## Empty convex polytopes in random point sets

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#### Abstract

Given a set P of points in  $\mathbb{R}^d$ , a convex hole (alternatively, empty convex polytope) of P is a convex polytope with vertices in P, containing no points of P in its interior. Let R be a bounded convex region in  $\mathbb{R}^d$ . We show that if P is a set of n random points chosen independently and uniformly over R, then the expected number of vertices of the largest hole of P is  $\Theta(\log n/(\log \log n))$ , regardless of the shape of R. This generalizes the analogous result proved for the case d=2 by Balogh, González-Aguilar, and Salazar.

#### Introduction

Given a set P of points in  $\mathbb{R}^d$ , a convex hole (alternatively, empty convex polytope) of P is a convex polytope with vertices in P, containing no points of P in its interior.

Recently, we showed that the expected size of the largest convex hole in a random n-point set in the plane is  $\Theta(\log n/\log\log n)$  [3]. One anonymous referee of this paper asked if this could be generalized to d>2 dimensions. Joe O'Rourke asked the same question in MathOverflow, and Douglas Zare replied that the  $\Omega(\log n/\log\log n)$  lower bound carries over easily to the d-dimensional case [10]. At the end of his reply, Zare wrote: "I don't know whether their harder upper bound of the same form also extends to higher dimensions, but I suspect that it does."

Our aim in this note is to show that, indeed, the upper bound also holds for higher dimensions. Thus, our main result is:

**Theorem 1** Let  $d \geq 2$  be an integer, and let R be a bounded convex region in  $\mathbb{R}^d$ . Let  $R_n$  be a set of n points chosen independently and uniformly at random from R, and let  $HOL(R_n)$  denote the random variable

that measures the number of vertices of the largest convex hole in  $R_n$ . Then

$$\mathbf{E}(\operatorname{Hol}(R_n)) = \Theta\left(\frac{\log n}{\log \log n}\right).$$

Moreover, a.a.s.

$$\operatorname{Hol}(R_n) = \Theta\left(\frac{\log n}{\log\log n}\right).$$

The proof, which is an immediate consequence of Theorems 2 and 3 below, follows very closely the main ideas of the proof of [3, Theorem 3]. Indeed, the strategy and the main ideas are so close that it seems best to follow as closely as possible the structure of [3]. As we shall see below, some of the results proved in [3] follow without any modification to arbitrary dimensions. The main adaptations needed are:

- 1. a generalization of the results in [3, Section 2] to d > 2 dimensions, to approximate convex sets in  $\mathbb{R}^d$  with lattice polytopes; and
- 2. an adaptation to d > 2 dimensions of the results on the probability that a random n-point set is in convex position, from the exact results of Valtr [11, 12] in  $\mathbb{R}^2$  to the asymptotic results of Bárány [6] in  $\mathbb{R}^d$ , for any  $d \geq 2$ .

The workhorse for the proof of Theorem 1 for arbitrary regions R is the following statement, which takes care of the particular case in which R is a parallelotope.

**Theorem 2** Let R be a parallelotope in  $\mathbb{R}^d$ . Let  $R_n$  be a set of n points chosen independently and uniformly at random from R, and let  $HoL(R_n)$  denote the random variable that measures the number of vertices of the largest convex hole in  $R_n$ . Then

$$\mathbf{E}(\operatorname{Hol}(R_n)) = \Theta\left(\frac{\log n}{\log \log n}\right).$$

Moreover, a.a.s.

$$\operatorname{Hol}(R_n) = \Theta\left(\frac{\log n}{\log\log n}\right).$$

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The other essential fact is that the order of magnitude of the expected number of vertices of the largest convex hole is independent of the shape of R:

**Theorem 3** There exist absolute constants b, b' with the following property. Let R and S be bounded convex regions in  $\mathbb{R}^d$ . Let  $R_n$  (respectively,  $S_n$ ) be a set of n points chosen independently and uniformly at random from R (respectively, S). Let  $Hol(R_n)$  (respectively,  $Hol(S_n)$ ) denote the random variable that measures the number of vertices of the largest convex hole in  $R_n$  (respectively,  $S_n$ ). Then, for all sufficiently large n,

$$b \le \frac{\mathbf{E}(\mathrm{HOL}(R_n))}{\mathbf{E}(\mathrm{HOL}(S_n))} \le b'.$$

Moreover, there exist absolute constants c, c' such that a.a.s.

$$c \le \frac{\operatorname{Hol}(R_n)}{\operatorname{Hol}(S_n)} \le c'.$$

We remark that Theorem 3 is in line with the following result proved by Bárány and Füredi [5]: the expected number of empty simplices in a set of n points chosen uniformly and independently at random from a convex set A with non-empty interior in  $\mathbb{R}^d$  is  $\Theta(n^d)$ , regardless of the shape of A.

