## UTV Expansion Pack

## Special-Purpose

# Rank Revealing Algorithms 

Version 1.0 for Matlab 6.5

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#### Abstract

This collection of Matlab software supplements and complements the package UTV Tools from 1999, and includes implementations of special-purpose rank-revealing algorithms developed since the publication of the original package. We provide algorithms for computing and modifying symmetric rank-revealing VSV decompositions, we expand the algorithms for the ULLV decomposition of a matrix pair to handle interference-type problems with a rankdeficient covariance matrix, and we provide a robust and reliable Lanczos algorithm which - despite its simplicity - is able to capture all the dominant singular values of a sparse or structured matrix. These new algorithms have applications in signal processing, optimization and LSI information retrieval. The corresponding manuscript is:


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## 1. Introduction

The Matlab package UTV Tools [5] from 1999 provides a collection of algorithms for computing and modifying (i.e., up- and downdating) rank-revealing decompositions of general matrices. These decompositions have many applications in signal processing, where they are used as fast and reliable alternatives [9], [20] to the versatile but computationally expensive and hard-to-update singular value decomposition (SVD).

Since the publication of UTV Tools more work has been done in the area of rank-revealing decompositions and algorithms. This work is motivated by the interest in using specialized rank-revealing algorithms, designed to take advantage of the underlying structure of the problem in consideration. The present package provides Matlab implementations of some of these newly developed algorithms, with emphasis on algorithms that expand the application areas of the original package (hence the name of the new package). Similar to the first package, the routines in this package can be considered as templates for more specialized implementations, perhaps in other computer languages, that can exploit the computer hardware.

Symmetric rank-revealing VSV decompositions for semidefinite and indefinite matrices were developed in [13] and [16] to provide algorithms and decompositions that take into account the symmetry of the matrix. Compared to the general UTV decompositions, the VSV decompositions lead to savings in computer time as well as advantages in the approximation properties of reduced-rank matrix approximations derived from the symmetric decompositions. The rank and subspace information provided by the VSV decompositions have applications, e.g., in deflation methods for solving block-structured symmetric indefinite systems [7] arising in optimization algorithms and PDE solvers.

The rank-revealing ULLV decomposition was originally developed for revealing the rank of a matrix quotient, defined as the product of one matrix and the pseudoinverse of another matrix, and with applications in noise reduction problems with broadband noise where the noise covariance matrix has full rank. When the noise covariance matrix is rank deficient (which is the case for interference or narrow-band noise) then the correct matrix quotient involves a weighted pseudoinverse [10], and the corresponding ULLV decomposition must reflect this. The most convenient way to deal with the full-rank and the rank deficient cases is to provide two different ULLV algorithms for the two variations of the decomposition.

While rank-revealing decompositions are convenient tools for dense matrices, they may be less suited for large sparse or structured matrices. For this reason we also provide a Lanczos algorithm for computing the dominant singular triplets of a matrix. Our algorithm demonstrates that if such an algorithm is based on complete reorthogonalization and explicit restarts, then the code need not be very complicated. The core of our implementation requires less than 100 lines of Matlab code, has a simple structure, and is thus suited for implementation on dedicated hardware platforms (in contrast to many other sophisticated and much more general - implementations in mathematical software libraries).

In addition to the above algorithms, and for completeness, we provide implementations of a few simple and "heuristic" algorithms which will often reveal the numerical rank, but
without the safety (and slight overhead) of a genuine rank-revealing algorithm.
Finally we provide a few scripts that demonstrate the use of our functions in connection with rank-deficient KKT systems in optimization, noise and interference reduction in signal processing, and signal extraction in NMR signals.

In the following sections we summarize the algorithms, giving new theory where it is needed. We conclude with a few numerical examples and an overview of the new package. Throughout the paper, the norm $\|\cdot\|$ denotes the 2 -norm, while $I$ and $E$ denote the identity matrix and the exchange matrix (consisting of the columns of the identity matrix in reverse order). Moreover, $L$ and $R$ always denote lower and upper triangular matrices, $V$ is always an orthogonal matrix, and $\Omega$ is always a signature matrix. We also make use of Matlab's colon notation to indicate submatrices.

## 2. The Symmetric VSV Decomposition

The rank-revealing VSV decomposition of a symmetric matrix was first discussed by Luk and Qiao [16] (for Toeplitz matrices). A careful study of various algorithms based on initial triangular factorization can be found in [13], while a study of the accuracy of approximations based on the VSV decomposition is given in [4]

### 2.1. Definitions

Assume that the symmetric matrix $A \in \mathbb{R}^{n \times n}$ has numerical rank $k$, i.e., there is a welldefined gap between the $k$ th singular value $\sigma_{k}$ and the next. Then the rank-revealing VSV decomposition of $A$ takes the form

$$
A=V S V^{T}, \quad S=\left(\begin{array}{ll}
S_{11} & S_{12}  \tag{2.1}\\
S_{12}^{T} & S_{22}
\end{array}\right), \quad V=\left(V_{1}, V_{2}\right)
$$

where the symmetric matrix $S$ is partitioned such that it reveals the numerical rank of $A$, i.e., the singular values of the $k \times k$ leading submatrix $S_{11}$ approximate the first $k$ singular values of $A$, while the norms $\left\|S_{12}\right\|$ and $\left\|S_{22}\right\|$ of the off-diagonal and trailing blocks are both of the order $\sigma_{k+1}$, cf. [13], [16].

The matrix $V$ is orthogonal, and it is partitioned such that the column spaces of the two blocks $V_{1}$ and $V_{2}$ are approximations to the subspaces spanned by the first $k$ and the last $n-k$ right singular vectors of $A$, respectively. In the signal processing literature, these two subspaces are referred to as the signal and noise subspaces. See [4] concerning the accuracy of these approximations.

For practical purposes, we choose to compute and represent the matrix $S$ in the factored form

$$
\begin{equation*}
S=T^{T} \Omega T \tag{2.2}
\end{equation*}
$$

in which $T$ is an upper or lower triangular matrix, and $\Omega$ is a signature matrix, i.e., a diagonal matrix with $\pm 1$ on the diagonal, such that the inertia of $A$ is preserved in the inertia of $\Omega$. If $A$ is positive definite then $\Omega$ is the identity matrix.

For semidefinite matrices, it was found in [13] that the optimal form of $T$ is lower triangular, because this choice leads to more accurate approximations of the signal and noise subspaces. Our package therefore includes software for computing the VSV decomposition $A=V L^{T} L V^{T}$ of a symmetric semidefinite matrix. We provide two functions hvsvsd and Ivsvsd, optimized for the high-rank case ( $k \approx n$ ) and low-rank case ( $k \ll n$ ), respectively.

For indefinite matrices, the singular vector estimation (which is part of the VSV algorithm) is simpler when $T$ is upper triangular, while a lower triangular $T$ provides a decomposition that is consistent with the semidefinite case. We provide high-rank VSV algorithms for both forms: the functions hvsvid_L and hvsvid_R compute the lower triangular form $A=V L^{T} \Omega L V^{T}$ and the upper triangular form $A=V R^{T} \Omega L R^{T}$, respectively. In the

## Low-Rank VSV Algorithm Ivsvid:

1. Compute $A=V L^{T} \Omega L V^{T}$ and let $k \leftarrow 1$.
2. Condition estimation: let $\widetilde{\sigma}_{k}$ estimate $\left\|L(k: n, k: n)^{T} \Omega(k: n, k: n) L(k: n, k: n)\right\|$ and let $w_{k}$ estimate the corresponding right singular vector.
If $\widetilde{\sigma}_{k}<\tau$ then $k \leftarrow k-1$, exit.
Revealment: determine orthogonal $P_{k}$ such that $P_{k}^{T} w_{k}=(1,0, \ldots, 0)^{T}$; update $L(k: n, k: n) \leftarrow L(k: n, k: n) P_{k}$ and $V(:, k: n) \leftarrow V(:, k: n) P_{k} ;$ update $L(k: n, k: n) \leftarrow Q_{k}^{T} L(k: n, k: n), \Omega(k: n, k: n) \leftarrow Q_{k}^{T} \Omega(k: n, k: n) Q_{k}$, where the hypernormal $Q_{k}$ ensures that the updated $L$ is triangular.
3. Deflation: let $k \leftarrow k+1$.
4. Go to step 2.

High-Rank VSV Algorithms hvsvid_T with $\mathrm{T}=\mathrm{L}$ or R :
Compute $A=V T^{T} \Omega T V^{T}$ and let $k \leftarrow 1$.
2. Condition estimation: let $\widetilde{\sigma}_{k}$ estimate $\sigma_{\min }\left(T(1: k, 1: k)^{T} \Omega(1: k, 1: k) T(1: k, 1: k)\right)$
and let $w_{k}$ estimate the corresponding right singular vector.
If $\widetilde{\sigma}_{k}>\tau$ then exit.
Revealment: determine orthogonal $P_{k}$ such that $P_{k}^{T} w_{k}=(0, \ldots, 0,1)^{T}$;
update $T(1: k, 1: k) \leftarrow T(1: k, 1: k) P_{k}$ and $V(:, 1: k) \leftarrow V(:, 1: k) P_{k} ;$
update $T(1: k, 1: k) \leftarrow Q_{k}^{T} T(1: k, 1: k)$ and $\Omega(1: k, 1: k) \leftarrow Q_{k}^{T} \Omega(1: k, 1: k) Q_{k}$,
where the hypernormal $Q_{k}$ ensures that the updated $T$ is triangular.
7. Deflation: let $k \leftarrow k-1$.
8. Go to step 2.

Figure 2.1: The VSV algorithms for symmetric indefinite matrices.
low-rank case the dilemma vanishes, and the function lvsvid computes the lower triangular form.

An alternative, but more expensive, approach to computing a high-rank indefinite VSV decomposition with a lower triangular $T$ is to first compute the upper triangular form $A=$ $V R^{T} \Omega R V^{T}$ and then compute the QR factorization $R^{T}=Q L^{T}$ which yields the lower triangular form $A=(V Q) L^{T} \Omega L(V Q)^{T}$. This approach is easy to implement using Matlab's qr function, but it is more expensive and therefore we do not provide an implementation.

### 2.2. Algorithms

The generic algorithm for computing the VSV decomposition of a symmetric semidefinite matrix is quite simple, because the singular values of $T$ are the square roots of the singular values of $A$ when $\Omega=I$. First we compute the symmetrically pivoted Cholesky factorization $\Pi A \Pi^{T}=\bar{C}^{T} \bar{C}$ (we use rook pivoting as implemented in [14]), followed by the computation of the rank-revealing ULV decomposition $E \bar{C} E=\widehat{U} L \widehat{V}^{T}$ (using high-rank and low-rank functions from UTV Tools). As a result, we obtain the desired decomposition $A=(\Pi \widehat{V}) L^{T} L(\Pi \widehat{V})^{T}$.

The generic algorithm for indefinite matrices starts with a symmetrically pivoted LDL ${ }^{T}$ factorization $\Pi A \Pi^{T}=\bar{L} \bar{D} \bar{L}^{T}$, using the rook pivoting implemented in [14]. Next, the middle block diagonal matrix $\bar{D}$ is replaced by the signature matrix, $\bar{L} \bar{D} \bar{L}^{T}=\widehat{W} \widehat{L} \widehat{\Omega} \widehat{L}^{T} \widehat{W}^{T}$, where $\widehat{W}$ is an orthogonal block diagonal matrix with $1 \times 1$ and $2 \times 2$ blocks on the diagonal; see $\S 4.2$ in [13].

Finally we reveal the rank of the product $\widehat{L} \widehat{\Omega} \widehat{L}^{T}$, by "peeling off" the small or large singular values one at a time. Our algorithms take basis in the following two reformulations and partitionings with $R=\widehat{L}^{T}, L=E \widehat{L}^{T} E$ and $\Omega=E \widehat{\Omega} E$ :

$$
\begin{align*}
R^{T} \widehat{\Omega} R & =\left(\begin{array}{cc}
R_{11}^{T} & 0 \\
R_{12}^{T} & R_{22}^{T}
\end{array}\right)\left(\begin{array}{cc}
\widehat{\Omega}_{1} & 0 \\
0 & \widehat{\Omega}_{2}
\end{array}\right)\left(\begin{array}{cc}
R_{11} & R_{12} \\
0 & R_{22}
\end{array}\right) \\
& =\left(\begin{array}{cc}
R_{11}^{T} \widehat{\Omega}_{1} R_{11} \\
R_{12}^{T} R_{1} R_{11}^{T} & R_{12}^{T} \widehat{\Omega}_{1} R_{12}+R_{12}^{T} \widehat{\Omega}_{2} R_{22}
\end{array}\right)  \tag{2.3}\\
L^{T} \Omega L & =\left(\begin{array}{cc}
L_{11}^{T} & L_{21}^{T} \\
0 & L_{22}^{T}
\end{array}\right)\left(\begin{array}{cc}
\Omega_{1} & 0 \\
0 & \Omega_{2}
\end{array}\right)\left(\begin{array}{cc}
L_{11} & 0 \\
L_{21} & L_{22}
\end{array}\right) \\
& =\left(\begin{array}{cc}
L_{11}^{T} \Omega_{1} L_{11}+L_{21}^{T} \Omega_{2} L_{21} & L_{21}^{T} \Omega_{2} L_{22} \\
L_{22}^{T} \Omega_{2} L_{21} & L_{22}^{T} \Omega_{2} L_{22}
\end{array}\right) . \tag{2.4}
\end{align*}
$$

The indefinite VSV algorithms are summarized in Fig. 2.1. Following the ideas from [3] in the low-rank case, we can now determined orthogonal transformations such that they, when applied symmetrically to $L^{T} \Omega L$, ensure that the singular values of the ( 1,1 )-block in (2.4) approximate the largest singular values of $A$. The construction of these transformations involves the computation of the largest singular value and corresponding right singular vector of the submatrix $L_{22}^{T} \Omega_{2} L_{22}$, and the user can choose between power iterations and the Lanczos method. At the same time, hypernormal rotations are used to maintain the triangular form of $L$.

In the high-rank case we follow the ideas from [16] and construct orthogonal transformations which ensure that the smallest singular values of $A$ are revealed in the (2,2)-block in (2.3) and (2.4). This involves the computation of the smallest singular value and corresponding right singular vector of the (1,1)-block. In the upper triangular case (2.3) this is done by means of inverse iterations applied to the submatrix $R_{11}^{T} \widehat{\Omega}_{1} R_{12}$.

