

Three-Way Symbolic Tree-Maps and Ultrametrics

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Abstract: Three-way dissimilarities are a generalization of (two-way) dissimilarities which can be used to indicate the lack of homogeneity or resemblance between any three objects. Such maps have applications in cluster analysis and have been used in areas such as psychology and phylogenetics, where three-way data tables can arise. Special examples of such dissimilarities are three-way tree-metrics and ultrametrics, which arise from leaf-labelled trees with edges labelled by positive real numbers. Here we consider three-way maps which arise from leaf-labelled trees where instead the interior *vertices* are labelled by an *arbitrary* set of values. For unrooted trees, we call such maps *three-way symbolic tree-maps*; for rooted trees, we call them *three-way symbolic ultrametrics* since they can be considered as a generalization of the (two-way) symbolic ultrametrics of Böcker and Dress. We show that, as with two- and three-way tree-metrics and ultrametrics, three-way symbolic tree-maps and ultrametrics can be characterized via certain k -point conditions. In the unrooted case, our characterization is mathematically equivalent to one presented by Gurvich for a certain class of edge-labelled hypergraphs. We also show that it can be decided whether or not an arbitrary three-way symbolic map is a tree-map or a symbolic ultrametric using a triplet-based approach that relies on the so-called BUILD algorithm

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for deciding when a set of 3-leaved trees or *triplets* can be displayed by a single tree. We envisage that our results will be useful in developing new approaches and algorithms for understanding 3-way data, especially within the area of phylogenetics.

Keywords: Three-way dissimilarity; Three-way symbolic map; Symbolic ultrametric; Ultrametric; Tree-metric; Phylogenetic tree.

1. Introduction

Three-way dissimilarities are a generalization of (two-way) dissimilarities, which can be used to indicate the lack of homogeneity or resemblance between any three objects in a given set (Joly and Le Calvé, 1995). They have applications in areas such as psychology (Heiser and Bennani, 1997) and phylogenetics (Levy, Yoshida, and Pachter, 2006), where they have been used to cluster data presented in the form of three-way data tables. Various special classes of three-way dissimilarities have been introduced (see e.g. Chepoi and Fichet, 2007; Hayashi, 1972; Heiser and Bennani, 1997; Joly and Le Calvé, 1995). These include three-way dissimilarities that arise from leaf-labelled trees, where the edges are weighted by positive real numbers. These so-called *three-way tree-metrics* and *three-way ultrametrics*, which arise from unrooted and rooted trees, respectively, generalize their much studied two-way counterparts (cf. Chepoi and Fichet, 2007, for an overview).

Intriguingly, in Böcker and Dress (1998), Böcker and Dress showed that the concept of ultrametricity for dissimilarities can be naturally extended to include two-way symmetric maps whose range is an arbitrary set of symbols. In particular, they introduced the concept of a *symbolic ultrametric* (a two-way map arising from a rooted, vertex-labelled tree via the least common ancestor map), and characterized them in terms of a 3- and a 4-point condition (see Section 2 for full details), a result which had in fact been discovered independently in another guise by V. Gurvich (1984) (see Section 2 for details). These conditions generalize the well-known 3-point condition for ultrametricity (cf. e.g. Semple and Steel, 2003, Chapter 7.2). Symbolic ultrametrics have been found to have interesting connections with cograph theory (Hellmuth et al., 2013), game theory (Gurvich, 1984; Gurvich, 2009), as well as applications within phylogenetics (Hellmuth et al., 2015; Lafond and El-Mabrouk, 2015). Therefore, it is of interest to understand how the theory of symbolic ultrametrics can be extended to three-way maps, as these may lead to useful new applications in these areas (e.g. see the last section for a potential application in phylogenetics).

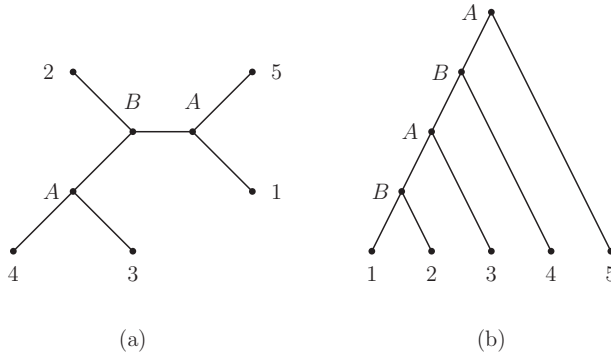


Figure 1. Two trees which give rise to (a) a three-way symbolic tree-map and (b) a three-way symbolic ultrametric.