Remark (Proof of Theorem 1). Theorem 1 is an immediate consequence of Theorems 2 and 3.

The proof of Theorem 2 is in Section 1. As we explain in Section 2, the proof of Theorem 3 is totally analogous to the proof of [3, Theorem 2].

We make a few final remarks before we move on to the proofs. For the rest of the paper we let Vol(U)denote the volume of a region U in  $\mathbb{R}^d$ . We also note that, throughout the paper, by  $\log x$  we mean the natural logarithm of x. Finally, since we only consider sets of points chosen independently and uniformly at random from a region, for brevity we simply say that such point sets are chosen at random from this region.

### 1 Proof of Theorem 2

We start by noting that if Q, Q' are two regions such that Q' is obtained from Q by an affine transformation, then  $HoL(Q_n) = HoL(Q'_n)$ . Thus we may assume without loss of generality that R is the isothetic unit area square centered at the origin.

We prove the lower and upper bounds separately. More specifically, we prove that for all sufficiently large n:

$$\Pr\left(\operatorname{HoL}(R_n) \ge \frac{1}{2} \frac{\log n}{\log \log n}\right) \ge 1 - n^{-2}.$$
 (1)

$$\Pr\left(\text{Hol}(R_n) \le d(2 + 2d^2) \frac{\log n}{\log \log n}\right) \ge 1 - n^{-1}.$$
 (2)

We note that (1) and (2) imply immediately the a.a.s. part of Theorem 2. Now the  $\Omega(\log n/\log\log n)$  part of the theorem follows from (1), since  $\operatorname{HoL}(R_n)$  is a non-negative random variable, whereas the  $O(\log n/\log\log n)$  part follows from (2), since  $\operatorname{HoL}(R_n)$  is bounded by above by n.

Thus we complete the proof by showing (1) and (2).

Proof of (1)

Let  $R_n$  be a set of n points chosen at random from R. We prove that a.a.s.  $R_n$  has an empty convex polytope of size at least  $\frac{\log n}{2\log\log n}$ . Consider the 2-dimensional projection  $\pi$ :

Consider the 2-dimensional projection  $\pi$ :  $\mathbb{R}^d \to \mathbb{R}^2$  defined by  $(x_1, x_2, x_3, x_4, \ldots, x_d) \to (x_1, x_2, 0, 0, \ldots, 0)$ . Note that  $\pi(R_n)$  is a set of n points chosen (independently and uniformly) at random from the unit square. Thus it follows from Eq. (1) in [3] that a.a.s.  $\pi(R_n)$  has a convex hole H of size at least  $\frac{\log n}{2\log\log n}$ . Clearly,  $\pi^{-1}(H)$  is an empty convex polytope of  $R_n$  of size at least  $\frac{\log n}{2\log\log n}$ .

Proof of (2)

Let  $R_n$  be a set of n points chosen at random from R. We remark that throughout the proof we always implicitly assume that n is sufficiently large.

We shall use the following easy consequence of Chernoff's bound. This is derived immediately, for instance, from Theorem A.1.11 in [1].

**Lemma 4** Let  $Y_1, \ldots, Y_m$  be mutually independent random variables with  $\Pr(Y_i = 1) = p$  and  $\Pr(Y_i = 0) = 1 - p$ , for  $i = 1, \ldots, m$ . Let  $Y := Y_1 + \ldots + Y_m$ .

$$\Pr(Y \ge (3/2)pm) < e^{-pm/16}.$$

Let S be the isothetic d-cube of volume  $3^d$ , also (as R) centered at the origin.

We need the following result on approximating convex sets by lattice parallelotopes.

Claim A. For each positive integer d > 0 there exist integers  $f_1(d)$  and  $f_2(d)$  with the following property. Let H be a convex set in  $\mathbb{R}^d$ . Then there exists a lattice parallelotope  $Q_1$  such that  $H \subseteq Q_1$  and  $\operatorname{Vol}(Q_1) \leq (f_1(d)+1)\operatorname{Vol}(H)$ . Moreover, if  $\operatorname{Vol}(H) \geq 2^{d-1} \cdot 1000/n$ , then there is a lattice parallelotope  $Q_0$  such that  $Q_0 \subseteq H$  and  $\operatorname{Vol}(Q_0) \geq (f_2(d)-1)\operatorname{Vol}(H)$ .