In the lower triangular case (2.4), however, it is impractical to apply inverse iterations to the submatrix $L_{11}^{T} \Omega_{1} L_{11}+L_{21}^{T} \Omega_{2} L_{21}$ because we do not have a useful factorization of this matrix. The inverse iterations are much easier to use when we can ignore the second term, but unfortunately this is not always the case: according to Thm. 4.3 in [13] \| $L_{21}^{T} \Omega_{2} L_{22} \|$ and $\left\|L_{22}^{T} \Omega_{2} L_{22}\right\|$ are guaranteed to be small, but this does not imply that $\left\|L_{21}^{T} \Omega_{2} L_{21}\right\|$ is small. Our solution is to apply a single step of block QR refinement to $L$, as described below; this ensures that the norm $\left\|L_{21}\right\|$ of the off-diagonal block in $L$ is always small.

### 2.3. Hypernormal Rotations and Their Break-Down

Hypernormal transformations are introduced in [2], and their use in our VSV algorithms is discussed in [13]. The "building blocks" of hypernormal transformations are Givens and hyperbolic rotations, the latter performing the transformation (for $|\alpha|>|\beta|)$ :

$$
\left(\begin{array}{rr}
c & -s \\
-s & c
\end{array}\right)\binom{\alpha}{\beta}=\binom{\left(\alpha^{2}-\beta^{2}\right)^{1 / 2}}{0}
$$

where $c^{2}-s^{2}=1$. The hyperbolic transformation is not defined when $|\alpha|=|\beta|$, and it is has large elements $|\alpha|$ and $|\beta|$ when $|\alpha|$ is close to $|\beta|$.

In our algorithms, the hyperbolic transformations are used to annihilate fill in the triangular matrix during the revealment steps (see Fig. 2.1). Consider the following situation,
where a right Givens rotation has introduced a nonzero element " $*$ " in position $(i+1, i)$ :

$$
\begin{array}{rlcccccc} 
& & \times & \times & \times & \times & \times & \times \\
& & & \times & \times & \times & \times & \times \\
i & \rightarrow & & & \times & \times & \times & \times \\
i+1 & \rightarrow & & & * & \times & \times & \times \\
& & & & & & \times & \times \\
& & & & & & & \times
\end{array}
$$

If the fill satisfies $\left|r_{i+1, i}\right| \approx\left|r_{i i}\right|$ then we introduce large elements in the updated $R$ which cancel in the product $R^{T} \Omega R$; an undesirable situation in numerical computations. Our remedy is to detect this situation and resort to a "fix." When $\left\|r_{i+1, i}|-| r_{i i}\right\|<10^{-5} \| R(i$ : $i+$ $1, i: i+1) \|$, we perform a cyclic permutation of columns $i$ through $i+2$, leading to the form

$$
\begin{array}{rccccccc} 
& & \times & \times & \times & \times & \times & \times \\
& & & \times & \times & \times & \times & \times \\
i & \rightarrow & & & \times & \times & \times & \times \\
i+1 & \rightarrow & & & \times & \times & * & \times \\
& & & & & \times & & \times \\
& & & & & & & \times
\end{array}
$$

after which we use hypernormal transformations to annihilate the two elements below the diagonal in columns $i$ and $i+1$. We then return to the condition estimation step and restart the revealment process. If we only permuted columns $i$ and $i+1$ then the difficulty would arise again in the restarted revealment step.

### 2.4. Block QR Refinement

In analogy with block QR refinement of UTV decompositions, we can apply a similar algorithm to the VSV decompositions in order to reduce the norm of the off-diagonal blocks. We discuss the algorithm for the upper triangular version only; the algorithm for the lower triangular version is practically the same.

Given $R$ partitioned as in (2.3) we first apply a sequence of right orthogonal transformations to annihilate the submatrix $R_{12}$, thus filling out the elements in the $(2,1)$-block. These elements, in turn, are annihilated by means of left hypernormal transformations which create new elements in the ( 1,2 )-block.

We now justify this approach when applied to $S=R^{T} \Omega R$. Let $R=L_{B} Q_{B}$ where $L_{B}$ is lower block triangular and $Q_{B}$ is orthogonal; then $S=\left(R^{T} \Omega L_{B}\right) Q_{B}$ and a block QR step consists of formally forming the product $S_{B} \equiv Q_{B}\left(R^{T} \Omega L_{B}\right)=L_{B}^{T} \Omega L_{B}$. Next, let $L_{B}=H_{B} R_{B}$ where $R_{B}$ is upper triangular and $H_{B}$ is hypernormal with $H_{B}^{T} \Omega H_{B}=\Omega_{B}$ (in which $\Omega_{B}$ is a new signature matrix). Inserting this we obtain $S_{B}=R_{B}^{T} \Omega_{B} R_{B}$, showing that the new factors $R_{B}$ and $\Omega_{B}$ indeed correspond to performing a block QR step on $S$.

The block QR refinement is implemented in the function vsv_qrit which determines whether it is applied to a semidefinite or an indefinite matrix and, in the latter case, whether it is applied to the $L$ or $R$ version.

### 2.5. Rank-One Modifications

We also provide algorithms for rank-one modifications of the form

$$
\begin{equation*}
A^{\prime}=A+\omega v v^{T} \tag{2.5}
\end{equation*}
$$

where $\omega= \pm 1$ and $v$ a vector. Equation (2.5) can be recast as

$$
A^{\prime}=V\binom{T}{v^{T} V}^{T}\left(\begin{array}{cc}
\Omega & 0 \\
0 & \omega
\end{array}\right)\binom{T}{v^{T} V} V^{T}
$$

When $A$ is semidefinite and $\omega=1$, the updating is equivalent to a rank-one update of the ULV decomposition where the numerical rank cannot decrease. The updating is implemented in function vsvsd_up and uses functions from UTV Tools.

When $\omega=-1$ or when $A$ is indefinite then the numerical rank of $A^{\prime}$ can stay the same, or it can increase or decrease by one. Then the updated factors are computed by applying left hypernormal rotations to annihilate the row $v^{T} V$. This modification algorithm is implemented in the two functions vsvid_L_mod and vsvid_R_mod.

For efficiency reasons one should avoid to apply the rank-revealing post processing to the full $S$ matrix. We partition $v^{T} V=d^{T}=\left(d_{1}^{T}, d_{2}^{T}\right)$ according to (2.3) and apply right Givens rotations $G$ such that

$$
d_{2}^{T} G^{T}=e_{2}^{T}=\left(\left\|d_{2}\right\|, 0, \ldots, 0\right)^{T}
$$

At the same time we apply left hypernormal rotations $H$ to maintain the triangular form of the (2,2)-block. Introducing $R_{12}^{\prime}=R_{12} G, R_{22}^{\prime}=H^{T} R_{22} G$ and $\Omega_{2}^{\prime}=H^{T} \Omega_{2} H$, we now have

$$
\begin{aligned}
& V^{T} A^{\prime} V=\left(\begin{array}{cc}
I & 0 \\
0 & G
\end{array}\right)\left(\begin{array}{cc}
R_{11} & R_{12}^{\prime} \\
0 & R_{22}^{\prime} \\
d_{1}^{T} & e_{2}^{T}
\end{array}\right)^{T}\left(\begin{array}{ccc}
\Omega_{1} & 0 & 0 \\
0 & \Omega_{2}^{\prime} & 0 \\
0 & 0 & \omega
\end{array}\right)\left(\begin{array}{cc}
R_{11} & R_{12}^{\prime} \\
0 & R_{22}^{\prime} \\
d_{1}^{T} & e_{2}^{T}
\end{array}\right)\left(\begin{array}{cc}
I & 0 \\
0 & G
\end{array}\right)^{T} \\
& =\left(\begin{array}{cc}
R_{11}^{T} \Omega_{1} R_{11} & R_{11}^{T} \Omega_{1} R_{12} G+\omega d_{1} e_{2}^{T} \\
G^{T} R_{12}^{T} \Omega_{1} R_{11}+\omega e_{2} d_{1}^{T} & G^{T}\left(R_{12}^{T} \Omega_{1} R_{12}+R_{22}^{T} \Omega_{2} R_{22}\right) G+\omega e_{2} e_{2}^{T}
\end{array}\right) .
\end{aligned}
$$

Since $S=R^{T} \Omega R$ reveals the rank of $A$, we know that both norms $\left\|R_{11}^{T} \Omega_{1} R_{12} G\right\|=\left\|R_{11}^{T} \Omega_{1} R_{12}\right\|$ and $\left\|G^{T}\left(R_{12}^{T} \Omega_{1} R_{12}+R_{22}^{T} \Omega_{2} R_{22}\right) G\right\|=\left\|R_{12}^{T} \Omega_{1} R_{12}+R_{22}^{T} \Omega_{2} R_{22}\right\|$ are small. Hence, due to the structure of $d_{1} e_{2}^{T}$ and $e_{2} e_{2}^{T}$, it is not possible to have any elements of large magnitude in the last $n-k-1$ rows or columns of the above matrix. Therefore, once $d^{T}$ has been annihilated, it suffices to reveal the rank of the leading $(k+1) \times(k+1)$ submatrix of the updated $S$ factor.

## 3. The Gap-Revealing QLP Factorization

Stewart [21] introduced the so-called QLP factorization

$$
\begin{equation*}
A=Q L P^{T} \tag{3.1}
\end{equation*}
$$

in which $Q$ and $P$ are orthogonal, and $L$ is lower triangular. The factorization is gap revealing in the sense that the absolute values of the diagonal elements of $L$ often track the singular values of $A$; but there is no guarantee that this is always the case. Hence, the factorization is not rank-revealing in the strict sense used in this package.

To compute the QLP decomposition, we compute a pivoted QR factorization $A \Pi_{P}=Q R$ followed by a second pivoted QR factorization $R^{T} \Pi_{Q}=P L^{T}$, and thus $A=\left(Q \Pi_{Q}\right) L\left(\Pi_{P} P\right)^{T}$. For high-rank matrices, this is easy to implement with Matlab's QR factorization, and it is implemented in function hqlp.

For low-rank matrices, Huckaby and Chan [15] implemented an algorithm using interleaved left and right Householder transformations. The algorithm essentially produces one row of $L$ at a time, starting from the top, and stops as soon as a gap is revealed. A this stage, the heuristic is that if we compute the full QLP factorization then the norms of the $(2,1)$ - and (2,2)-blocks of $L$ will be small. Hence we can neglect these blocks and return the low-rank approximation

$$
\begin{equation*}
A_{k}=Q(:, 1: k) L(1: k, 1: k) P(:, 1: k)^{T}, \tag{3.2}
\end{equation*}
$$

where the gap appears between singular values $\sigma_{k}$ and $\sigma_{k+1}$. This algorithm is implemented in function Iqlp, and we emphasize that is computes the rank- $k$ matrix approximation in (3.2), not a full factorization.

## 4. The ULLIV Decomposition

The ULLV decomposition of a matrix pair $(A, B)$ with $A \in \mathbb{R}^{m \times n}, B \in \mathbb{R}^{p \times n}$ was originally defined for the case $m \geq n \geq \operatorname{rank}(A)$ and $p \geq n=\operatorname{rank}(B)$ in [17]. Algorithms for computing and modifying this decomposition are included in UTV Tools.

In certain applications, such as interference reduction [10], [12], the matrix $B$ does not have full column rank. This led Luk and Qiao [18] to define an alternative decomposition, which we shall refer to as the ULLIV decomposition. Assume again that $m \geq n \geq \operatorname{rank}(A)$ while $B$ has full row rank, i.e., $\operatorname{rank}(B)=p<n$. Then the ULLIV decomposition takes the form

$$
A=U_{A} L_{A}\left(\begin{array}{cc}
L & 0  \tag{4.1}\\
0 & I
\end{array}\right) V^{T}, \quad B=U_{B}(L, 0) V^{T}
$$

in which $I$ is an identity matrix of order $n-p, U_{A} \in \mathbb{R}^{m \times n}$ has orthonormal columns, and $U_{B} \in \mathbb{R}^{p \times p}$ and $V \in \mathbb{R}^{n \times n}$ are orthogonal. Moreover, $L_{A} \in \mathbb{R}^{n \times n}$ and $L \in \mathbb{R}^{p \times p}$ are both lower triangular.

As shown in [10], when $\operatorname{rank}(B)<n$ it is the matrix quotient $A B_{A}^{\dagger}$ that is required, and whose numerical rank should be revealed. Here $B_{A}^{\dagger}$ is the $A$-weighted pseudoinverse of $B$. Given the ULLIV decomposition (4.1) it is proved in [10] that

$$
A B_{A}^{\dagger}=U_{A} L_{A}(1: p, 1: p) U_{B}^{T},
$$

showing that the leading $p \times p$ block of $L_{A}$ must be rank revealing. To compute such a ULLIV decomposition, we start with the QR factorization $B=(L, 0) V^{T}$ followed by the QR factorization $A V\left(\begin{array}{cc}L^{-1} & 0 \\ 0 & I\end{array}\right)=U_{A} L_{A}$. Setting $U_{B}=I$ we thus have an initial decomposition, which is then made rank-revealing by applying the similar steps from the ULLV algorithm to $L$ and $U_{B}$ as well as to the first $p$ columns of $L_{A}, U_{A}$ and $V$. This algorithm is implemented in function ulliv.

An efficient algorithm for updating the ULLIV decomposition when a row $a^{T}$ is appended to $A$ is described in [18]. The algorithm takes its basis in the formulation

$$
\binom{A}{a^{T}}=\widetilde{U}_{A}\binom{L_{A}}{d^{T}}\left(\begin{array}{cc}
L & 0 \\
0 & I
\end{array}\right) V^{T}, \quad B=U_{B}(L, 0) V^{T}
$$

with

$$
\widetilde{U}_{A}=\left(\begin{array}{cc}
U_{A} & 0 \\
0 & 1
\end{array}\right), \quad d^{T}=a^{T} V\left(\begin{array}{cc}
L^{-1} & 0 \\
0 & I
\end{array}\right) .
$$

In the first stage, the partial row $d(p+1: n)^{T}$ is annihilated by means of left Givens rotations which are absorbed in $\widetilde{U}_{A}$; these rotations maintain the small and rank-revealing elements in $L_{A}(1: p, 1: p)$. In the second stage, the remaining elements of the row $d^{T}$ are annihilated by interleaved right and left Givens rotations - this stage is identical to the ULLV updating algorithm from UTV tools, and modifies $U_{B}$ and $L$ as well as the first $p$ columns of $\widetilde{U}_{A}, L_{A}$
and $V$. Small elements are maintained in rows $k+2$ to $p$ of $L_{A}$, where $k$ is the numerical rank. In the third stage, which is also identical to that of ULLV updating, the numerical rank of the updated $L_{A}(1: p, 1: p)$ is revealed. The complete algorithm is implemented in function ulliv_up_A.