In this paper, we shall address this question. Let X be a set (of taxa) of size at least 3 and let M be a set (of symbols) of size at least 2. A (*three-way*) *symbolic map* is a map $\delta : \binom{X}{3} \rightarrow M$. For example, consider the unrooted tree in Figure 1(a) with leaf-set $X = \{1, 2, 3, 4, 5\}$ whose interior vertices are labelled by elements from the set $M = \{A, B\}$. This tree gives rise naturally to a three-way symbolic map from X to the set M ; to each triple of leaves we assign the element of M which labels the vertex lying on all shortest paths between any two of these three leaves (e.g., the triple $\{1, 3, 5\}$ is assigned the symbol A). We call symbolic maps that arise in this way *three-way symbolic tree-maps*.

In Section 3, we show that a three-way symbolic tree-map uniquely determines its underlying labelled tree (Proposition 3.2), and also give a 4- and 5-point characterization for such maps (see Theorem 3.3). This result is mathematically equivalent to Gurvich (1984, Theorem 5), but for completeness we provide its proof. Our characterization for three-way symbolic tree-maps is analogous to the well-known 4-point condition for tree-metrics (cf. Semple and Steel, 2003, Chapter 7.1), and also generalizes the conditions presented in Herrmann, Huber, and Moulton (2012, Theorem 7) for determining when a three-way dissimilarity arises from a tree. To prove Theorem 3.3 we introduce a symbolic variant of the Farris transform (Semple and Steel, 2003, p. 149), which allows us to apply the main result from Böcker and Dress, 1998). We conclude the section with a description of how our result is related to the ones presented in Gurvich (1984).

In Section 4, we turn our attention to obtaining three-way symbolic maps from rooted trees. Consider the rooted tree in Figure 1(b). A symbolic ultrametric can be associated to this tree by defining the

value for each pair of leaves to be the symbol labelling the least common ancestor vertex of these two leaves. Therefore, a natural way to define a three-way symbolic ultrametric could be to take the value of each triple of leaves to be the *set* consisting of the symbols labelling the least common ancestor of all pairs of leaves in the triple (for example, we would assign the set $\{A, B\}$ to the triple $1, 2, 5$). However, this does not suffice to capture the tree (see Section 4).

Even so, as we shall see, if we consider the values of the triples to be *multisets* instead of sets (for example, we would assign the multiset $\{A, A, B\}$ to the triple $1, 2, 5$ in Figure 1(b)), then we can in fact recover the underlying labelled tree in case $|X| \geq 5$ (Theorem 4.4). We call maps obtained in this way *three-way symbolic ultrametrics*. In Section 5, we give 3-, 4- and 5-point conditions which ensure that a three-way symbolic map that maps into the set of size 3 multisets of a set of symbols is a symbolic ultrametric. This is somewhat surprising since for three-way dissimilarities, a 6-point condition is required to ensure that they can be represented by a rooted tree in an analogous way (cf. Herrmann et al., 2012, Theorem 7).

We conclude the paper by considering an alternative approach for deciding whether or not a three-way symbolic map is a tree-map or symbolic ultrametric. This approach is based on the BUILD algorithm (Aho et al., 1981), which can be used to decide when a set of *triplets* (i. e. resolved rooted leaf-labelled trees each with three leaves) is displayed by some supertree or not. Applying this algorithm to three-way symbolic maps has the advantage that only sets of size three (as opposed to sets of size up to five) need to be considered so as to determine if a three-way symbolic map is a tree-map or a symbolic ultrametric. This could potentially lead to practical algorithms for performing this task. In Section 7, we present some future directions.

2. Preliminaries

For a set $\{x_1, \dots, x_k\}$, $k \geq 1$, in the powerset $\mathcal{P}(X)$ of X and a map $\delta : \mathcal{P}(X) \rightarrow M$, we will write $\delta(x_1, \dots, x_k)$ instead of $\delta(\{x_1, \dots, x_k\})$.