**Sketch of Proof.** By a the theorem of M. Balla [2], for every convex compact set  $H \subset \mathbb{R}^d$  there exists a parallelotope P such that  $P \subset H \subset dP = \widehat{P}$  where dP is the image of P under a homothety with ratio d. This implies that  $d^{-d}\operatorname{Vol}(P) \leq \operatorname{Vol}(H) \leq d^d\operatorname{Vol}(P)$ .

For each vertex  $v_i$ ,  $i = 0, ..., 2^d$ , of  $\widehat{P}$ , let us denote by  $Q_{v_i}$  the parallelotope with side length 2/n with

facets parallel to the facets of  $\widehat{P}$  that has  $v_i$  as one of its vertices and  $\widehat{P} \cap Q_{v_i} = \{v_i\}$ . Observe that each  $Q_{v_i}$  contains a d-ball of diameter 2/n and for that, there is a lattice point  $v_i'$  in the interior of each  $Q_{v_i}$ . Let  $Q_1$  be the convex hull of the points  $v_1', \ldots, v_{2^d}'$ . Note that  $d(v_i, v_i') \leq \frac{2}{n} \sqrt{d}$  for each  $i = 1, \ldots, 2^d$ , this implies that  $\varrho(\widehat{P}, Q_1) \leq \frac{2}{n} \sqrt{d}$  where  $\varrho(\cdot, \cdot)$  is the Hausdorff metric. Then

$$Vol(Q_1) \leq Vol(\widehat{P}) + \frac{2}{n}\sqrt{d} \cdot Surf(\widehat{P}) + Vol(B)$$
  
$$\leq d^d Vol(H) + \frac{2}{n}\sqrt{d} \cdot Surf(\widehat{P}) + Vol(B)$$
  
$$\leq (f_1(d) + 1) Vol(H)$$

where B is the d-ball of diameter  $\frac{2}{n}\sqrt{d}$  and  $\mathrm{Surf}(\cdot)$  is the volume (d-1)-dimensional.

Now, for each vertex  $w_i$ ,  $i=0,\ldots,2^d$ , of P consider the parallelotope  $Q_{w_i}$  with side length 2/n with facets parallel to the facets of P that has  $w_i$  as one of its vertices and  $Q_{w_i} \subset P$ . Because each  $Q_{w_i}$  contains a d-ball of diameter  $\frac{2}{n}$ , then there exists a lattice point in each  $Q_{w_i}$ , let  $w_i'$  this point. The existence of these points is guaranteed provided that  $\operatorname{Vol}(H) \geq 2^{d-1} \cdot 1000/n$  Let  $Q_0$  be the convex hull of the points  $w_1',\ldots,w_{2^d}'$ . Note that  $d(w_i,w_i') \leq \frac{2}{n}\sqrt{d}$  for each  $i=1,\ldots,2^d$ , this implies that  $\varrho(P,Q_0) \leq \frac{2}{n}\sqrt{d}$ . Then  $\operatorname{Vol}(P) - \frac{2}{n}\sqrt{d} \cdot \operatorname{Surf}(Q_0) - \operatorname{Vol}(B) \leq \operatorname{Vol}(Q_0)$  This implies

$$Vol(Q_0) \geq Vol(P) - \frac{2}{n}\sqrt{d} \cdot Surf(Q_0) - Vol(B)$$

$$\geq Vol(P) - \frac{2}{n}\sqrt{d} \cdot Surf(P) - Vol(B)$$

$$\geq d^{-d} Vol(H) - \frac{2}{n}\sqrt{d} \cdot Surf(H) - Vol(B)$$

$$\geq (f_2(d) - 1)Vol(H).$$

For the rest of the proof, for simplicity we define  $f_3(d) := f_3(d)$ , where  $f_1(d)$  and  $f_2(d)$  are as in Claim A

Since there are  $(9n+1)^d < (10n)^d$  lattice points, out of which  $(n+1)^d$  are in R, it follows that there are fewer than  $\binom{(10n)^d}{2^d} < (10n)^{d2^d}$  lattice d-parallelotopes in total, and fewer than  $\binom{n^d}{2^d} < n^{d2^d}$  lattice d-parallelotopes all of whose vertices are in R.