When a row $b^{T}$ is appended to $B$ then we must distinguish whether the rank increases or stays the same, because of our assumption that $B$ has full row rank. An updating algorithm for the former case is described in [18]; we found it convenient to augment the algorithm with an additional stage, which simplifies the rank-revealing step. This algorithm takes its basis in the formulation

$$
A=U_{A}\left(L_{A}, 0\right)\left(\begin{array}{cc}
L & 0 \\
0 & I \\
f_{1}^{T} & f_{2}^{T}
\end{array}\right) V^{T}, \quad\binom{B}{b^{T}}=\widetilde{U}_{B} \widetilde{I}\left(\begin{array}{cc}
L & 0 \\
0 & I \\
f_{1}^{T} & f_{2}^{T}
\end{array}\right) V^{T}
$$

with

$$
\left(f_{1}^{T}, f_{2}^{T}\right)=b^{T} V, \quad \widetilde{U}_{B}=\left(\begin{array}{cc}
U_{B} & 0 \\
0 & 1
\end{array}\right), \quad \widetilde{I}=\left(\begin{array}{ccc}
I_{p} & 0 & 0 \\
0 & 0 & 1
\end{array}\right) .
$$

In the first stage, we use right Givens rotations to reduce the partial row $f_{2}^{T}$ to the form $(\phi, 0, \ldots, 0)$ with $\phi=\left\|f_{2}\right\|$. These rotations are also applied to the columns of $V$, and due to the presence of the matrix $I$ the same rotations are propagated from the right to the last $n-p$ columns of $L_{A}$. The resulting fill is annihilated again by means of left Givens rotations which are absorbed in $U_{A}$.

We now interchange rows $p+1$ and $n+1$ of the third factor, as well as the same columns of the second factors. This results in factors with the structure

$$
\begin{aligned}
& \binom{B}{b^{T}}=\widetilde{U}_{B}\left(\begin{array}{cc}
I_{p+1} & 0 \\
0 & 0
\end{array}\right)\left(\begin{array}{cccccccccc}
\times & & & & & & & & & \\
\times & \times & & & & & & & & \\
\times & \times & \times & & & & & & & \\
\times & \times & \times & \times & & & & & & \\
\times & \times & \times & \times & \times & & & & & \\
\times & \times & \times & \times & \times & \times & & & & \\
\times & \times & \times & \times & \times & \times & \times & & & \\
& & & & & & & 1 & 1 & \\
& & & & & & & & & \\
& & & & & & 1 & & &
\end{array}\right) .
\end{aligned}
$$

Note the zero column and the spike in the second factor of $A$. The $\varepsilon$ symbols represent rows $k+1, \ldots, p$ with small elements that reveal the numerical rank $k$. In order to maintain as many small elements as possible in $L_{A}$ we augment the Luk-Qiao algorithm by first performing a cyclic downshift of rows $k+1$ through $p+1$, and then we annihilate the resulting horizontal
spike in row $k+1$ by means of right Givens rotations:

$$
\begin{aligned}
& \rightarrow\left(\begin{array}{cccccccc}
\times & & & & & & & \\
\times & \times & & & & & & \\
\times & \times & \times & & & & & \\
\times & \times & \times & \times & & & & \\
\varepsilon & \varepsilon & \varepsilon & \varepsilon & \varepsilon & & & \\
\varepsilon & \varepsilon & \varepsilon & \varepsilon & \varepsilon & \varepsilon & & \\
\varepsilon & \varepsilon & \varepsilon & \varepsilon & \varepsilon & \varepsilon & & \\
\times & \times & \times & \times & \times & \times & & \\
\times & \times & \times & \times & \times & \times & \times & \times \\
\times & \times & \times & \times & \times & \times & \times & \times \\
\times & \times & \times
\end{array}\right)
\end{aligned}
$$

When these rotations are propagated into $L$ from the left, they create fill which is annihilated again by means of right Givens rotations. The triangular structure of $L_{A}$ and $L$ is thus maintained, and $L_{A}$ has small elements in rows $k+2$ through $p+1$.

The remaining steps are identical to the algorithm from [18]. The single " 1 " in the bottom row of the third factor is chased to the left by interleaved swaps of neighbor columns and Givens rotations applied to the rows to annihilate the fill. These transformations are propagated to the left and right, and they create fill in the second factor of $A$ which, in turn, is annihilated with left Givens rotations. A final Givens rotation applied to rows 1 and $n+1$ annihilates the " 1 " in position ( $n+1,1$ ), and a simple scaling restores the identity matrix $I_{p+1}$.

Having thus obtained a zero bottom row in the third factor we neglect the rightmost column of the first two factors. The result is a new factorization in which $L$ and $I$ have dimensions $(p+1) \times(p+1)$ and $(n-p-1) \times(n-p-1)$, respectively. Finally we reveal the rank of the updated $L_{A}(1: p+1,1: p+1)$ by at most two rank-revealing steps. The complete algorithm is implemented in function ulliv_up_B.

If the rank of $B$ is known to stay the same after the updating, then the algorithm should accommodate this fact. Ideally, the element $\phi$ in the reduced $f_{2}^{T}$ should be zero, or very small. When this is the case, it is simple to annihilate $f_{1}^{T}$ by means of strategic Givens rotations. Otherwise it is required to apply a ULV rank-revealing step to the $(p+1) \times(p+1)$ matrix

$$
\left(\begin{array}{cc}
L & 0 \\
f_{1}^{T} & \phi
\end{array}\right) .
$$

This step is guaranteed to produce small elements in the bottom row which can therefore safely be neglected. When applied to the full third factor, a single fill is created in position $(p+1, p)$, and we can use strategic Givens rotations to chase this element to the left. In both cases the resulting $U_{B}$ has dimensions $(p+1) \times p$, and thus neither the updated $B$ nor the updated $U_{B}$ conform to our requirements of the ULLIV decomposition. We do not provide implementations.

## 5. A Lanczos Algorithm with Reorthogonalization and Restarts

We now describe our Lanczos-based routine Isvdrr for computing the largest $p$ singular values $\sigma_{i}$ and associated right singular vectors $v_{i}$ of a matrix $A$. Our algorithm uses full reorthogonalization and explicit restarts, and it is based on the work in [6] with modifications that make it more reliable and possibly faster.

If we apply $k$ steps of the Lanczos algorithm $[8, \S \S 9.1-2]$ to the matrix $A^{T} A$ (by first multiplying with $A$ and then by $A^{T}$ ), cf. Fig. 5.1, then in exact arithmetic we produce an $n \times k$ matrix $V_{k}$ with orthonormal columns, and a $k \times k$ symmetric semidefinite tridiagonal matrix $T_{k}$, such that $A V_{k}=V_{k} T_{k}$. Then it is well known that some of the large eigenvalues of $T_{k}$ will approximate some of the large eigenvalues of $A^{T} A$. Since these eigenvalues are the squares of the singular values of $A$, we thus have a basic procedure for iteratively computing approximations to the large singular values of a matrix.

More precisely, let $T_{k}=S_{k} \Theta_{k} S_{k}^{T}$ denote the eigenvalue decomposition of $T_{k}$, and let $\theta_{i}^{(k)}$ denote these eigenvalues. Moreover, let $y_{i}^{(k)}$ denote the columns of the matrix $Y_{k}=V_{k} S_{k}$. Then $\left(\theta_{i}^{(k)}, y_{i}^{(k)}\right)$ are called the Ritz pairs associated with the $k$ th step of the Lanczos process, and some of the Ritz pairs will approximate some of the eigenpairs of $A^{T} A$.

Since this Lanczos algorithm is based on the implicit formation of the matrix $A^{T} A$, there is no guarantee that it can provide accurate estimates of the small singular values of $A$ in finite-precision computations. This does not cause a problem here, however, because our algorithm is intended solely for the computation of the largest singular values.

A more severe difficulty with finite-precision computations in the Lanczos algorithm is that the Lanczos vectors (the columns of $V_{k}$ ) lose orthogonality as the Ritz values converge. This, in turn, leads to various difficulties with repeated and spurious eigenvalues of $T_{k}$ that do not represent approximations to eigenvalues of $A^{T} A$. A number of sophisticated remedies have been proposed for overcoming these difficulties, many of them involving partial or selective reorthogonalizations, combined with methods for monitoring the accuracy of the Ritz values, cf. $[8, \S 9.2]$. With the inclusion of these techniques, the Lanczos algorithm can be used as a general-purpose method for computing, in principle, any portion of the eigenvalue spectrum of $A$.

Our goal here, on the other hand, is to provide a simple Lanczos algorithm solely for computing the largest $p$ singular values of a matrix, suited as a basis for dedicated hardware implementations. For this reason, we use complete reorthogonalization among the Lanczos vectors (which takes place after step 5 in the getrtzp algorithm in Fig. 5.1). As long as $p$ is not large, the additional computational work involved in this approach is acceptable, the actual code is very simple, and the storage requirements for the Lanczos vectors and the converged Ritz vectors are known a priori.

Our stopping criterion for the Lanczos process is based on an estimate of the error in $\left(\theta_{i}^{(k)}\right)^{1 / 2}$, when considered an approximation to $\sigma_{j}$; different indices $i$ and $j$ are needed

## Basic Lanczos Algorithm getrtzp:

Initialization: $\beta_{0} \leftarrow 0 ; v_{0} \leftarrow 0 ; v_{1} \leftarrow$ initial vector.
For $k=1, \ldots, \ell$
$w \leftarrow A^{T}\left(A v_{k}\right)-\beta_{k-1} v_{k-1} ;$
$T_{k, k} \leftarrow \alpha_{k} \leftarrow v_{k}^{T} w ;$
$w \leftarrow \alpha_{k} v_{k}$;
$T_{k, k+1}=T_{k+1, k} \leftarrow \beta_{k} \leftarrow\|w\| ;$
$v_{k+1} \leftarrow w / \beta_{k} ;$
$T_{k}=S_{k} \Theta_{k} S_{k}^{T}$ (eigenvalue decomposition).
Use error estimates $e_{i}^{(k)}(5.1)$ to identify $n_{\mathrm{c}}$ converged Ritz pairs.
Lanczos SVD Alg. w/ Reorthogonalization and Restarts Isvdrr:
Initialization: $v_{\text {init }} \leftarrow A^{T} e ; n_{\text {crp }} \leftarrow 0 ; \mathcal{R} \mathcal{P} \leftarrow \emptyset$ (no Ritz pairs).
2. While $n_{\text {crp }}<k$
use getrtzp to compute $n_{\mathrm{c}}$ Ritz pairs;
$\mathcal{R} \mathcal{P} \leftarrow \mathcal{R} \mathcal{P} \cup\{$ set of new Ritz pairs $\} ;$
$n_{\text {crp }} \leftarrow n_{\text {crp }}+n_{\mathrm{c}} ;$
$v_{\text {init }} \leftarrow v_{k+1}$ from getrtzp.
For $i=1, \ldots, \rho_{0}$
use getrtzp to compute $n_{\mathrm{c}}$ Ritz pairs;
9. if necessary, swap new Ritz pair(s) with pair(s) in $\mathcal{R P}$.

Figure 5.1: Top: the basic algorithm getrtzp for computing Ritz pairs of $A^{T} A$. Bottom: the complete Lanczos algorithm Isvdrr, in which $e=(1, \ldots, 1)^{T}, \mathcal{R} \mathcal{P}$ is the set of converged Ritz pairs, and $n_{\text {crp }}$ is the total number of converged Ritz pairs.
because there is no guarantee that the Ritz values converge in the "natural order." From Theorem 9.1.2 in [8], we know that the error $\sigma_{j}^{2}-\theta_{i}^{(k)}$ in the $i$ th Ritz value is bounded above as

$$
\left|\sigma_{j}^{2}-\theta_{i}^{(k)}\right| \leq\left|\beta_{k} s_{k i}\right|, \quad i=1, \ldots, k
$$

where $\beta_{k}$ is the bottom off-diagonal element of $T_{k}$, and $s_{k i}$ is the $i$ th element of the bottom row of $S_{k}$. If we write $\left(\theta_{i}^{(k)}\right)^{1 / 2}=\tilde{\sigma}_{j}^{(k)}=\sigma_{j}+\delta_{j}^{(k)}$, then

$$
\left|\sigma_{j}^{2}-\theta_{i}^{(k)}\right|=\left|\sigma_{j}+\tilde{\sigma}_{j}^{(k)}\right|\left|\sigma_{j}-\tilde{\sigma}_{j}^{(k)}\right|=\left|2 \tilde{\sigma}_{j}^{(k)}-\delta_{j}^{(k)}\right|\left|\delta_{j}^{(k)}\right| \approx 2 \tilde{\sigma}_{j}^{(k)}\left|\delta_{j}^{(k)}\right|
$$

showing that the quantity

$$
\begin{equation*}
e_{i}^{(k)}=\left|\beta_{k} s_{k i}\right|\left(\theta_{i}^{(k)}\right)^{-1 / 2}, \quad i=1, \ldots, k \tag{5.1}
\end{equation*}
$$

provides an estimate of the error in $\tilde{\sigma}_{i}^{(k)}=\left(\theta_{i}^{(k)}\right)^{1 / 2}$ considered as an approximation to a singular value $\sigma_{j}$ of $A$. Our criterion for accepting a Ritz value as converged is therefore

$$
\begin{equation*}
e_{i}^{(k)}<\tau \sigma_{\max } \tag{5.2}
\end{equation*}
$$

where $\tau$ is a user-specified tolerance, and $\sigma_{\max }$ is an estimate of the largest singular value $\sigma_{1}$.
From Thm. 8.1.2 in [8] we also know that if $\tilde{v}_{i}^{(k)}$ is the approximate eigenvector associated with $\theta_{i}^{(k)}$ then

$$
\left\|A^{T} A \tilde{v}_{i}^{(k)}-\theta_{i}^{(k)} \tilde{v}_{i}^{(k)}\right\|=\left|\beta_{k} x_{k i}\right|=e_{i}^{(k)} \tilde{\sigma}_{i}^{(k)}
$$

showing that small error estimates $e_{i}^{(k)}$ guarantee small residuals. Furthermore, according to Thm. 11.7.2 in [19], small residuals imply a small subspace angle between the subspaces spanned by the exact and approximate eigenvectors.

Unfortunately, the number of Lanczos iterations needed to capture $p$ singular values, within the accuracy estimates provided by (5.1), may exceed $p$ by a large factor. The cure to this difficulty is to restart the Lanczos process with an initial vector that is orthogonal to the set of converged Ritz vectors. This is easily archived in our algorithm, where the Ritz vectors are explicitly saved.

Let $n_{\text {crp }}$ denote the total number of converged Ritz pairs, and let $\ell_{0}$ be a fixed number greater than $p$, chosen by the user. Each time the Lanczos process is (re)started, we perform $\ell$ iterations. In our algorithm, one can either choose $\ell=\ell_{0}$ or $\ell=\ell_{0}-n_{\text {crp }}$. The latter choice, which is default, ensures that a total of $\ell_{0}$ Lanczos vectors are used.