A *symbolic ultrametric* (Böcker and Dress, 1998) is a 2-way symbolic map $D : \binom{X}{2} \rightarrow M$ satisfying:

- (U1) For all three distinct elements $x, y, z \in X$, at least two of the three values $D(x, y)$, $D(y, z)$ and $D(x, z)$ are the same.
- (U2) There exists no four pairwise distinct elements $x, y, z, u \in X$ such that $D(x, y) = D(y, z) = D(z, u) \neq D(z, x) = D(x, u) = D(u, y)$.

Suppose that T is a tree. Then we denote by $L(T)$ the set of leaves of T and by $V^\circ(T) := V(T) - L(T)$ the set of *internal vertices* of T . If T is rooted then we denote by ρ_T the root of T . Moreover, for any two distinct leaves x and y in T , we define the *least common ancestor* $\text{lca}_T(x, y)$ of x and y in T to be the last vertex in T that lies on both of the paths which start at ρ_T and end in x and in y . If $\text{lca}_T(x, y)$ is adjacent with both x and y then we call the set $\{x, y\}$ a *cherry* of T . We also say that vertex v in T lies *below* a vertex $w \neq v$ in T if w lies on the path from the root of T to v .

A (*rooted/unrooted*) *phylogenetic tree* T on X is a (*rooted/unrooted*) tree with leaf-set X that does not contain vertices of degree two in case T is unrooted and no vertex with indegree and outdegree one in case T is rooted. Note that we will only use the terms rooted or unrooted in case it is not clear from the context which type of tree we are considering. Two phylogenetic trees T and T' on X are *isomorphic* if there exists a bijection $V(T) \rightarrow V(T')$ that induces a graph isomorphism between T and T' that is the identity on X (i.e. the map which takes every element in X to itself). In case T is a rooted phylogenetic tree on X , and Y is a subset of X with size at least two, we let T_Y denote the phylogenetic tree spanned by Y (obtained by suppressing vertices with indegree and outdegree one), and say that T_Y is *induced by* Y .

A *labelled (rooted/unrooted) tree* \mathcal{T} on X is a pair (T, t) , where T is a (*rooted/unrooted*) phylogenetic tree on X , and t is a *labelling map on* M , that is, a map from the internal vertices of T to a set M of symbols. If $t(u) \neq t(v)$ for every $u \neq v$ contained in the same edge of T , we say that \mathcal{T} is *discriminating*. A labelled rooted tree $\mathcal{T} = (T, t)$ on X is a *representation* of a (two-way) symbolic map $D : \binom{X}{2} \rightarrow M$ (or \mathcal{T} *represents* D) if for all distinct $x, y \in X$, we have $D(x, y) = t(\text{lca}_T(x, y))$.

Theorem 2.1 [Böcker and Dress, 1998] *Let $D : \binom{X}{2} \rightarrow M$ be a symbolic map. There exists a discriminating labelled rooted tree \mathcal{T} that represents D if and only if D is a symbolic ultrametric. If this holds, then such a tree is necessarily unique.*

Interestingly, Theorem 2.1 appeared in a different guise in Gurvich (1984) (see also Gurvich, 2000, for more details) in the context of game theory. Within this context, the leaves of the tree T are seen as end of game situations, the label set M corresponds to a set of players, and a directed path from the root of T to a leaf is a sequence of plays. As we will come back to this correspondence in Section 3, we

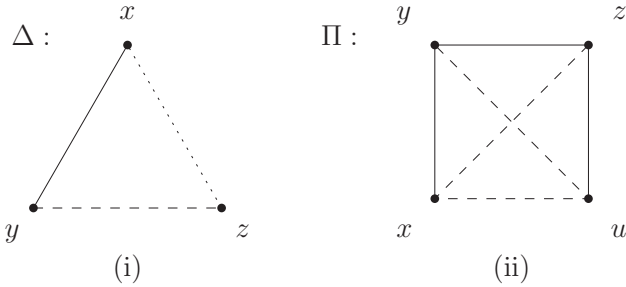


Figure 2. (i) An edge-colored graph Δ on $X = \{x, y, z\}$ adapted from Gurvich (1984, Figure 1). (ii) An edge-colored graph Π on $X = \{x, y, z, u\}$ adapted from the same figure. Colors are represented in terms of different edge styles (plain, dashed and dotted).

now review some relevant terminology and results presented in Gurvich (1984).