Claim B. With probability at least  $1 - n^{-10}$  every lattice parallelotope Q with  $Vol(Q) < 20f_3(d) \log n/n$  satisfies that  $|R_n \cap Q| \le (3/2) \cdot 20f_3(d) \log n$ .

**Proof.** We note that, since  $(f_1(d) + 1)/f_2(d) > 1$ , it follows that  $20f_3(d) > d2^d + 10$ . Let Q be a lattice parallelotope with  $\operatorname{Vol}(Q) < 20 f_3(d) \log n/n$ , and let  $Z = Z(Q) \subseteq R$  be any lattice parallelotope containing Q, with  $\operatorname{Vol}(Z) = 20 f_3(d) \log n/n$ . Let  $X_Q$  (respectively,  $X_Z$ ) denote the random variable that

measures the number of points of  $R_n$  in Q (respectively, Z). We apply Lemma 4 with  $p = \operatorname{Vol}(Z)$  and m = n, to obtain  $\Pr(X_Z \geq (3/2) \cdot 20 \, f_3(d) \log n) < e^{-(3/2)20 \, f_3(d)/24 \log n} = n^{-(5/4)f_3(d)}$ . Since  $Q \subseteq Z$ , it follows that  $\Pr(X_Q \geq (3/2) \cdot 20 \, f_3(d) \log n) < n^{-(5/4)f_3(d)}$ . As the number of choices for Q is at most  $(10n)^{d2^d}$ , with probability at least  $(1-(10n)^{d2^d} \cdot n^{-(5/4)f_3(d)}) > 1-n^{-10}$ , no such Q contains more than  $(3/2) \cdot 20 \, f_3(d) \log n$  points of  $R_n$ .

A polytope is empty if its interior contains no points of  $R_n$ .

**Claim C.** With probability at least  $1 - n^{-10}$ , there is no empty lattice parallelotope  $Q \subseteq R$  with  $Vol(Q) \ge 20 (d2^d + 10) \log n/n$ .

**Proof.** The probability that a fixed lattice parallelotope  $Q \subseteq R$  with  $\operatorname{Vol}(Q) \ge d2^d + 10 \log n/n$  is empty is  $(1 - \operatorname{Vol}(Q))^n < n^{-d2^d + 10}$ . Since there are fewer than  $n^{d2^d}$  lattice parallelotopes in R, it follows that the probability that at least one of the lattice parallelotope with area at least  $(d2^d + 10) \log n/n$  (and hence with area at least  $20 (d2^d + 10) \log n/n$ ) is empty is less than  $n^{d2^d} \cdot n^{-(d2^d + 10)} \le n^{-10}$ .

For the rest of the proof, we let H be a maximum size convex hole of  $R_n$ .

Claim D. With probability at least  $1 - n^{-10}$  we have  $Vol(Q_1) < f_3(d) \log n/n$ .

**Proof.** Suppose first that  $\operatorname{Vol}(H) < 2^{d-1}1000/n$ . Then  $\operatorname{Vol}(Q_1) \le 2^{d-1} \cdot 1000(f_1(d)+1)/n$ . Since this is obviously smaller than  $f_3(d)\log n/n$ , in this case we are done. Now suppose that  $\operatorname{Vol}(H) \ge 2^{d-1} \cdot 1000/n$ , so that  $Q_0$  (from Claim A) exists. Moreover,  $\operatorname{Vol}(Q_1) \le (f_1(d)+1)\operatorname{Vol}(H)$ . Since  $Q_0 \subseteq H$ , and H is a hole of  $R_n$ , it follows that  $Q_0$  is empty. Thus, by Claim C, with probability at least  $1-n^{-10}$  we have that  $\operatorname{Vol}(Q_0) < (d2^d+10)\log n/n$ . Now since  $\operatorname{Vol}(Q_1) < (f_1(d)+1)\operatorname{Vol}(H)$  and  $\operatorname{Vol}(Q_0) \ge f_2(d) \cdot \operatorname{Vol}(H)$ , it follows that  $\operatorname{Vol}(Q_1) \le (f_1(d)+1)\operatorname{Vol}(Q_0)/f_2(d)$ . Thus with probability at least  $1-n^{-10}$  we have that  $\operatorname{Vol}(Q_1) \le ((f_1(d)+1)(d2^d+10)/f_2(d))\log n/n = f_3(d)\log n/n$ .