When we have reached a stage where $n_{\operatorname{crp}}=p$ Ritz values $\theta_{i}^{(k)}$ have converged according to (5.2), there is no guarantee that we have computed approximations to the desired $p$ largest singular values. Our heuristic remedy for this difficulty is to restart the Lanczos process additional $\rho_{0}$ times, where $\rho_{0}$ is a small number (the default is $\rho_{0}=2$ ). Experiments in [6] show that these additional restarts indeed improve the reliability of the algorithm, at little extra cost.

Upon completion our algorithm Isvdrr returns approximations $\tilde{\sigma}_{i}$ and $\tilde{v}_{i}$ to the largest $p$ singular values and corresponding right singular vectors. Approximations to the left singular vectors can then be computed as $\tilde{u}_{i}=A \tilde{v}_{i} / \tilde{\sigma}_{i}$, and we emphasize that these vectors are not orthonormal. If an SVD routine is available, one can instead compute a diagonal matrix $\widehat{\Sigma}$ and two matrices $\widehat{U}$ and $\widehat{V}$ with orthonormal columns satisfying $A \widehat{V}=\widehat{U} \widehat{\Sigma}$, by means of the following procedure:

1. $A\left(\tilde{v}_{1}, \ldots, \tilde{v}_{p}\right)=\hat{U} \widehat{\Sigma} \widehat{V}^{T} \quad$ (SVD computation),
2. $\widehat{V} \leftarrow\left(\tilde{v}_{1}, \ldots, \tilde{v}_{p}\right) \widehat{V}$.

More details about this approach can be found in [6], which also includes numerical results concerning the accuracy and efficiency of the Isvdrr algorithm.

## 6. Numerical Examples

We conclude with three demonstrations of the use of our package. The first example is available in the script vsviddemo, and shows how the three routines hvsvid_L, hvsvid_R and Ivsvid can be used to compute symmetric indefinite VSV decompositions (2.1) of rank deficient KKT matrices of the form

$$
A=\left(\begin{array}{cc}
M & N^{T} \\
N & 0
\end{array}\right) .
$$

In our test problems, $M$ is an $m \times m$ symmetric semidefinite and rank-deficient matrix, and $N=\Theta M$ with $\Theta$ a $q \times m$ random matrix. The resulting matrix $A$ has dimension $n=m+q$, and it is rank deficient and symmetric indefinite.

We generate 500 test matrices, and for each matrix and each VSV decomposition we compute the numerical rank $r$, the backward error $\left\|A-V T^{T} \Omega T V^{T}\right\|$, and the subspace angle between the numerical null space (spanned by the last $n-r$ singular vectors of $A$ ) and the approximate null space spanned by the last $n-r$ columns of $V$. A typical example of the results from such a test is shown in Fig. 6.1 for $m=12$ and $q=2$. Occasionally, the errors are in the range $10^{-13}-10^{-9}$, while most of the errors are less than $10^{-13}$.

The second test problem is available in the function ullivdemo and illustrates the use of the ULLIV decomposition is noise and interference reduction. We generate a clean signal $s \in \mathbb{R}^{N}$ of length $N=350$ consisting of a sum of 9 sinusoids with unit amplitude; see the DFT spectrum in the top of Fig. 6.2. The noisy signal is generated by adding white noise and an interfering signal to the clean signal; the white noise is generated by the Matlab command $0.5^{*} \operatorname{randn}(\mathrm{~N}, 1)$, and the interfering signal is a sum of 16 sinusoids with amplitude 0.2 . The DFT spectrum of the noisy signal is shown in the middle of Fig.6.2.

The filtered signal is then computed by means of the subspace method described in [12]. This involves the computation of the ULLIV decomposition (4.1) of two Hankel matrices $A$ and $B$, the first being $311 \times 40$ and consisting of the noisy signal, and the second being $32 \times 40$ and representing the signal subspace of the interfering signal. From the ULLIV decomposition and the numerical rank $p$ we then construct the matrix

$$
X=U_{A} \Psi L_{A}\left(\begin{array}{cc}
L & 0 \\
0 & I
\end{array}\right) V^{T}, \quad \Psi=\operatorname{diag}\left(I_{p}, 0\right) .
$$

Finally we construct the filtered signal by averaging along the antidiagonals of $X$. The DFT spectrum of the filtered signal is shown in the bottom of Fig.6.2, and we see that we have indeed reduced the interference while maintaining most of the clean signal.

In the third test problem, which is available in the script Isvdrrdemo, we compute the $k$ largest singular values and right singular vectors of a complex Toeplitz matrix of size $513 \times 512$ constructed from an NMR signal available as Data Set 002 at the BioSource database [1] of MRS signals. See, e.g., [22] for an application of such computations. The computations are performed for $k=5,10,15$ and 20 , and for each $k$ we compare the errors in the singular values and vectors, as well as the computing times, with those of the Matlab function svds.


Figure 6.1: Backward errors and subspace angles for 500 KKT test problems of dimension $14 \times 14$.


Figure 6.2: DFT spectra of the clean, noisy and filtered signals in ullivdemo test problem.


Figure 6.3: Computing times and errors for the largest $k$ singular values and right singular vectors, computed by means of our Lanczos algorithm Isvdrr as well as Matlab's Lanczos algorithm implemented in the svds function. The bottom left figure reports max $\left|\sigma_{i}-\tilde{\sigma}_{i}\right|$, $i=1, \ldots, p$, while the bottom right figure reports the subspace angle between the subspaces spanned by the dominant $p$ right singular vectors and the approximations $\tilde{v}_{i}, \ldots, \tilde{v}_{p}$.

We used the tolerance $\tau=10^{-10}$ in the convergence criterion (5.2), and the matrix-vector multiplications with the Toeplitz matrix are done via Matlab's fft function. Figure 6.3 shows that for this test problem (using a 1.4 GHz Pentium), Isvdrr computes accurate singular values and right singular vectors faster than the general-purpose Lanczos routine svds. For comparison, Matlab's dense SVD routine computed all the singular values and right singular vectors in 8.5 seconds.

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## 7. Overview of Routines

|  | Demo Functions and Solvers |
| :--- | :--- |
| Isvdrrdemo | Demonstrates the use of the restarted Lanczos algorithm <br> implemented in Isvdrr and applied to NMR data from [1] |
| tvsv | Solves a symmetric num. rank-deficient system of equations <br> ullivdemo |
| Demonstrates the use of the high-rank ULLIV algorithm ulliv, <br> applied to noise and interference reduction |  |
| vsviddemo | Demonstrates the use of the indefinite VSV algorithms hvsvid_L, <br> hvsvid_R and lvsvid applied to symmetric indefinite KKT systems |


| Rank-Revealing Symmetric VSV Algorithms |  |  |
| :--- | :--- | :---: |
| hvsvid_L | High-rank algorithm for indefinite matrix, $L$ version |  |
| hvsvid_R | High-rank algorithm for indefinite matrix, $R$ version |  |
| hvsvsd | High-rank algorithm for semidefinite matrix |  |
| lvsvid | Low-rank algorithm for indefinite matrix |  |
| lvsvsd | Low-rank algorithm for semidefinite matrix |  |
| vsv_qrit | Block QR refinement of VSV decomposition |  |


|  | VSV Up- and Downdating |
| :--- | :--- |
| vsvid_L_mod | Rank-one modification of indefinite matrix, $L$ version |
| vsvid_R_mod | Rank-one modification of indefinite matrix, $R$ version |
| vsvsd_up | Rank-one update of semidefinite matrix |

QLP Algorithms
hqlp High-rank gap-revealing QLP factorization Iqlp QLP matrix approximation of a low-rank matrix

|  | ULLIV Algorithms for a Matrix Pair $(A, B)$ |
| :--- | :--- |
| ulliv | Rank-revealing ULLIV decomposition, $B$ has full row rank |
| ulliv_up_A | Rank-one update of the $A$ matrix |
| ulliv_up_B | Rank-one update of the $B$ matrix (rank increases) |


| Lanczos | Algorithm with Reorthogonalization and Restarts |
| :--- | :--- |
| getrtzp | Computation of Ritz pairs |
| Isvdrr | Driver routine for Lanczos algorithm |
| tprod | Toeplitz matrix-vector multiplication using FFT |
| tprodinit | Initialization routine for tprod |


| Misc. Tools |  |
| :--- | :--- |
| agl5 | Ashcraft-Grimes-Lewis $5 \times 5$ test matrix for LDL $^{T}$ factorization |
| app_hyp | Apply a stabilized hyperbolic rotation |
| app_qrot | Apply a quadratic rotation |
| csi-10-10 | Mat-file with NMR data for lsvdrrdemo |
| gen_hyp | Generate a stabilized hyperbolic rotation |
| gen_qrot | Generate a quadratic rotation |
| hvsvid_cdef | Column deflation of upper triang. matrix in indef. VSV decomp. |
| hvsvid_rdef | Row deflation of lower triang. matrix in indef. VSV decomp. |
| lvsvid_cdef | Column deflation of lower triang. matrix in indef. VSV decomp. |
| TOTinviter | Inverse iterations applied to $T^{T} \Omega T$ |
| TOTlanczos | Lanczos method applied to $T^{T} \Omega T$ |
| TOTpowiter | Power iterations applied to $T^{T} \Omega T$ |
| vsvid_ip | Interim processor for indefinite VSV decomposition |


| Two routines from the Matrix Computation Toolbox [14] |  |  |
| :--- | :--- | :---: |
| cholp | Pivoted Cholesky factorization |  |
| Idlt_symm | LDL $^{T}$ factorization with symmetric or rook pivoting |  |

## agl5

## Purpose

Ashcraft-Grimes-Lewis 5-times-5 test problem for $\mathrm{LDL}^{T}$ factorization

## Synopsis

$[A, L, D]=a g l 5$

## Description

Generates a 5 -times- 5 test problem for the $\mathrm{LDL}^{T}$ factorization, such that L has a large entry when partial pivoting (Bunch-Kaufman) is used.

## References

[1] C. Ashcraft, R.G. Grimes and J.G. Lewis, "Accurate symmetric indefinite linear equation solvers," SIAM J. Matrix Anal. Appl., 20 (1999), pp. 513-561.

## app_hyp

## Purpose

Apply a stabilized hyperbolic rotation (left/right).
Synopsis
[u1,u2] = app_hyp(v1,v2,c,s,sgn)

## Description

Apply a stabilized hyperbolic rotation, defined by the parameters c and s, from the left to the row vectors v 1 and v 2 such that

$$
\begin{aligned}
& \text { [u1] }=\text { [ ch -sh ] [v1] } \\
& \text { [u2] [-sh ch ] [v2] }
\end{aligned}
$$

or from the right to the column vectors v1 and v2 such that

$$
\begin{array}{r}
\left.[\mathrm{u} 1 \mathrm{u} 2]=\left[\begin{array}{ll}
\mathrm{v} 1 & \mathrm{v} 2
\end{array}\right] \begin{array}{cc}
{\left[\begin{array}{cc}
\mathrm{ch} & -\mathrm{sh}
\end{array}\right]} \\
{[-\mathrm{sh}} & \mathrm{ch}
\end{array}\right]
\end{array}
$$

where $\mathrm{ch}=1 / \mathrm{s}$ and $\mathrm{sh}=\mathrm{c} / \mathrm{s}$.

## See Also

gen_hyp - Determine a 2-by-2 stabilized hyperbolic rotation.

## References

[1] S.T. Alexander, C.-T. Pan \& R.J. Plemmons, "Analysis of recursive least squares hyperbolic rotation algorithms for signal processing," Lin. Alg. Appl. 98 (1998), 3-40.

## app_qrot

## Purpose

Apply a quadratic rotation (left/right).
Synopsis
[u1,u2,dd1,dd2] = app_qrot(v1,v2,c,s,d1,d2,sgn)

## Description

Apply a quadratic rotation $H$, defined by the parameters c and s, from the left to the row vectors v 1 and v 2 such that

$$
\begin{array}{cc}
{[\mathrm{u} 1]} \\
{[\mathrm{u} 2]}
\end{array} \quad \mathrm{H}[\mathrm{v} 1]
$$

or from the right to the column vectors v1 and v2 such that

$$
\left[\begin{array}{ll}
u 1 & u 2
\end{array}\right]=\left[\begin{array}{ll}
u 1 & u 2
\end{array}\right] H^{\prime}
$$

Also update the signature matrix:

```
[dd1 0 ] = sgn*[[d1 0 ]
[ 0 dd2] [ 0 0 d2]
```

where sgn is determined by the rotation.

## See Also

gen_qrot $\quad-\quad$ Determine a 2 -by-2 quadratic rotation.

## gen_hyp

## Purpose

Determine a 2-by-2 stabilized hyperbolic rotation matrix.

## Synopsis

$[c, s, r, s g n]=$ gen_hyp $(a, b)$

## Description

Compute a stabilized hyperbolic rotation to annihilate b using a, i.e., compute parameters $\mathrm{c}, \mathrm{s}$, and r such that

$$
\begin{aligned}
& \text { [ ch -sh ] [a] = [r] with [ ch -sh ] S [ ch -sh ] = sgn*S } \\
& {\left[\begin{array}{ll}
- \text { sh } & \mathrm{ch}
\end{array}\right][\mathrm{b}] \quad[0] \quad[-\mathrm{sh} \text { ch }] \quad[-\mathrm{sh} \text { ch }]}
\end{aligned}
$$

where $\mathrm{ch}=1 / \mathrm{s}$ and $\mathrm{sh}=\mathrm{c} / \mathrm{s}$, and where the signature matrix S is either

$$
\begin{aligned}
S= & {\left[\begin{array}{lll}
-1 & 0
\end{array}\right] } \\
& {\left[\begin{array}{lll}
0 & 1
\end{array}\right] }
\end{aligned} \quad \text { or } \quad S=\left[\begin{array}{rrr}
1 & 0
\end{array}\right] .
$$

## See Also

app_hyp - Apply a stabilized hyperbolic rotation.

## References

[1] S.T. Alexander, C.-T. Pan \& R.J. Plemmons, "Analysis of recursive least squares hyperbolic rotation algorithms for signal processing," Lin. Alg. Appl. 98 (1998), 3-40.

## gen_qrot

## Purpose

Determine a 2 -by- 2 quadratic rotation matrix.

## Synopsis

[c,s,r,sgn] = gen_qrot(a,b,d1,d2)

## Description

Compute a real quadratic (Givens or hyperbolic) rotation H to annihilate b using a , i.e., compute $\mathrm{c}, \mathrm{s}$, and r such that

```
H [a] = [r] with H' S H = sgn*S
[b] [0]
```

where the signature matrix S is

$$
\begin{aligned}
S= & {\left[\begin{array}{cc}
d 1 & 0
\end{array}\right], \quad \mathrm{d} 1, \mathrm{~d} 2=+1,-1 . } \\
& {\left[\begin{array}{ll}
0 & d 2
\end{array}\right] }
\end{aligned}
$$

See Also
app_qrot - Apply a quadratic rotation.