Suppose that H is an *edge-labelled graph* on X , that is a graph with vertex set $V(H) = X$ and edge set $E(H) = \binom{X}{2}$, equipped with a map $D : E(H) \rightarrow M$ for a given, nonempty set M . Then D is a (two-way) symbolic map on X , and any symbolic map $D : \binom{X}{2} \rightarrow M$ can trivially be seen as an edge-labelled graph (H, D) on X .

An edge-labelled graph (H, D) on X is said to be *linked* if for all $m \in M$, the graph H_m obtained from H by removing all edges $e \in E(H)$ for which $D(e) = m$ holds is connected. For example, both graphs Δ and Π depicted in Figure 2 are linked. If (H, D) does not contain any linked subgraph, it is said to be *separated*.

The following result from Gurvich (1984) links the property for an edge labelled graph to be separated with the representability of the symbolic map it induces. For this, it relies on the equivalence between the following three statements for a symbolic map $D : \binom{X}{2} \rightarrow M$ (see Gurvich, 1984, Theorem 2 for the equivalence between (ii) and (iii), and Gurvich, 1984, Theorem 4) for the equivalence between (i) and (ii), where a discriminating labelled rooted tree is called a *positional structure*, or *PS* for short):

- (i) There exists a (unique) discriminating labelled rooted tree \mathcal{T} that represents D .
- (ii) The edge-labelled graph (H, D) is separated.
- (iii) The edge-labelled graph (H, D) does not contain any subgraph isomorphic to Δ or Π (depicted in Figure 2).

As mentioned above, this result, and in particular the equivalence between conditions (i) and (iii), provides a direct equivalent to Theo-

rem 2.1. Indeed, as is easy to see a symbolic map $D : X^2 \rightarrow M$ satisfies (U1) (*resp.* (U2)) if and only if the edge-labelled graph (H, D) does not contain a subgraph isomorphic to Δ (*resp.* Π), implying that condition (iii) and the property of being a symbolic ultrametric are equivalent.

3. Three-Way Symbolic Tree-Maps

We begin by considering three-way symbolic maps that arise from labelled unrooted trees. Such a tree $\mathcal{T} = (T, t)$ clearly gives rise to a three-way symbolic map $\delta_{\mathcal{T}} : \binom{X}{3} \rightarrow M$ by putting, for all $x, y, z \in X$, $\delta_{\mathcal{T}}(x, y, z) = t(\text{med}_T(x, y, z))$, where $\text{med}_T(x, y, z)$ denotes the *median vertex* of x, y, z in T (that is, the unique vertex lying on the paths from x to y , from x to z and from y to z , respectively).

If for a three-way symbolic map $\delta : \binom{X}{3} \rightarrow M$, there exists a labelled unrooted tree \mathcal{T} such that $\delta = \delta_{\mathcal{T}}$, we say that δ is a *three-way symbolic tree-map (on X)*, and that \mathcal{T} is a *representation* of δ (or \mathcal{T} *represents* δ). We now characterize such maps. To do this, we define a *symbolic Farris transform*, the definition of which is adapted from the well-known Farris transform (Gurvich, 2009, p. 149) as follows.

Suppose $\mathcal{T} = (T, t)$ is a labelled unrooted tree on X where $|X| \geq 4$. Put $\delta = \delta_{\mathcal{T}}$. Pick a leaf $r \in X$, and define a rooted phylogenetic tree T_r on $X - \{r\}$ as follows: direct all edges of T away from r , and remove r and its outgoing edge. This induces a bijection ψ_r from the set of internal vertices of T to the set of internal vertices of T_r . Hence the map $t_r : V(T_r) \rightarrow M$ which takes any internal vertex v of T_r to M given by $t_r(v) = t(\psi_r^{-1}(v))$ is well-defined, and the pair $\mathcal{T}_r = (T_r, t_r)$ is a labelled rooted tree.

Now, suppose that δ is the three-way symbolic tree-map that is represented by \mathcal{T} , and that D_r is the symbolic ultrametric on X that is represented by \mathcal{T}_r .