Claim E. For each fixed integer d > 0, there exist a universal positive constant  $c_2 := c_2(d)$  with the following property. Let K be any convex polytope in  $\mathbb{R}^d$ . Then the probability that r points chosen at random from K are in convex position is at most  $(c_2 n^{\frac{2}{d-1}})^{-n}$ .

**Proof.** This is an immediate consequence of [6] (see for instance [4, Theorem 2.1]).  $\Box$ 

Claim F. With probability at least  $1 - 2n^{-2}$  the random point set  $R_n$  satisfies that no lattice parallelotope Q with  $\operatorname{Vol}(Q) < 20 \, f_3(d) \log n/n$  contains  $d(2+2d^2) \log n/(\log \log n)$  points of  $R_n$  in convex position.

**Proof.** Let Q be a lattice parallelotope with  $\operatorname{Vol}(Q) < 20 \, f_3(d) \log n/n$ . By Claim B, with probability at least  $1 - n^{-10}$  we have  $|R_n \cap Q| \leq (3/2) \cdot 20 \, f_3(d) \log n$ . Thus it suffices to show that the probability that there exists a lattice parallelotope Q with with  $|R_n \cap Q| \leq (3/2) \cdot 20 \, f_3(d) \log n$  and  $d(2 + 2d^2) \log n/(\log \log n)$  points of  $R_n$  in convex position is at most  $n^{-2}$ .

Let  $c_2 := c_2(d)$  be as in Claim E. Thus the expected number of r-tuples of  $R_n$  in Q in convex position is at most

$$\begin{split} \binom{|R_n \cap Q|}{r} \left(c_2 r^{\frac{2}{d-1}}\right)^{-r} \\ & \leq \binom{(3/2) \cdot f_3(d) \log n}{r} \left(c_2 r^{\frac{2}{d-1}}\right)^{-r} \\ & \leq \left(\frac{e \cdot (3/2) \cdot F_2 \log n}{r}\right)^r \left(c_2 r^{\frac{2}{d-1}}\right)^{-r} \\ & < \left(c_3 \log n \cdot r^{-1 - \frac{2}{d-1}}\right)^r, \end{split}$$

where  $c_3 := 3ef_3(d)/2c_2$ .

Since there are at most  $n^{d2^d}$  choices for Q, it follows that the expected total number of such r-tuples with  $r = d(2+2d^2)\log n/\log\log n$  is at most  $n^{d2^d} \cdot \left(c_3\log n \cdot r^{-1-\frac{2}{d-1}}\right)^r < n^{-2}$  (this last inequality follows from an elemenary but long manipulation). This completes the proof, since it follows that the probability that such an r-tuple exists is at most  $n^2$ .

To finish the proof of (2), recall that H is a maximum size empty convex polytope of  $R_n$ , and that  $H \subseteq Q_1$ . It follows immediately from Claims D and F that with probability at least  $1 - n^{-1}$  the parallelotope  $Q_1$  does not contain a set of  $d(2 + 2d^2) \log n/(\log \log n)$  points of  $R_n$  in convex position. In particular, with probability at least  $1 - n^{-1}$  the size of H is at most  $d(2 + 2d^2) \log n/(\log \log n)$ .

# 2 Proof of Theorem 3

The proof of Theorem 3 is totally analogous to the proof of [3, Theorem 2]. Indeed, in that proof, essentially all the arguments are independent of the dimension. The only adaptation that needs to be done is that we need a version of [3, Corollary 6] for d > 2 dimensions. We recall that [3, Corollary 6]

claims that if H is a closed convex set in  $\mathbb{R}^2$ , then there exist rectangles U, K such that  $U \subseteq H \subseteq K$ ,  $Vol(U) \geq Vol(H)/8$ , and  $Vol(K) \leq 2Vol(H)$ .

A d-dimensional analogue of this statement follows from the following result in [9]: if H is a convex body in  $\mathbb{R}^d$ , then H contains a parallelotope P such that some translate of dP contains K. Indeed, this implies at once that if H is a closed convex set in  $\mathbb{R}^d$ , then there exist parallelotopes U, K such that  $U \subseteq H \subseteq K$ ,  $\operatorname{Vol}(U) \geq \operatorname{Vol}(H)/d$ , and  $\operatorname{Vol}(K) \leq d \cdot \operatorname{Vol}(H)$ .

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