## getrtzp

## Purpose

Compute (additional) Ritz pairs for a cross-product matrix

## Synopsis

[rtzvals,rtzvecs,errests,nconv,num_iter,vnext,smax] = ...
getrtzp(A,k,ncrp,Vk,max_iter,vinit,tol,smax)

## Description

Applies Lanczos iterations to the matrix $\mathrm{A}^{\prime} * \mathrm{~A}$; the Lanczos vectors are explicitly reorthogonalized internally, as well as with respect to an existing set of converged vectors. A singular value estimate is considered as converged when an error estimate is below smax*tol.

## Input Parameters

| A | matrix; |
| :--- | :--- |
| k | number of desired singular values; |
| ncrp | number of converged Ritz pairs so far; |
| Vk | previously converged Ritz vectors; |
| max_iter | max. no. of Lanczos iterations in this call to getrtzp; |
| vinit | start vector; |
| tol | relative residual tolerance; |
| smax | current estimate of largest singular value; |

## Output Parameters

| rtzvals | converged Ritz values og A' $*$ A; |
| :--- | :--- |
| rtzvecs | corresponding converged Ritz vectors; |
| errests | corresponding error estimates of singular values; |
| nconv | number of converged Ritz pairs; |
| num_iter | number of iterations used in the call to getrtzp; |
| vnext | vector for next call to getrtzp; |
| smax | updated value of smax; |

## References

[1] R.D. Fierro and E.P. Jiang, "Lanczos and the Riemannian SVD in information retrieval applications," Numer. Lin. Alg. Appl., to appear.

## hqlp

## Purpose

High-rank gap-revealing QLP factorization.

## Synopsis

$[p, L, P, Q]=h q l p(A)$
$[p, L, P, Q]=$ hqlp(A,gap_tol)

## Description

Computes a gap-revealing factorization $\mathrm{A}=\mathrm{Q} * \mathrm{~L} * \mathrm{P}^{\prime}$, in which Q and P are orthogonal matrices, and L is a lower triangular matrix whose diagonal elements sometimes approximate the singular values of $A$. Also returns the largest $p$ such that $\operatorname{abs}(\mathrm{L}(\mathrm{p}, \mathrm{p}) / \mathrm{L}(\mathrm{p}+1, \mathrm{p}+1))>$ gap_tol. Designed for high-rank matrices; use lqlp for lowrank matrices.

## Input Parameters

A general matrix;
gap_tol tolerance for gap detection;
Defaults $\quad$ gap_tol $=\min (\operatorname{size}(A)) /$ eps;

## Output Parameters

$\mathrm{p} \quad$ estimate of numerical rank of A ;
$\mathrm{L} \quad$ lower triangular matrix in $\mathrm{A}=\mathrm{Q} * \mathrm{~L} * \mathrm{P}^{\prime}$;
P right orthogonal matrix;
Q left orthogonal matrix;

## Algorithm

A pivoted QR factorization $\mathrm{A} * \mathrm{Pi}_{-} \mathrm{p}=\mathrm{Q} * \mathrm{R}$ is followed by a pivoted QR factorization $\mathrm{R}^{\prime} * \mathrm{Pi}_{-\mathrm{q}}=\mathrm{P} * \mathrm{~L}{ }^{\prime}$; thus $\mathrm{A}=\left(\mathrm{Q} * \mathrm{Pi}_{\_\mathrm{q}}\right) * \mathrm{~L} *\left(\mathrm{Pi} \mathrm{i}_{-\mathrm{p}} * \mathrm{P}\right)^{\prime}$. The diagonal elements of L sometimes track the singular values of A , but this is not guaranteed; hence the factorization cannot be guaranteed to reveal rank.

## See Also

Iqlp - pivoted QLP matrix approximation with interleaved factorizations.

## References

[1] G.W. Stewart, "The QLP approximation to the singular value decomposition," SIAM J. Sci. Comp., 20 (1999), pp. 1336-1348.

## hvsvid_L

## Purpose

High-rank revealing decomp. of a sym. indef. matrix, $L$ version.

## Synopsis

[p,L,Omega, V] = hvsvid_L(A)
[p,L,Omega,V] = hvsvid_L(A,tol_rank)
[p,L,Omega, V$]=$ hvsvid_L(A,tol_rank,max_iter)
[p,L,Omega,V] = hvsvid_L(A,tol_rank,max_iter,fixed_rank)

## Description

Computes a rank-revealing VSV decompostion $\mathrm{A}=\mathrm{V} *\left(\mathrm{~L}^{\prime} *\right.$ Omega $\left.* \mathrm{~L}\right) * \mathrm{~V}^{\prime}$ of a symmetric indefinite n-by-n matrix. Only the upper triangular part needs to be specified. Optimized for matrices whose rank $p$ close to $n$. Function hvsvid_R computes the $R$ version.

## Input Parameters

A
tol_rank
max_iter max. number of inverse iterations per deflation step, used in the singular vector estimator;
fixed_rank deflate to the fixed rank given by fixed_rank instead of using the rank decision tolerance;

Defaults tol_rank $=\mathrm{n} *$ norm $(\mathrm{A}, 1) * e \mathrm{ps}$;
max_iter $=5$;

## Output Parameters

$p \quad$ numerical rank of A;
$\mathrm{L} \quad$ lower triangular matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{~L}^{\prime} *\right.$ Omega $\left.* \mathrm{~L}\right) * \mathrm{~V}^{\prime}$;
Omega signature matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{~L}^{\prime} *\right.$ Omega $\left.* \mathrm{~L}\right) * \mathrm{~V}^{\prime}$;
V

$$
\text { orthogonal matrix in } \mathrm{A}=\mathrm{V} *\left(\mathrm{~L}^{\prime} * \text { Omega } * \mathrm{~L}\right) * \mathrm{~V}^{\prime}
$$

## Algorithm

The symmetric indefinite matrix A is preprocessed by a pivoted LDL' factorization. An interim stage (where D is made diagonal) is followed by a rank-revealing ULV-like decomposition, using inverse iterations for singular vector estimation.

See Also
hvsvid_R - High-rank revealing VSV alg. for sym. indef. matrices, R version hvsvsd - High-rank revealing VSV alg. for symmetric semidefinite matrices Ivsvid - Low-rank revealing VSV alg. for symmetric indefinite matrices Ivsvsd - Low-rank revealing VSV alg. for symmetric semidefinite matrices

## References

[1] P.C. Hansen \& P.Y. Yalamov, "Computing Symmetric Rank-Revealing Decompositions via Triangular Factorization", SIAM J. Matrix Anal. Appl., 23 (2001), pp. 443-458.

## hvsvid_R

## Purpose

High-rank revealing decomp. of a sym. indef. matrix, R version.

## Synopsis

$[p, R, O m e g a, V]=$ hvsvid_R(A)
$[p, R, O m e g a, V]=$ hvsvid_R(A,tol_rank)
$[p, R, O m e g a, V]=$ hvsvid_R(A,tol_rank,max_iter)
$[p, R, O m e g a, V]=$ hvsvid_R(A,tol_rank,max_iter,fixed_rank)

## Description

Computes a rank-revealing VSV decompostion $\mathrm{A}=\mathrm{V} *\left(\mathrm{R}^{\prime} *\right.$ Omega $\left.* \mathrm{R}\right) * \mathrm{~V}^{\prime}$ of a symmetric indefinite n-by-n matrix. Only the upper triangular part needs to be specified. Optimized for matrices whose rank p is close to $n$. Functions hvsvid and lvsvid compute the L-version of the decomposition.

## Input Parameters

A
tol_rank
max_iter
fixed_rank deflate to the fixed rank given by fixed_rank instead of using the rank decision tolerance;

Defaults tol_rank $=\mathrm{n} *$ norm $(\mathrm{A}, 1) * e \mathrm{ps}$;
max_iter $=5$;

## Output Parameters

$p$ numerical rank of $A$;
$\mathrm{R} \quad$ upper triangular matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{R}{ }^{\prime} *\right.$ Omega $\left.* \mathrm{R}\right) * \mathrm{~V}^{\prime}$;
Omega signature matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{R}^{\prime} *\right.$ Omega $\left.* \mathrm{R}\right) * \mathrm{~V}^{\prime}$;
V
orthogonal matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{R}^{\prime} *\right.$ Omega $\left.* \mathrm{R}\right) * \mathrm{~V}^{\prime}$;

## Algorithm

The symmetric indefinite matrix A is preprocessed by a pivoted LDL' factorization. An interim stage (where D is made diagonal) is followed by a rank-revealing URV-like decomposition, using inverse iterations for singular vector estimation.

See Also
hvsvid_L - High-rank revealing VSV alg. for sym. indef. matrices, L version
hvsvsd - High-rank revealing VSV alg. for symmetric semidefinite matrices
Ivsvid - Low-rank revealing VSV alg. for symmetric indefinite matrices
Ivsvsd - Low-rank revealing VSV alg. for symmetric semidefinite matrices

## References

[1] P.C. Hansen \& P.Y. Yalamov, "Computing Symmetric Rank-Revealing Decompositions via Triangular Factorization", SIAM J. Matrix Anal. Appl., 23 (2001), pp. 443-458.

## hvsvid cdef

## Purpose

Deflate one column of R in the high-rank URV-like VSV algorithm.

## Synopsis

[R,Omega, V ,fail] = hvsvid_cdef(R,Omega, $\mathrm{V}, \mathrm{r}, \mathrm{vmin}$ )

## Description

Given the decomposition $\mathrm{V} * \mathrm{R}{ }^{\prime} *$ Omega $* \mathrm{R} * \mathrm{~V}^{\prime}$, the function deflates the last column of $R(1: r, 1: r)$. vmin is an estimate of the right singular vector of $\mathrm{R}(1: r, 1: r)^{\prime} * \operatorname{Omega}(1: r, 1: r) * \mathrm{R}(1: r, 1: r)$ associated with the smallest singular value sigma_r. On return, the norm of the last column of $\mathrm{R}(1: r, 1: r)^{\prime} * \operatorname{Omega}(1: r, 1: r) * R(1: r, 1: r)$ is of the order sigma_r. The matrix V can be left out by inserting an empty matrix [].

## Input Parameters

R
Omega
V
r
vmin
upper triangular matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{R}^{\prime} *\right.$ Omega $\left.* \mathrm{R}\right) * \mathrm{~V}^{\prime}$;
signature matrix in $A=V *\left(R^{\prime} *\right.$ Omega $\left.* R\right) * V^{\prime}$; orthogonal matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{R}^{\prime} * O \operatorname{meg} a * \mathrm{R}\right) * \mathrm{~V}^{\prime}$; size of submatrix to be deflated;
estimate of the smallest right singular vector of the product $\mathrm{R}(1: r, 1: r)^{\prime} * \operatorname{Omega}(1: r, 1: r) * \mathrm{R}(1: r, 1: r)$;

## Output Parameters

$\mathrm{R} \quad$ upper triangular matrix in resulting $\mathrm{A}=\mathrm{V} *\left(\mathrm{R}^{\prime} *\right.$ Omega $\left.* \mathrm{R}\right) * \mathrm{~V}^{\prime}$;
Omega signature matrix in resulting $A=V *\left(R^{\prime} *\right.$ Omega $\left.* R\right) * V^{\prime}$;
$\mathrm{V} \quad$ orthogonal matrix in resulting $\mathrm{A}=\mathrm{V} *\left(\mathrm{R}{ }^{\prime} *\right.$ Omega $\left.* \mathrm{R}\right) * \mathrm{~V}^{\prime}$;
fail true if a hypernormal rotation is ill defined;

## References

[1] P.C. Hansen \& P.Y. Yalamov, "Computing Symmetric Rank-Revealing Decompositions via Triangular Factorization", SIAM J. Matrix Anal. Appl., 23 (2001), pp. 443-458.

## hvsvid_rdef

## Purpose

Deflate one row of L in the high-rank ULV-like VSV algorithm.

## Synopsis

[L,V,Omega] = hvsvid_rdef(L,V,Omega,r,vmin)

## Description

Given the decomposition $\mathrm{V} * \mathrm{~L}^{\prime} * \mathrm{Omega} * \mathrm{~L} * \mathrm{~V}^{\prime}$, the function deflates the last row of $\mathrm{L}(1: r, 1: r)$. vmin is an estimate of the right singular vector of
$\mathrm{L}(1: r, 1: r)^{\prime} * \operatorname{Omega}(1: r, 1: r) * \mathrm{~L}(1: r, 1: r)$ associated with the smallest singular value sigma_r. On return, the norm of the last column of $\mathrm{L}(1: r, 1: r)^{\prime} * \operatorname{Omega}(1: r, 1: r) * \mathrm{~L}(1: r, 1: r)$ is of the order sigma_r. The matrix V can be left out by inserting an empty matrix [].

## Input Parameters

```
L lower triangular matrix in A = V *(L'*Omega*L)*V';
Omega signature matrix in A = V *(L'*Omega*L)*V';
V orthogonal matrix in A = V *(L'*Omega*L)*V';
r size of submatrix to be deflated;
vmin estimate of the smallest right singular vector of
    the product L(1:r,1:r)'*Omega(1:r,1:r)*L(1:r,1:r);
```


## Output Parameters

$\mathrm{L} \quad$ upper triangular matrix in resulting $\mathrm{V} *\left(\mathrm{~L}^{\prime} *\right.$ Omega $\left.* \mathrm{~L}\right) * \mathrm{~V}^{\prime}$;
$\mathrm{V} \quad$ orthogonal matrix in resulting $\mathrm{V} *(\mathrm{~L} * *$ Omega $* \mathrm{~L}) * \mathrm{~V}^{\prime}$;
Omega signature matrix in resulting $\mathrm{V} *(\mathrm{~L} * *$ Omega $* \mathrm{~L}) * \mathrm{~V}^{\prime}$;

## References

[1] P.C. Hansen \& P.Y. Yalamov, "Computing Symmetric Rank-Revealing Decompositions via Triangular Factorization", SIAM J. Matrix Anal. Appl., 23 (2001), pp. 443-458.

## hvsvsd

## Purpose

High-rank revealing decompostion of a symmetric semidefinite matrix.

## Synopsis

$[p, L, V]=\operatorname{hvsvsd}(A)$
$[p, L, V]=$ hvsvsd(A,tol_rank)
$[p, L, V]=$ hvsvsd(A,tol_rank,fixed_rank)

## Description

Computes a rank-revealing VSV decomposition $\mathrm{A}=\mathrm{V} *\left(\mathrm{~L}^{\prime} * \mathrm{~L}\right) * \mathrm{~V}^{\prime}$ of a symmetric semidefinite n-by-n matrix. Only the upper triangular part needs to be specified. Optimized for matrices whose rank p is close to n .