Lemma 3.1 *For all $x, y \in X - \{r\}$ with $|X| \geq 4$, we have $D_r(x, y) = \delta(x, y, r)$.*

Proof. It suffices to note that via the symbolic Farris transform, the median vertex of x, y and r in T becomes the least common ancestor of x and y in T_r . Denoting the latter by v , we then have $D_r(x, y) = t_r(v) = t(\psi_r^{-1}(v)) = t(\text{med}_T(x, y, r)) = \delta(x, y, r)$. ■

Motivated by this observation, for a three-way symbolic map $\delta : \binom{X}{3} \rightarrow M$ and some $r \in X$ where $|X| \geq 4$, we define the map

$$\delta_r : \binom{X - \{r\}}{2} \rightarrow M; \quad \delta_r(x, y) = \delta(x, y, r),$$

for all $x, y \in X$ distinct (which can be considered as a symbolic analogue of the Farris transform as defined in Gurvich (2009, p. 149)). Using Lemma 3.1, we can now prove a uniqueness result.

Proposition 3.2 *Let $\delta : \binom{X}{3} \rightarrow M$ be a three-way symbolic tree-map where $|X| \geq 4$. There exists a unique discriminating labelled unrooted tree \mathcal{T} that represents δ .*

Proof. Let $r \in X$ and consider the map $\delta_r : \binom{X - \{r\}}{2} \rightarrow M$. By Lemma 3.1, δ_r is a symbolic ultrametric, and thus, admits a unique discriminating representation \mathcal{T}_r . Moreover, this representation is obtained from a representation of δ , using the symbolic Farris transform. This operation is clearly invertible, and preserves the property of being discriminating. Thus, the labelled unrooted tree \mathcal{T} obtained from \mathcal{T}_r by inverting the symbolic Farris transform is necessarily the only discriminating representation of δ .
■

We now characterize three-way symbolic tree maps. As we shall explain below, an equivalent characterization appears in Gurvich (1984) in the guise of Theorem 5 of that paper. For the sake of completeness, we present a proof within our framework. Subsequent to this, we explain how the approach in Gurvich (1984) relates to ours.

Theorem 3.3 *Suppose that $|X| \geq 4$ and that $\delta : \binom{X}{3} \rightarrow M$ is a three-way symbolic map. Then δ is a three-way symbolic tree-map if and only if δ satisfies the following two conditions:*

(M1) *For all $\{x, y, z, u\} \in \binom{X}{4}$, either*

$$\delta(x, y, z) = \delta(x, y, u) = \delta(x, z, u) = \delta(y, z, u)$$

or two of these four are equal and so are the remaining two.

(M2) *There does not exist $\{x, y, z, u, v\} \in \binom{X}{5}$ such that*

$$\delta(v, x, y) = \delta(v, y, z) = \delta(v, z, u) \neq \delta(v, z, x) = \delta(v, x, u) = \delta(v, u, y).$$

In order to prove Theorem 3.3, we start with a useful lemma.

Lemma 3.4 *Suppose that $|X| \geq 4$ and that $\delta : \binom{X}{3} \rightarrow M$ is a three-way symbolic map satisfying (M1) and (M2). Then for all $r \in X$, the map δ_r is a symbolic ultrametric.*

Proof. Let $r \in X$. We need to show that δ_r satisfies properties (U1) and (U2).

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To see that δ_r satisfies (U1), consider three elements $x, y, z \in X - \{r\}$. Since δ satisfies (M1) the set $\{\delta(r, x, y), \delta(r, x, z), \delta(r, y, z)\}$ contains at most two distinct elements. As this set is precisely the set $\{\delta_r(x, y), \delta_r(x, z), \delta_r(y, z)\}$, (U1) follows.

To see that (U2) holds, assume for contradiction that there exist four pairwise distinct elements $x, y, z, u \in X - \{r\}$ such that $\delta_r(x, y) = \delta_r(y, z) = \delta_r(z, u) \neq \delta_r(z, x) = \delta_r(x, u) = \delta_r(u, y)$. This implies $\delta(r, x, y) = \delta(r, y, z) = \delta(r, z, u) \neq \delta(r, z, x) = \delta(r, x, u) = \delta(r, u, y)$, which is impossible in view of (M2).

