{i\} + z_\mu = z_{\mu \cup \nu}.$$

To avoid conflict with (***), we must have:

$$i \in \mu \text{ implies } [1] \mu \cup v = 5 \text{ and } 4 \cong v \cup \{i\}, \text{ or}$$

$$[2] \mu \cup v = v \cup \{i\}.$$

We also have $j \in v \setminus \mu$ and the trick above can be applied again to produce

$$j \in v \text{ implies } [3] \mu \cup v = 5 \text{ and } 4 \cong \mu \cup \{j\}, \text{ or}$$

$$[4] \mu \cup v = \mu \cup \{j\}.$$

Now [0] makes [2] equivalent to $\mu = \{i\}$, and [4] equivalent to $v = \{j\}$. Thus [1] and [4] are incompatible as well as [2] and [3]. Our initial assumptions deny the conjunction of [2] and [4], so we must have [1] and [3]. But this forces $3 \cong |v|$ and $|\mu|$ which contradicts [0]. This concludes the proof.

5. Stable Boolean configurations. Throughout this section, \mathfrak{d} will be a prime, Boolean, and stable PC in a modular lattice, L . The lattice homomorphisms,

$$z, u: 2^5 \rightarrow L,$$

produce Boolean congruences on 2^5 which are, of course, determined by their respective ideals, $\text{Id}(z)$ and $\text{Id}(u)$, of subsets congruent to \emptyset . Now $\{i\} \in \text{Id}(z) \Leftrightarrow z_i = z_o \Leftrightarrow$ for all $j \neq i, z_{ij} = z_j \Leftrightarrow$ for all $j \neq i, u_{ij} = u_j \Leftrightarrow u_i = u_o \Leftrightarrow \{i\} \in \text{Id}(u)$. Therefore $\text{Id}(z) = \text{Id}(u)$, and by factoring out this ideal we produce, for some r with $0 \leq r \leq 5$, lattice embeddings

$$z', u': 2^r \rightarrow L.$$

Our first result shows that this r can be further restricted.

Lemma 5.1. *If \mathfrak{d} is Boolean and stable, then the set $\{I_\mu: \mu \subseteq 5\}$ consists of 2^r planes for some $r, 0 \leq r \leq 3$.*

Proof. Let i, j, k, m, n be distinct members of 5 , and assume that for all $s \neq n, z_{sn} > z_n$. From 2.3 and stability, this implies that for all $s \neq n, u_n z_{sn} = u_n d_{sn}$. 2.3 also says that $\{u_n d_{sn}: s \neq n\}$ are points in general position in I_n . But \mathfrak{d} is Boolean, and therefore

$$u_i d_{in} \cong u_i z_{in} (u_j z_{jn} + u_k z_{kn} + u_m z_{mn}) \cong z_{in} z_{jkmn} = z_n.$$

This is a contradiction.

Thus for every $n \in 5$, there exists an $s \neq n$ such that $z_{sn} = z_n$. Again since \mathfrak{d} is Boolean this implies that for every $n \in 5$, there exists an $s \neq n$ such that $z_s = z_o$. Elementary counting now produces two distinct $s \in 5$ with $z_s = z_o$.

We may therefore replace 5 by r for $0 \leq r \leq 3$, and assume that we have lattice monomorphisms,

$$z, u: 2^r \rightarrow L,$$

that satisfy:

(1) $I_\mu := u_\mu/z_\mu$ is a non-degenerate projective plane for all $\mu \subseteq r$;

(2) For all $\emptyset \subseteq \mu < \nu \subseteq r$, $z_\mu u_\nu < z_\nu$ and $u_\mu + z_\nu < u_\nu$.

We define $L' := \cup \{I_\mu : \mu \subseteq r\}$. Clearly L' is a sublattice of L of finite length.

Lemma 5.2. L' is simple if and only if for all $i \in r$, $z_i \cong u_{\neg i}$.

Proof. If our condition fails, then, for the offending $i \in r$, L' is the disjoint union of the filter, $\uparrow z_i$ and the ideal, $\downarrow u_{\neg i}$. Thus L' is not simple.

Conversely, assume the condition holds. We proceed by induction on r . If $r=0$, then L' is a non-degenerate projective plane and hence simple. If $0 < r \leq 3$, take a prime quotient q/p in L' , and let θ be the congruence it generates. Since $L' = \uparrow z_i \cup \downarrow u_{\neg i}$ and $z_i \cong u_{\neg i}$, we must have this quotient in $\uparrow z_i$ or in $\downarrow u_{\neg i}$. By induction, θ collapses either the filter or the ideal. Since $z_i < u_{\neg i}$, induction applies also to the other part and θ collapses all of L' . Therefore L' is simple.

Theorem 5.3. Suppose \mathbf{V} is a variety of modular lattices and assume that there exists a Boolean, prime PC in some member of \mathbf{V} . Then there exists in \mathbf{V} a simple non-Arguesian lattice of length $3+r$, with $0 \leq r \leq 3$, and a Boolean, stable, prime PC, \mathbf{d} , in L with the following properties:

(1) L is generated by $\{d_{ij} : i \neq j \text{ in } 5\}$;

(2) The set $\{I_\mu : \mu \subseteq 5\}$ consists of precisely 2^r planes.

Proof. By 4.3. there exists in some member L of \mathbf{V} a PC, \mathbf{d} , that is prime, Boolean, and stable. By 5.1, the set $\{I_\mu : \mu \subseteq 5\}$ consists of 2^r distinct planes for some r , with $0 \leq r \leq 3$. We may assume without loss of generality that L is generated by the PC and is therefore the union of the planes I_μ .

Since L is obviously of finite length, we may assume that its length is as small as possible. We claim that in this case L is simple. To see this, we consider a homomorphism $\varphi: L \rightarrow S$, where S is simple and φ does not identify d_{01} and d_{*01} . Clearly $\varphi(\mathbf{d})$ is a (non-Arguesian) PC in S and in fact also prime and Boolean (since $\varphi(z_\mu(\mathbf{d})) = z_\mu(\varphi(\mathbf{d}))$ and similarly for the u 's). The length of S therefore cannot be less than that of L . This makes φ an isomorphism and L simple.

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