## Input Parameters

A
tol_rank
fixed_rank

Defaults
symmetric semidefinite matrix; rank decision tolerance; deflate to the fixed rank given by fixed_rank instead of using the rank decision tolerance; tol_rank $=\mathrm{n} *$ norm $(\mathrm{A}, 1) * \mathrm{eps} ;$

## Output Parameters

$p$ numerical rank of $A$;
$\mathrm{L} \quad$ lower triangular matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{~L}^{\prime} * \mathrm{~L}\right) * \mathrm{~V}^{\prime}$;
$\mathrm{V} \quad$ orthogonal matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{~L}^{\prime} * \mathrm{~L}\right) * \mathrm{~V}^{\prime}$;

## Algorithm

The symmetric semidefinite matrix A is preprocessed by a pivoted Cholesky factorization and then postprocessed by a high-rank-revealing ULV decomposition. An indefinite matrix results in an error message during the Cholesky factorization.

## See Also

hvsvid_L - High-rank revealing VSV alg. for sym. indef. matrices, L version hvsvid_R - High-rank revealing VSV alg. for sym. indef. matrices, R version Ivsvid - Low-rank revealing VSV alg. for symmetric indefinite matrices Ivsvsd - Low-rank revealing VSV alg. for symmetric semidefinite matrices

## References

[1] P.C. Hansen \& P.Y. Yalamov, "Computing Symmetric Rank-Revealing Decompositions via Triangular Factorization", SIAM J. Matrix Anal. Appl., 23 (2001), pp. 443-458.

## lqlp

## Purpose

Pivoted QLP matrix approximation with interleaved factorizations.

## Synopsis

$[p, L, P, Q$, gap_ratio $]=\operatorname{lq} \mid p(A)$
$[p, L, P, Q$, gap_ratio $]=\operatorname{lq} \mid p(A$, tol_gap $)$
$[p, L, P, Q$, gap_ratio $]=\operatorname{lqlp}(A$, tol_gap,fixed_rank)

## Description

Computes a rank-p pivoted QLP matrix approximation of an m-by-n matrix $A(m \geq n)$ satisfying $A * P=Q * L$ with a lower triangular p-by-p matrix $L$. The rank (or stopping point) p is either fixed_rank or is dynamically determined by tol_gap. The absolute value of the diagonal elements of L are approximations to the first p singular values of A, while the columns of Q and P approximate the first p left and right singular vectors of A .

## Input Parameters

tol_gap truncate the decomposition after compuing a rank-p approximation to A , where p is the smallest
integer such that $\operatorname{abs}(\mathrm{L}(\mathrm{p}, \mathrm{p}) / \mathrm{L}(\mathrm{p}+1, \mathrm{p}+1)) \geq$ tol_gap;
fixed_rank ignore tol_gap and truncate the decomposition after
computing an approximation to A of rank fixed_rank.
Defaults tol_gap $=\mathrm{n} / \mathrm{eps}$;
fixed_rank $=\mathrm{n}$.

## Output Parameters

smallest integer such that $\operatorname{abs}(\mathrm{L}(\mathrm{p}, \mathrm{p}) / \mathrm{L}(\mathrm{p}+1, \mathrm{p}+1))$ i tol_gap (or fixed_rank);
L p-by-p lower triangular matrix whose diagonal elements, in absolute value, track the largest $p$ singular values of $A$;
$\mathrm{P}, \mathrm{Q} \quad$ matrices with p orthonormal columns;
gap_ratio $\quad \operatorname{abs}(\mathrm{L}(\mathrm{p}+1, \mathrm{p}+1) / \mathrm{L}(\mathrm{p}, \mathrm{p}))$, that is, the ratio of the first approximate singular value excluded to the last one included, empty if $\mathrm{p}=\mathrm{n}$ or fixed_rank $=\mathrm{n}$.

## Algorithm

The first $p$ rows and columns of the pivoted QR factorization of A are computed, $\mathrm{A} * \mathrm{Pi}_{-} 1=\mathrm{Q} * \mathrm{R}$. Then the pivoted QR factorization of $\mathrm{R}^{\prime}$ is computed, $\mathrm{R}{ }^{\prime} * \mathrm{Pi}_{-} 2=\mathrm{P} * \mathrm{~L}^{\prime}$, where $L$ ' is p-by-p. The rank $p$ is either fixed_rank or is determined dynamically by the following. The computation of R is stopped after k 1 rows and columns, where $\operatorname{abs}(\mathrm{R}(\mathrm{k} 1, \mathrm{k} 1) / \mathrm{R}(\mathrm{k} 1-1, \mathrm{k} 1-1)) \leq$ tol_gap. Then rows and columns of L' are computed to see whether $\operatorname{abs}(\mathrm{L}(\mathrm{j}, \mathrm{j}) / \mathrm{L}(\mathrm{j}-1, \mathrm{j}-1)) \leq$ tol_gap for any $\mathrm{j} \leq \mathrm{k} 1$. If so, the algorithm halts after computing these j rows and columns of L , and the final approximation to the SVD of A is rank j . If not, the next k 2 rows and columns of R are computed until tol_gap is achieved, then corresponding rows and columns of $\mathrm{L}(\mathrm{k} 1+\mathrm{j}, 1 \leq \mathrm{j} \leq \mathrm{k} 2)$ are computed to see whether tol_gap holds at any j and the computation can be halted. If not, more rows and columns of R are computed, etc. This process is called interleaving.

## See Also

hqlp - high-rank gap-revealing QLP factorization.

## References

[1] G.W. Stewart, "The QLP approximation to the singular value decomposition," SIAM J. Sci. Comp., 20 (1999), pp. 1336-1348.
[2] D.A. Huckaby and T.F. Chan, "Stewart's pivoted QLP decomposition for low-rank matrices," Tech. Report CAM-02-54, Dept. Mathematics, UCLA, 2002.

## Isvdrr

## Purpose

Lanczos SVD with reorthogonalization and explicit restarts

## Synopsis

[sk,Vk,nits,nrst,errests] = Isvdrr(A,k,max_iter,max_restarts,conserve,safety)

## Description

Computes approximations to the k dominant singular triplets, using the Lanczos method with reorthogonalization and explicit restarts. The stopping criterion is that the error in each computed singular value must be smaller than tol times the largest singular value.

## Input Parameters

A
k number of desired singular triplets;
tol tolerance for relative residual of triplets;
max_iter maximum number of Lanczos iterations per restart;
max_restarts maximum number of Lanczos restarts;
conserve if 1 then max_iter Lanczos vectors are used in total, otherwise max_iter vectors are used in each restart;
safety perform additional safety restarts, after k Ritz values have converged;

Defaults $\quad$ tol $=1 \mathrm{e}-4$;
max_iter $=\min (2 * \mathrm{k}, \mathrm{n})$
max_restarts $=100$;
conserve $=1$;
safety $=2$;

## Output Parameters

sk vector of singular value estimates;
Vk matrix of right singular vector approximations;
nits number of times A and A' combined have been referenced;
nrst number of restarts;
errests vector of error estimates for singular values;

## Matrix Representation

The input parameter A can be either a matrix (dense or sparse) or a structure that holds information about a Toeplitz matrix: A.col $=$ first column of A, A.row $=$ first row of A .

## References

[1] R.D. Fierro and E.P. Jiang, "Lanczos and the Riemannian SVD in information retrieval applications," Numer. Lin. Alg. Appl., to appear.

## Ivsvid

## Purpose

Low-rank revealing decompostion of a symmectric indefinite matrix.

## Synopsis

$[p, L, O m e g a, V]=\operatorname{Ivsvid}(A)$
$[p, L, O m e g a, \mathrm{~V}]=\operatorname{lvsvid}(\mathrm{A}$, tol_rank)
$[p, L, O m e g a, V]=$ Ivsvid(A,tol_rank,max_iter)
[p,L,Omega,V] = Ivsvid(A,tol_rank,max_iter,est_type)
[p,L,Omega,V] = Ivsvid(A,tol_rank,max_iter,est_type,fixed_rank)

## Description

Computes a rank-revealing VSV decompostion $A=V *\left(\mathrm{R}^{\prime} * \operatorname{Omega} * \mathrm{R}\right) * \mathrm{~V}^{\prime}$ of a symmetric indefinite n-by-n matrix. Only the upper triangular part needs to be specified. Optimized for matrices whose rank p is small compared to $n$.

## Input Parameters

A
tol_rank
max_iter max. number of power/Lanczos iterations per deflation step, used in the singular vector estimator;
est_type if true, then estimate singular vectors by means of the Lanczos procedure, else use the power method;
fixed_rank deflate to the fixed rank given by fixed_rank instead of using the rank decision tolerance;

Defaults tol_rank $=\mathrm{n} *$ norm $(\mathrm{A}, 1) * e \mathrm{ps}$;
max_iter $=5$;
est_type $=0($ power method $) ;$

## Output Parameters

$p \quad$ numerical rank of A;
$\mathrm{L} \quad$ lower triangular matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{~L}^{\prime} * \operatorname{Omega} * \mathrm{~L}\right) * \mathrm{~V}^{\prime}$;
Omega signature matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{~L}^{\prime} *\right.$ Omega $\left.* \mathrm{~L}\right) * \mathrm{~V}^{\prime}$;
$\mathrm{V} \quad$ orthogonal matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{~L}^{\prime} *\right.$ Omega $\left.* \mathrm{~L}\right) * \mathrm{~V}^{\prime}$;
See Also
hvsvid - High-rank revealing VSV alg. for symmetric indefinite matrices
hvsvid_R - High-rank revealing VSV alg., sym. indef. matrices, R version
hvsvsd - High-rank revealing VSV alg. for symmetric semidefinite matrices
lvsvsd - Low-rank revealing VSV alg. for symmetric semidefinite matrices

## References

[1] P.C. Hansen \& P.Y. Yalamov, "Computing Symmetric Rank-Revealing Decompositions via Triangular Factorization", SIAM J. Matrix Anal. Appl., 23 (2001), pp. 443-458.

## Ivsvid_cdef

## Purpose

Deflate one column of L in the low-rank ULV-like VSV algorithm.

## Synopsis

$[\mathrm{L}$, Omega, V$]=$ Ivsvid_cdef(L,Omega, V, ,r,vmax $)$

## Description

Given the VSV decomposition $\mathrm{V} * \mathrm{~L}{ }^{\prime} *$ Omega $* \mathrm{~L} * \mathrm{~V}^{\prime}$, the function deflates the first column of $L(r: n, r: n)$. vmax is an estimate of the right singular vector of
$\mathrm{L}(\mathrm{r}: \mathrm{n}, \mathrm{r}: \mathrm{n})^{\prime} *$ Omega(r:n,r:n) $* \mathrm{~L}(\mathrm{r}: \mathrm{n}, \mathrm{r}: \mathrm{n})$ associated with the largest singular value sigma_1. On return, the norm of the first column of $\mathrm{L}(\mathrm{r}: \mathrm{n}, \mathrm{r}: \mathrm{n})^{\prime} * \operatorname{Omega}(\mathrm{r}: \mathrm{n}, \mathrm{r}: \mathrm{n}) * \mathrm{~L}(\mathrm{r}: \mathrm{n}, \mathrm{r}: \mathrm{n})$ is of the order sigma_1. The matrix V can be left out by inserting an empty matrix [].

## Input Parameters

```
L lower triangular matrix in A = V *(L'*Omega*L)*V';
```

Omega signature matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{~L}^{\prime} *\right.$ Omega $\left.* \mathrm{~L}\right) * \mathrm{~V}^{\prime}$;
$\mathrm{V} \quad$ orthogonal matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{~L}^{\prime} *\right.$ Omega $\left.* \mathrm{~L}\right) * \mathrm{~V}^{\prime}$;
$r$ size of submatrix to be deflated;
vmin estimate of the smallest right singular vector of
the product $\mathrm{L}(1: r, 1: r)^{\prime} * \operatorname{Omega}(1: r, 1: r) * \mathrm{~L}(1: r, 1: r)$;

## Output Parameters

$\mathrm{L} \quad$ lower triangular matrix in resulting $\mathrm{A}=\mathrm{V} *\left(\mathrm{~L}^{\prime} *\right.$ Omega $\left.* \mathrm{~L}\right) * \mathrm{~V}^{\prime}$;
Omega signature matrix in resulting $\mathrm{A}=\mathrm{V} *\left(\mathrm{~L}^{\prime} *\right.$ Omega $\left.* \mathrm{~L}\right) * \mathrm{~V}^{\prime}$;
$\mathrm{V} \quad$ orthogonal matrix in resulting $\mathrm{A}=\mathrm{V} *\left(\mathrm{~L}^{\prime} *\right.$ Omega $\left.* \mathrm{~L}\right) * \mathrm{~V}^{\prime}$;
$\mathrm{V} \quad$ rthogonal matrix in resulting $\mathrm{A}=\mathrm{V} *\left(\mathrm{~L}^{\prime} *\right.$ Omega $\left.* \mathrm{~L}\right) * \mathrm{~V}^{\prime}$;

## References

[1] R.D. Fierro and P.C. Hansen, "Low-Rank Revealing UTV Decompositions", Numerical Algorithms, 15 (1997), pp. 37-55.

## Ivsvsd

## Purpose

Low-rank revealing decompostion of a symmetric semidefinite matrix.

## Synopsis

$$
\begin{aligned}
& {[p, L, V]=\operatorname{lvsvsd}(A)} \\
& {[p, L, V]=\operatorname{lvsvsd}(A, \text { tol_rank })} \\
& {[p, L, V]=\operatorname{lvsvsd}(A, \text { tol_rank,_max_iter })} \\
& {[p, L, V]=\operatorname{lvsvsd}(A, \text { tol_rank,max_iter,est_type })} \\
& {[p, L, V]=\operatorname{lvsvsd}(A, \text { tol_rank,_max_iter,est_type,fixed_rank })}
\end{aligned}
$$

## Description

Computes a rank-revealing VSV decompostion $A=V *\left(L^{\prime} * L\right) * V^{\prime}$ of a symmetric semidefinite n-by-n matrix. Only the upper triangular part of needs to be specified. Optimized for matrices whose rank p is small compared to n .

## Input Parameters

A
tol_rank rank decision tolerance;
max_iter max. number of inverse iterations per deflation step, used in the singular vector estimator;
est_type if true, then estimate singular vectors by means of the Lanczos procedure, else use the power method;
fixed_rank deflate to the fixed rank given by fixed_rank instead of using the rank decision tolerance;

Defaults tol_rank $=\mathrm{n} *$ norm $(\mathrm{A}, 1) * e \mathrm{ps}$;
max_iter $=5$;
est_type $=0($ power method $) ;$

## Output Parameters

| $p$ | numerical rank of $A ;$ |
| :--- | :--- |
| $L$ | lower triangular matrix in $A=V *\left(L^{\prime} * \mathrm{~L}\right) * \mathrm{~V}^{\prime} ;$ |
| V | orthogonal matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{~L}^{\prime} * \mathrm{~L}\right) * \mathrm{~V}^{\prime} ;$ |

## Algorithm

The symmetric semidefinite matrix A is preprocessed by a pivoted Cholesky factorization and then postprocessed by a low-rank revealing ULV decomposition, using either power or Lanczos iterations to estimate the dominant singular vectors. An indefinite matrix results in an error message during the Cholesky factorization.