■

Note that the converse of the Lemma 3.4 is not true in general. Consider for example the sets $X = \{1, \dots, n\}$, $n \geq 4$, $M = \{A, B\}$, and the map $\delta : \binom{X}{3} \rightarrow M$ defined for $x, y, z \in X$ by putting $\delta(x, y, z) = A$ if $1 \in \{x, y, z\}$ and $\delta(x, y, z) = B$ otherwise. Clearly, δ does not satisfy (M1), as we have $\delta(1, 2, 3) = \delta(1, 2, 4) = \delta(1, 3, 4) \neq \delta(2, 3, 4)$. However, we have $\delta_1(x, y) = A$ for all $x, y \in X - \{1\}$, which is clearly a symbolic ultrametric. In fact, for any $2 \leq k \leq n$ we have $\delta_k(x, y) = A$ if $1 \in \{x, y\}$ and $\delta_k(x, y) = B$ otherwise and, so, δ_k is also a symbolic ultrametric on $X - \{k\}$.

Armed with Lemma 3.4, we can now prove Theorem 3.3.

Proof. Assume first that δ is a three-way symbolic tree-map, and denote by $\mathcal{T} = (T, t)$ its representation. To see that δ satisfies (M1), consider four pairwise distinct elements $x, y, z, u \in X$. Two cases may occur. If $\text{med}_{\mathcal{T}}(x, y, z) = \text{med}_{\mathcal{T}}(x, y, u) = \text{med}_{\mathcal{T}}(x, z, u) = \text{med}_{\mathcal{T}}(y, z, u)$, it follows immediately that $\delta(x, y, z) = \delta(x, y, u) = \delta(x, z, u) = \delta(y, z, u)$. Otherwise, there exists two pairs, say $\{x, y\}$ and $\{z, u\}$, such that the path between x and y and the path between z and u are disjoint. In this case, we have $\text{med}_{\mathcal{T}}(x, y, z) = \text{med}_{\mathcal{T}}(x, y, u) \neq \text{med}_{\mathcal{T}}(x, z, u) = \text{med}_{\mathcal{T}}(y, z, u)$. If $t(\text{med}_{\mathcal{T}}(x, y, z)) = t(\text{med}_{\mathcal{T}}(x, z, u))$, it follows that $\delta(x, y, z) = \delta(x, y, u) = \delta(x, z, u) = \delta(y, z, u)$. Otherwise, we have $\delta(x, y, z) = \delta(x, y, u) \neq \delta(x, z, u) = \delta(y, z, u)$. Thus, δ satisfies (M1).

If $|X| = 4$ then it is straight forward to check that the theorem holds. So assume that $|X| \geq 5$. To see that δ satisfies (M2), assume for contradiction that there exist pairwise distinct $x, y, z, u, v \in X$ such that $\delta(v, x, y) = \delta(v, y, z) = \delta(v, z, u) \neq \delta(v, z, x) = \delta(v, x, u) = \delta(v, u, y)$. We can apply the symbolic Farris transform to \mathcal{T} and v , thus obtaining a labelled rooted tree \mathcal{T}_v . By Lemma 3.1, \mathcal{T}_v is a representation of δ_v , implying that δ_v is a symbolic ultrametric. But, by definition, δ_v satisfies $\delta_v(x, y) = \delta_v(y, z) = \delta_v(z, u) \neq \delta_v(z, x) = \delta_v(x, u) = \delta_v(u, y)$, which contradicts (U2).

Conversely, assume that δ satisfies Properties (M1) and (M2), and let $r \in X$. By Lemma 3.4, the map δ_r is a symbolic ultrametric. Thus, there exists a labelled rooted tree $\mathcal{T}_r = (T_r, t_r)$ on $X - \{r\}$ representing δ_r . Consider the labelled unrooted tree $\mathcal{T} = (T, t)$ on X defined as follows. First, add a new vertex r to T_r and the edge $\{\rho_{T_r}, r\}$. Then consider all edges in the resulting tree to be undirected. Let $t : V^o(T) \rightarrow M$ denote the map given by $t(v) = t_r(v)$, for all $v \in V^o(T)$. We claim that for all $\{x, y, z\} \in \binom{X}{3}$, we have $\delta(x, y, z) = t(\text{med}_T(x, y, z))$, that is, \mathcal{T} is a representation of δ . To prove this, it suffices to consider two cases. Suppose $\{x, y, z\} \in \binom{X}{3}$.