## See Also

hvsvid - High-rank revealing VSV alg. for symmetric indefinite matrices hvsvid_R - High-rank revealing VSV alg., sym. indef. matrices, R version hvsvsd - High-rank revealing VSV alg. for symmetric semidefinite matrices Ivsvid - Low-rank revealing VSV alg. for symmetric indefinite matrices

## References

[1] P.C. Hansen \& P.Y. Yalamov, "Computing Symmetric Rank-Revealing Decompositions via Triangular Factorization", SIAM J. Matrix Anal. Appl., 23 (2001), pp. 443-458.

## TOTinviter

## Purpose

Inverse iterations for T'*Omega*T.
Synopsis
[smin,vmin] $=$ TOTinviter(T,Omega,itmax)

## Description

Inverse iterations on the product $\mathrm{T}^{\prime} *$ Omega $* \mathrm{~T}$ to compute the smallest singular value and the corresponding singular vector.

## Input Parameters

T
Omega
triangular matrix;
diagonal matrix;
itmax maximum number of iterations;

## Output Parameters

smin estimate of smallest singular value;
vmin estimate of corresponding singular vector;

## TOTlanczos

## Purpose

Symmetric Lanczos procedure for T* $*$ Omega*T.

## Synopsis

[smax,vmax] = TOTlanczos(T,Omega,itmax)
[smax,vmax] = TOTlanczos(T,Omega,itmax,reorth)

## Description

Computes an estimate of the largest singular value and the associated singular vector of the matrix T'*Omega*T using a maximum of itmax Lanczos iterations. If reorth $=$ 1 , then MGS reorthogonalization is used.

## Input Parameters

T matrix;

Omega diagonal matrix;
itmax maximum number of iterations;
reorth MGS reorthogonalization if true;
Defaults reorth $=1$ (reorthogonalization).

## Output Parameters

smin estimate of smallest singular value;
vmin estimate of corresponding singular vector;
See Also
TOTpowiter - Power iterations for $\mathrm{T}^{*} *$ Omega $* \mathrm{~T}$

## References

[1] G.H. Golub and C.F. Van Loan, "Matrix Computations", Johns Hopkins University Press, 3. Ed., 1996; p. 480.

## TOTpowiter

## Purpose

Power iterations for $\mathrm{T}^{*} *$ Omega*T.
Synopsis
[smin,vmin] = TOTpowiter(T,Omega,itmax)

## Description

Power iterations on the product $\mathrm{T}^{\prime} *$ Omega $* \mathrm{~T}$ to compute the largest singular value and the corresponding singular vector.

## Input Parameters

T triangular matrix;
Omega diagonal matrix;
itmax maximum number of iterations;

## Output Parameters <br> smin estimate of largest singular value; <br> vmin estimate of corresponding singular vector;

See Also<br>TOTlanczos - Lanczos procedure for T'*Omega*T

## tprod

## Purpose

Toeplitz matrix-vector multiplication via FFT.
Synopsis
$\mathrm{y}=\operatorname{tprod}(\operatorname{lambda}, \mathrm{m}, \mathrm{n}, \mathrm{x}$, transp$)$

## Description

This routine computes $\mathrm{T} * \mathrm{x}$ for tranps $=0$ or $\mathrm{T}^{\prime} * \mathrm{x}$ for value $=1$, where T is an m -by-n Toeplitz matrix, using the eigenvalues lambda of a related circulant matrix computed ny means of tprodinit.

## Input Parameters

 lambda eigenvalue vector need for the FFT; $\mathrm{m}, \mathrm{n}$ dimensions of Toeplitz matrix; $x \quad$ vector to be multiplied by T ;transp if 0 or nonexisting, multiply with T , ortherwise multiply with T';

## Output Parameters

y matrix-vector product

## Algorithm

Let lambda be the eigenvalues of a circulant matrix derived from T (see tprodinit); then the product $\mathrm{T}^{*} \mathrm{x}$ consists of the first m elements of the vector ifft (lambda. ${ }^{*} \mathrm{fft}([\mathrm{x}, \mathrm{z}])$ ), where z is a vector of zeros, while the product $\mathrm{T}^{* *} \mathrm{x}$ consists of the first n elements of the vector ifft(conj(lambda).*fft([x,z])).

## See Also

tprodinit - Initialization routine for tprod

## References

[1] P.C. Hansen, "Decconvolution and regularization with Toeplitz matrices," Numer. Algo. 29 (2002), pp. 323-378.

## tprodinit

## Purpose

Initialization routine for tprod (Toeplitz matrix-vector product)
Synopsis
lambda $=$ tprodinit(colT,rowT)

## Description

lambda contains the eigenvalues of a related circulant matrix, needed for matrix-vector multiplication with the Toeplitz matrix specified by colT and rowT. The length of lmbda is a powere of 2 .

## Input Parameters

colT contains the first column of the Toeplitz matrix; row T contains the first row of the Toeplitz matrix;

## Output Parameters

lambda contains the eigenvalues of an extended circulant matrix;

## Algorithm

lambda contains the eigenvalues of the circulant matrix C whose first column is $\mathrm{c}=$ [colT;z;rowT(end:-1:2)], where z is a zero columns with dimensions such that length(c) is a power of 2 (for efficiency).

See Also
tprod - Toeplitz matrix-vector product, using tprodinit

## References

[1] P.C. Hansen, "Decconvolution and regularization with Toeplitz matrices," Numer. Algo. 29 (2002), pp. 323-378.

## tvSv

## Purpose

Solves a rank-deficient system using the VSV decomposition.

## Synopsis

$x_{\text {_tvs }}=\operatorname{tvsv}(\mathrm{L}, \mathrm{V}, \mathrm{p}, \mathrm{b})$
x_tvsv $=\operatorname{tvsv}(\mathrm{R}$, Omega, $\mathrm{V}, \mathrm{p}, \mathrm{b})$
$x_{-}$tvsv $=\operatorname{tvsv}(\mathrm{L}$, Omega, $\mathrm{V}, \mathrm{p}, \mathrm{b})$

## Description

Solves the symmetric and near-rank deficient system of equations $\mathrm{A} x=\mathrm{b}$, using the rank-revealing VSV decomposition of A. Three decompositions of A can be used:

```
A = V *L'*L *V' (semidefinite A, lower triangular L)
A = V*R'*Omega*R*V' (indefinite A, upper triangular R)
A = V*L'*Omega*L*V' (indefinite A, lower triangular L)
```


## Input Parameters

Semindefinite case:
L lower triangular matrix;
V orthognal matrix;
p numerical rank;
b right-hand side;
Indefinite case:
L or $\mathrm{R} \quad$ lower or upper triangular matrix;
Omega signature matrix
V orthognal matrix;
p numerical rank;
b right-hand side;

## Output Parameters

x_tvsv truncated VSV solution;

## References

[1] P. C. Hansen and P. Y. Yalamov, "Computing Symmetric Rank-Revealing Decompositions via Triangular Factorization," SIAM J. Matrix Anal. Appl., 23 (2001), pp. 443-458.

## ulliv

## Purpose

High-rank-revealing ULLV algorithm, B full row rank.

## Synopsis

$[p, L A, L, V, U A, U B, v e c]=u l l i v(A, B)$
[p,LA,L,V,UA, UB,vec] =ulliv(A,B,tol_rank)
$[p, L A, L, V, U A, U B, v e c]=u l l i v\left(A, B, t o l \_r a n k, t o l \_r e f, m a x \_r e f\right)$
[p,LA,L,V,UA,UB,vec] = ulliv(A,B,tol_rank,tol_ref,max_ref,fixed_rank)

## Description

Computes a rank-revealing ULLIV decomposition of an mA -by-n matrix A with $\mathrm{mA} \geq$ n , and an mB -by-n matrix B with full row rank $\mathrm{mB} \leq \mathrm{n}$ :

```
A = UA*LA*[[ L 0 ] * V, and B = UB*[ L O l | *V,
    [ O I ]
```

Here, LA is n-by-n and $L$ is $m B$-by-mB and both are lower triangular; UA, UB and $V$ are unitary matrices, where only the first nA column of UA are computed.

The ULLV decomposition is a quotient ULV decomposition of the quotient $\mathrm{A} * \operatorname{pinv}(\mathrm{~B}) \_\mathrm{A}$, where $\operatorname{pinv}(\mathrm{B})_{-} \mathrm{A}$ is the A -weighted pseudoinverse of B :

```
A*pinv(B)_A = UA(:,1:mB)*LA(1:mB,1:mB)*UB'.
```

Hence the lower triangular matrix $\mathrm{LA}(1: \mathrm{mB}, 1: \mathrm{mB})$ reveals the numerical rank p of the matrix quotient.
Note that the algorithm is optimized for numerical rank p close to nB , and that this algorithm should not be used if $B$ is ill conditioned or rank deficient. Use the function ullv if B has full column rank.

```
Input Parameters
    A mA-by-n matrix ( }\textrm{mA}\geq\textrm{n}\mathrm{ );
    B mB-by-n matrix (mB \leqn);
    tol_rank rank decision tolerance;
    tol_ref upper bound on the 2-norm of the off-diagonal block
        LA(p+1:n,1:p) relative to the Frobenius-norm of LA;
    max_ref max. number of refinement steps per singular value
        to achieve the upper bound tol_ref;
    fixed_rank deflate to the fixed rank given by fixed_rank instead
        of using the rank decision tolerance;
    Defaults tol_rank = sqrt(n)*norm(A,1)*eps;
        tol_ref = 1e-04;
        max_ref = 0;
```


## Output Parameters

$\mathrm{p} \quad$ numerical rank of $\mathrm{A} * \operatorname{pinv}(\mathrm{~B})$ _A;
LA,L,V,UA,UB the ULLV factors;
vec a 5-by-1 vector with:
$\operatorname{vec}(1)=$ upper bound of $\operatorname{norm}(\mathrm{LA}(p+1: n, 1: p))$,
$\operatorname{vec}(2)=$ estimate of pth singular value,
$\operatorname{vec}(3)=$ estimate of $(p+1)$ th singular value,
$\operatorname{vec}(4)=$ a posteriori upper bound of num. nullspace angle,
$\operatorname{vec}(5)=$ a posteriori upper bound of num. range angle.

## Algorithm

First compute the LQ factorization $\mathrm{B}=[\mathrm{L}, 0] * \mathrm{~V}^{\prime}$ and then form the matrix $\mathrm{X}=$ $\mathrm{A} * \mathrm{~V} * \operatorname{inv}(\operatorname{diag}(\mathrm{~L}, \mathrm{I}))$, followed by the QR factorization $\mathrm{X}=\mathrm{UA} * \mathrm{LA}$. Thus, $\mathrm{A}=$ $\mathrm{UA} * \mathrm{LA} * \operatorname{diag}(\mathrm{~L}, \mathrm{I}) * \mathrm{~V}^{\prime}$ and $\mathrm{B}=\mathrm{L} * \mathrm{~V}^{\prime}$. Then deflation and refinement (optional) are employed to produce a rank-revealing decomposition. The deflation procedure is based on the generalized LINPACK condition estimator, and the refinement steps on QRiterations.

## See Also

ullv $\quad-$ Rank-revealing ULLV algorithm, B full column rank.

## References

[1] F.T. Luk and S. Qiao, "A New Matrix Decomposition for Signal Processing", Automatica, 30 (1994), pp. 39-43.
[2] F.T. Luk and S. Qiao, "An adaptive algoithm for interference cancelling in array processing; in F.T. Luk (Ed.), "Advanced Signal Processing Algorithms, Architectures, and Implementations VI," SPIE Proceedings, Vol. 2846 (1996), pp. 151-161.

## ulliv_up_a

## Purpose

Update the A-part of the rank-revealing ULLIV decomposition.

## Synopsis

[p,LA,L,V,UA,UB,vec] = ulliv_up_A(p,LA,L,V,UA,UB,a)
[p,LA,L,V,UA,UB,vec] = ulliv_up_A(p,LA,L,V,UA,UB,a,tol_rank)
[p,LA,L,V,UA,UB,vec] = ulliv_up_A(p,LA,L,V,UA,UB,a,tol_rank,tol_ref,max_ref)
$[p, L A, L, V, U A, U B, v e c]=$ ulliv_up_A(p,LA,L,V,UA,UB,a,tol_rank,tol_ref,max_ref,fixed_rank)

## Description

Given a rank-revealing ULLIV decomposition of the mA-by-n matrix with $\mathrm{mA} \geq \mathrm{n}$, and the full-rank mB-by-n matrix $\mathrm{B}=\mathrm{UB} * \mathrm{~L} * \mathrm{~V}^{\prime}$ with $\mathrm{mB}<\mathrm{n}$, the function computes the updated decomposition


```
[ a ] [ O I ]
```

where a is the new row added to A. Note that B must have full row rank, that the row dimension of UA will increase by one, and that the matrices UA and UB can be left out by inserting an empty matrix [] while V is required.

## Input Parameters

$\mathrm{p} \quad$ numerical rank of $\mathrm{A} * \operatorname{pinv}(\mathrm{~B})$ _A revealed in LA;
LA,L,V,UA,UB the ULLIV factors of $A$ and B;

| a | the new row added to A ; |
| :---: | :---: |
| tol_rank | decision tolerance; |
| tol_ref | upper bound on the 2-norm of the off-diagonal block |
| max_ref | LA( $\mathrm{p}+1: \mathrm{mB}, 1: \mathrm{p}$ ) relative to the Frobenius-norm of LA; max. number of refinement steps per singular value to achieve the upper bound tol ref; |
| fixed_rank | if true, deflate to the fixed rank given by p instead of using the rank decision tolerance; |
| Defaults | $\begin{aligned} & \text { tol_rank }=\operatorname{sqrt}(\mathrm{n}) * \text { norm }(\mathrm{LA}, 1) * \mathrm{eps} ; \\ & \text { tol_ref }=1 \mathrm{e}-04 ; \\ & \text { max_ref }=0 ; \end{aligned}$ |

## Output Parameters

$\mathrm{p} \quad$ numerical rank of updated quotient;
LA,L,V,UA,UB the updated ULLV factors;
vec a 5-by-1 vector with:
$\operatorname{vec}(1)=$ upper bound of $\operatorname{norm}(\mathrm{LA}(\mathrm{p}+1: \mathrm{mB}, 1: \mathrm{p}))$,
$\operatorname{vec}(2)=$ estimate of pth singular value,
$\operatorname{vec}(3)=$ estimate of $(p+1)$ th singular value,
$\operatorname{vec}(4)=$ a posteriori upper bound of num. nullspace angle,
$\operatorname{vec}(5)=$ a posteriori upper bound of num. range angle.