Case (a): $\{x, y, z\} \subseteq X - \{r\}$. Without loss of generality, $\delta_r(x, z) = \delta_r(y, z) = t_r(u)$ and $\delta_r(x, y) = t_r(v)$, where u and v are vertices of T_r , and v is below or equal to u in T_r . In this case, $t(\text{med}_T(x, y, z))$ equals $t_r(v)$. By (M1) and since $\delta(x, z, r) = \delta(y, z, r)$, we have $\delta(x, y, z) = \delta(x, y, r) = t_r(v) = t(\text{med}_T(x, y, z))$. Thus, \mathcal{T} is a representation of δ in this case.

Case (b): $r \in \{x, y, z\}$, say $r = z$. If we denote by v the least common ancestor of x and y in T_r , then $t(\text{med}_T(x, y, z)) = t_r(v)$. Hence $\delta(x, y, r) = \delta_r(x, y) = t_r(v) = t(\text{med}_T(x, y, r))$. Thus, \mathcal{T} is a representation of δ in this case, too.

■

We next elaborate on the relationship between Theorem 3.3 and Theorem 5 in Gurvich (1984). As mentioned in Section 2, a symbolic two-way map $D : \binom{X}{2} \rightarrow M$ can be seen as an edge-labelled graph (H, D) . Similarly, a symbolic three-way map $\delta : \binom{X}{3} \rightarrow M$ can be seen as an edge-labelled 3-hypergraph (\mathcal{H}, δ) , where by 3-hypergraph, we mean that the edges of \mathcal{H} are sets of three vertices (instead of two for graphs). Within this context, the vertex set of a 3-hypergraph \mathcal{H} associated to a 3-way map δ is X , as in the case of two-way maps, and the edge set of \mathcal{H} is $\binom{X}{3}$.

Gurvich (1984) uses as a starting point for his characterization the equivalence, presented in Section 2, between symbolic 2-way maps that can be represented by a rooted tree and edge-labelled graphs that do not contain any subgraph isomorphic to the graphs Δ or Π (see Figure 2). The idea underlying Gurvich (1984, Theorem 5) is the following. From an edge-labelled 3-hypergraph (\mathcal{H}, δ) on X , we can pick an element $r \in X$ and consider the edge-labelled graph (H, δ_r) on $X - \{r\}$. It is then possible to highlight three edge-labelled 3-hypergraph $\delta_2, \delta_3, \delta_4$ with four vertices, that get transformed into edge-labelled graphs isomorphic to Δ via that operation, and one edge-labelled 3-hypergraph π with five

vertices, that gets transformed into an edge-labelled graph isomorphic to Π . The equivalence between the representability of δ by a labelled unrooted tree and the representability of δ_r by a labelled rooted tree for all $r \in X$ then leads to the conclusion that an edge-labelled 3-hypergraph (\mathcal{H}, δ) is representable if and only if it does not contain a sub(hyper)graph isomorphic to any of $\delta_2, \delta_3, \delta_4$ and π .

As it turns out, (\mathcal{H}, δ) contains a sub(hyper)graph isomorphic to one of $\delta_2, \delta_3, \delta_4$ (*resp.* to π) if and only if δ does not satisfy (M1) (*resp.* (M2)). This implies that Theorem 3.3 and Theorem 5 in Gurvich (1984) are equivalent.

Finally, note that a similar result also appears in Grunewald, Long, and Wu (2017). However, the arguments used by Grunewald et al. (2017) do not rely on the projection of a three-way map to a two-way map and of an unrooted tree to a rooted tree, as is the case both here and in Gurvich (1984).

4. Three-Way Symbolic Ultrametrics

In the last section, we considered the problem of deciding when a three-way symbolic map arises from a labelled unrooted tree. In this section, we start to consider this problem for their rooted counterparts. In particular, after defining the concept of a three-way symbolic ultrametric, we shall show that to determine whether or not a three-way symbolic map is a symbolic ultrametric, it suffices to consider its restriction to sets of size five.