## See Also

ulliv_up_B - Update the B-part of the rank-revealing ULLIV decomp.

## References

[1] F.T.Luk and S. Qiao, "A New Matrix Decomposition for Signal Processing", Automatica, 30 (1994), pp. 39-43.
[2] F.T.Luk and S. Qiao, "An adaptive algoithm for interference cancelling in array processing; in F.T. Luk (Ed.), "Advanced Signal Processing Algorithms, Architectures, and Implementations VI," SPIE Proceedings, Vol. 2846 (1996), pp. 151-161.

## ulliv_up_B

## Purpose

Update the B-part of the rank-revealing ULLIV decomposition.

## Synopsis

$[p, L A, L, V, U A, U B, v e c]=$ ulliv_up_B(p,LA,L,V,UA,UB,b)
$[p, L A, L, V, U A, U B, v e c]=$ ulliv_up_B(p,LA,L,V,UA,UB,b,tol_rank)
[p,LA,L,V,UA,UB,vec] = ulliv_up_B(p,LA,L,V,UA,UB,b,tol_rank,tol_ref,max_ref)
[p,LA,L,V,UA,UB,vec] = ulliv_up_B(p,LA,L,V,UA,UB,b,tol_rank,tol_ref,max_ref,fixed_rank)

## Description

Given a rank-revealing ULLIV decomposition of the mA -by-n matrix with $\mathrm{mA} \geq \mathrm{n}$, and the full-rank mB-by-n matrix $\mathrm{B}=\mathrm{UB} * \mathrm{~L} * \mathrm{~V}^{\prime}$ with $\mathrm{mB}<\mathrm{n}$, the function computes the updated decomposition


```
    [ 0 I ] [ b ]
```

where $b$ is the new row added to $B$. Note that the updated matrix [B;b] must have full row rank, that the row dimension of UB will increase by one, and that the matrices UA and UB can be left out by inserting an empty matrix [] while V is required.

## Input Parameters

$p \quad$ numerical rank of $A * \operatorname{pinv}(B)$ _A revealed in LA;
$L A, L, V, U A, U B$ the ULLIV factors of $A$ and $B$;

| b | the new row added to B; |
| :--- | :--- |
| tol_rank | rank decision tolerance; <br> upper bound on the 2-norm of the off-diagonal block |
| tol_ref | LA $(\mathrm{p}+1: \mathrm{mB}, 1: \mathrm{p})$ relative to the Frobenius-norm of LA; <br> max. number of refinement steps per singular value <br> to achieve the upper bound tol_ref; |
| maxed_rank | if true, deflate to the fixed rank given by p <br> instead of using the rank decision tolerance; |
| Defaults | tol_rank $=\operatorname{sqrt}(\mathrm{n}) *$ norm $(\mathrm{LA}, 1) * e p s ;$ <br> tol_ref $=1 \mathrm{e}-04 ;$ <br> max_ref $=0 ;$ |

## Output Parameters

$\mathrm{p} \quad$ numerical rank of updated quotient;
LA,L,V,UA,UB the updated ULLV factors;
vec a 5-by-1 vector with:
$\operatorname{vec}(1)=$ upper bound of $\operatorname{norm}(\mathrm{LA}(\mathrm{p}+1: \mathrm{mB}, 1: \mathrm{p}))$,
$\operatorname{vec}(2)=$ estimate of pth singular value,
$\operatorname{vec}(3)=$ estimate of $(p+1)$ th singular value,
$\operatorname{vec}(4)=$ a posteriori upper bound of num. nullspace angle,
$\operatorname{vec}(5)=$ a posteriori upper bound of num. range angle.

## See Also

ulliv_up_A - Update the A-part of the rank-revealing ULLIV decomp.

## References

[1] F.T.Luk and S. Qiao, "A New Matrix Decomposition for Signal Processing", Automatica, 30 (1994), pp. 39-43.
[2] F.T.Luk and S. Qiao, "An adaptive algoithm for interference cancelling in array processing; in F.T. Luk (Ed.), "Advanced Signal Processing Algorithms, Architectures, and Implementations VI," SPIE Proceedings, Vol. 2846 (1996), pp. 151-161.

## vsv_qrit

## Purpose

Refinement of VSV decomposition via block QR-iterations.

## Synopsis

[L] = vsv_qrit(p,num_ref,L)
$[\mathrm{L}, \mathrm{V}]=$ vsv_qrit(p,num_ref,L, [],V)
[L,Omega] = vsv_qrit(p,num_ref,L,Omega)
[L,Omega,V] = vsv_qrit(p,num_ref,L,Omega,V)
[R,Omega] = vsv_qrit(p,num_ref,R,Omega)
$[R, O m e g a, V]=$ vsv_qrit(p,num_ref,R,Omega,V)

## Description

Given a VSV decomposition with numerical rank p, of one of the forms

```
A = V*L'*L *V' (semidefinite A, lower triangular L)
A = V*L'*Omega*L*V' (indefinite A, lower triangular L)
A = V*R'*Omega*R*V' (indefinite A, upper triangular R)
```

the function refines the rank-revealing decomposition via num_ref steps of block QR iterations applied to the triangular matrix.

## Input Parameters

```
p numerical rank of A;
```

num_ref number of refinement iterations;
T triangur matrix ( L or R , depending on VSV decomposition);
Omega signature matrix (indef. case) or empty (semidef. case);
$\checkmark \quad$ orthogonal matrix;

## Output Parameters

$\mathrm{T} \quad$ refined triangular matrix
Omega or V refined Omega (indef. case) or refined V (semidef. case)
V refined V (indef. case)

## Algorithm

Refinement is identical to block QR iteration, in which the off-diagonal block of the triangular matrix is "flipped" to the diagonally opposite position and then back again.

## See Also

ulv_ref - Refine one column of L in the ULV decomposition.
urv_ref - Refine one column of $R$ in the URV decomposition.

## References

[1] R. Mathias and G.W. Stewart, "A Block QR Algorithm and the Singular Value Decomposition", Lin. Alg. Appl., 182 (1993), pp. 91-100.

## vsvid_L_mod

## Purpose

Rank-one modification of VSV decomp. of sym. indef. matrix, L version.

## Synopsis

[p,L,Omega,V] = vsvid_L_mod(p,L,Omega,V,omega,v)
[p,L,Omega,V] = vsvid_L_mod(p,L,Omega,V,omega,v,tol_rank)
[p,L,Omega,V] = vsvid_L_mod(p,L,Omega,V,omega,v,tol_rank,inv_iter)
[p,L,Omega,V] = vsvid_L_mod(p,L,Omega,V,omega,v,tol_rank,inv_iter,fixed_rank)

## Description

Given a rank-revealing VSV decomposition of a symmetric indefinite matrix $\mathrm{A}=$ $\mathrm{V} *\left(\mathrm{~L}^{\prime} *\right.$ Omega $\left.* \mathrm{~L}\right) * \mathrm{~V}^{\prime}$, the function computes the updated rank-revealing decomposition of the matrix A + omega*v*v', where omega $=+1$ or -1 .

## Input Parameters

| p | numerical rank of A; |
| :---: | :---: |
| L | lower triangular matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{~L}^{\prime} * \mathrm{Omega} * \mathrm{~L}\right) * \mathrm{~V}^{\prime}$; |
| Omega | signature matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{~L}^{\prime} *\right.$ Omega $\left.* \mathrm{~L}\right) * \mathrm{~V}^{\prime}$; |
| V omega | orthogonal matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{~L}^{\prime} *\right.$ Omega $\left.* \mathrm{~L}\right) * \mathrm{~V}^{\prime}$; update $(+1)$ or downdate $(-1)$; |
| v | rank-one update column vector; |
| tol_rank | rank decision tolerance; |
| inv_iter | number of inverse iterations per deflation step; |
| fixed_rank | deflate to the fixed rank given by fixed_rank instead of using the rank decision tolerance; |
| Defaults | tol_rank $=\mathrm{n} *$ norm $(\mathrm{L}, 1) *$ eps; |
|  | inv_iter $=5$; |

## Output Parameters

p numerical rank of modified matrix;

L updated lower triangular matrix;
Omega updated signature matrix;
$V$ updated orthogonal matrix;
See Also
vsvid_R_mod - Rank-one mod. of VSV decomp. of sym. indef. matrix, $R$ version.
vsvsd_up - Rank-one update of VSV decomposition of semidefinite matrix.

## References

[1] P.C. Hansen \& P.Y. Yalamov, "Computing Symmetric Rank-Revealing Decompositions via Triangular Factorization", SIAM J. Matrix Anal. Appl., 23 (2001), pp. 443-458.

## vsvid_R_mod

## Purpose

Rank-one modification of VSV decomp. of indef. matrix, R version.

## Synopsis

[p,R,Omega, V$]=$ vsvid_R_mod(p,R,Omega,V,omega,v)
$[p, R, O m e g a, V]=$ vsvid_R_mod(p,R,Omega,V,omega,v,tol_rank)
$[p, R, O m e g a, V]=$ vsvid_R_mod(p,R,Omega,V,omega,v,tol_rank,inv_iter)
$[p, R, O m e g a, V]=$ vsvid_R_mod(p,R,Omega,V,omega,v,tol_rank,inv_iter,fixed_rank)

## Description

Given a rank-revealing VSV decomposition of a symmetric indefinite matrix $\mathrm{A}=$ $\mathrm{V} *\left(\mathrm{R}{ }^{\prime} *\right.$ Omega $\left.* \mathrm{R}\right) * \mathrm{~V}^{\prime}$, the function computes the updated rank-revealing decomposition of the matrix $\mathrm{A}+$ omega $^{\mathrm{v}} * \mathrm{v}^{\prime}$, where omega $=+1$ or -1 .

## Input Parameters

| p | numerical rank of A; |
| :---: | :---: |
| R | upper triangular matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{R}^{\prime} *\right.$ Omega $\left.* \mathrm{R}\right) * \mathrm{~V}^{\prime}$; |
| Omega | signature matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{R}^{\prime} *\right.$ Omega $\left.* \mathrm{R}\right) * \mathrm{~V}^{\prime}$; |
| V | orthogonal matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{R}^{\prime} * \mathrm{Omega} * \mathrm{R}\right) * \mathrm{~V}^{\prime}$; |
| omega | update ( +1 ) or downdate ( -1 ); |
| $\checkmark$ | rank-one update column vector; |
| tol_rank | rank decision tolerance; |
| inv_iter | number of inverse iterations per deflation step; |
| fixed_rank | deflate to the fixed rank given by fixed_rank instead of using the rank decision tolerance; |
| Defaults | tol_rank $=\mathrm{n} *$ norm $(\mathrm{R}, 1) * \mathrm{eps}$; |
|  | $\text { inv_iter }=5$ |
| ut Param |  |
| p | numerical rank of updated A; |
| R | updated upper triangular; |
| Omega | updated signature matrix; |
| V | updated orthogonal matrix; |

See Also
vsvid_mod - Rank-one modification of VSV decomposition of indefinite matrix.
vsvsd_up - Rank-one update of VSV decomposition of semidefinite matrix.

## References

[1] P.C. Hansen \& P.Y. Yalamov, "Computing Symmetric Rank-Revealing Decompositions via Triangular Factorization", SIAM J. Matrix Anal. Appl., 23 (2001), pp. 443-458.

## vsvid_ip

## Purpose

Interim process for indefinite VSV algorithm.

## Synopsis

[W,C,Omega] = vsvid_ip(L,D)

## Input Parameters

L,D
Factors in $\mathrm{LDL}^{T}$ factorization;

## Output Parameters

W,C,Omega Matrices in $\mathrm{L} * \mathrm{D} * \mathrm{~L}^{\prime}=\mathrm{W} * \mathrm{C}^{\prime} *$ Omega $* \mathrm{C} * \mathrm{~W}^{\prime}$.

## Algorithm

The factorization $\mathrm{L} * \mathrm{D} * \mathrm{~L}^{\prime}$ is replaced with $\mathrm{W} * \mathrm{C}^{\prime} * \mathrm{Omega}^{2} * \mathrm{C} * \mathrm{~W}^{\prime}$, where W is orthogonal, C is upper triangular, and Omega is a signature matrix. This is accomplished via small eigenvalue decompositions of the 1-by-1 and 2-by- 2 blocks of $D$.

## References

[1] P.C. Hansen, \& P.Y. Yalamov, "Computing Symmetric Rank-Revealing Decompositions via Triangular Factorization", SIAM J. Matrix Anal. Appl., 23 (2001), pp. 443-458.

## vsvsd_up

## Purpose

Rank-one update of VSV decomposition of semidefinite matrix.

## Synopsis

$[p, L, V]=$ vsvsd_up $(p, L, V, v)$
$[p, L, V]=$ vsvsd_up( $p, L, V, v$, tol_rank)
$[p, L, V]=$ vsvsd_up $(p, R, V, v$, tol_rank,fixed_rank)

## Description

Given a rank-revealing VSV decomposition of a symmetric semidefinite matrix $\mathrm{A}=$ $\mathrm{V} *\left(\mathrm{~L}^{\prime} * \mathrm{~L}\right) * \mathrm{~V}^{\prime}$, the function computes the updated rank-revealing decomposition of $\mathrm{A}+$ $\mathrm{v} * \mathrm{v}^{\prime}$. Use function vsvid_mod with omega $=-1$ to downdate the VSV decomposition of a symmetric semidefinite matrix.

## Input Parameters

$p \quad$ numerical rank of A;
$\mathrm{L} \quad$ lower triangular matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{~L}{ }^{\prime} * \mathrm{~L}\right) * \mathrm{~V}^{\prime}$;
$\mathrm{V} \quad$ orthogonal matrix in $\mathrm{A}=\mathrm{V} *\left(\mathrm{~L}^{\prime} * \mathrm{~L}\right) * \mathrm{~V}^{\prime}$;
v rank-one update column vector;
tol_rank rank decision tolerance;
fixed_rank deflate to the fixed rank given by fixed_rank instead of using the rank decision tolerance;

Defaults tol_rank $=\mathrm{n} *$ norm $(\mathrm{L}, 1) * e \mathrm{ps} ;$

## Output Parameters

$\mathrm{p} \quad$ numerical rank of updated A ;
L updated lower triangular matrix;
V updated orthogonal matrix;
See Also
vsvid_L_mod - Rank-one mod. of VSV decomp. of sy,. indef. matrix, L version. vsvid_R_mod - Rank-one mod. of VSV decomp. of sy,. indef. matrix, R version.

## References

[1] P.C. Hansen \& P.Y. Yalamov, "Computing Symmetric Rank-Revealing Decompositions via Triangular Factorization", SIAM J. Matrix Anal. Appl., 23 (2001), pp. 443-458.