We begin by considering how to define a three-way symbolic ultrametric. If we consider 3 distinct leaves x, y, z of a rooted phylogenetic tree T on X , then we can clearly identify two internal vertices of the tree given by the set $\{\text{lca}_T(x, y), \text{lca}_T(x, z), \text{lca}_T(y, z)\}$ (in contrast to unrooted phylogenetic trees where we can identify only one, namely the median of the 3 leaves). A natural approach to obtain a three-way symbolic map δ from a labelled rooted tree $\mathcal{T} = (T, t)$ might therefore be to take $\delta(x, y, z)$ to be the set $\{t(\text{lca}_T(x, y)), t(\text{lca}_T(x, z)), t(\text{lca}_T(y, z))\}$, for $x, y, z \in X$ distinct. However, as can be seen in Figure 3, such a map does not necessarily uniquely capture \mathcal{T} . For this reason, we shall consider instead maps to multisets.

To formalize this, let $\mathcal{M} = \mathcal{M}_M$ denote the set of multisets $\{a, b, c\}$ with $a, b, c \in M$. As it will be useful later on, we shall also sometimes denote an element in \mathcal{M} as a sum. So, for example, for the element $\{a, a, b\} \in \mathcal{M}$ with $a, b \in M$, we sometimes also write $2a + b$.

Now, given a labelled rooted tree $\mathcal{T} = (T, t)$ on X , we define the three-way symbolic map

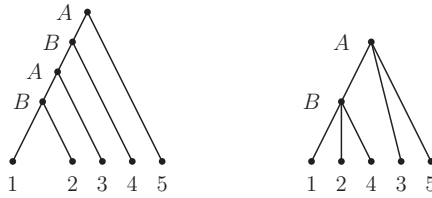


Figure 3. Two labelled rooted trees on $X = \{1, 2, 3, 4, 5\}$ with labelling maps t and t' on $M = \{A, B\}$, respectively, for which the sets $\{t(\text{lca}_T(x, y)), t(\text{lca}_T(x, z)), t(\text{lca}_T(y, z))\}$ and $\{t'(\text{lca}_T(x, y)), t'(\text{lca}_T(x, z)), t'(\text{lca}_T(y, z))\}$ coincide, for any three elements $x, y, z \in X$ distinct.

$\delta_{\mathcal{T}} : \binom{X}{3} \rightarrow \mathcal{M}$ by putting

$$\delta_{\mathcal{T}}(x, y, z) = \{t(\text{lca}_T(x, y)), t(\text{lca}_T(x, z)), t(\text{lca}_T(y, z))\}.$$

for all distinct $x, y, z \in X$. If for a three-way symbolic map $\delta : \binom{X}{3} \rightarrow \mathcal{M}$ there exists a labelled rooted tree $\mathcal{T} = (T, t)$ on X such that $\delta = \delta_{\mathcal{T}}$, then we call δ a *three-way symbolic ultrametric (on X)*. Thus, intuitively, δ is a three-way symbolic ultrametric if it can be represented by labelling a rooted tree on X in such a way that, for every 3-subset $\{x, y, z\}$ of X , $\delta(x, y, z)$ is the multiset consisting of the labels of the least common ancestors for all pairs of elements in $\{x, y, z\}$. In addition, we say that \mathcal{T} is a *representation* for δ (or that \mathcal{T} *represents* δ). We say that \mathcal{T} is *discriminating* if $t(u) \neq t(v)$, for every $u \neq v$ contained in the same edge in T . Note that we can think of δ as a symbolic analogue of a three-way perimeter map which arises from a weighted tree T by taking, for any three leaves of T , the length of subtree spanned by the those leaves (see e.g. Chepoi and Fichet, 2007). Also, note that by (U1), δ must satisfy the following property:

Lemma 4.1 *Let $\delta : \binom{X}{3} \rightarrow \mathcal{M}$ be a three-way symbolic ultrametric. Then, for any three distinct elements $x, y, z \in X$, the number of distinct elements in the multiset $\delta(x, y, z)$ is at most two.*

We now turn our attention to showing that we can determine whether or not a three-way symbolic map $\delta : \binom{X}{3} \rightarrow \mathcal{M}$ is a symbolic ultrametric by restricting δ to subsets of X with size five. To do this, we first need to introduce some additional notation. For a subset Y of X of size four or more, let $\delta|_Y$ denote the restriction of δ to $\binom{Y}{3}$, that is, the map obtained by restricting the map δ to the subset $\binom{Y}{3}$ of $\binom{X}{3}$. Note that if δ is a three-way symbolic ultrametric, then $\delta|_Y$ is a three-way symbolic ultrametric for all subsets $Y \subseteq X$ with $|Y| \geq 4$. Indeed,

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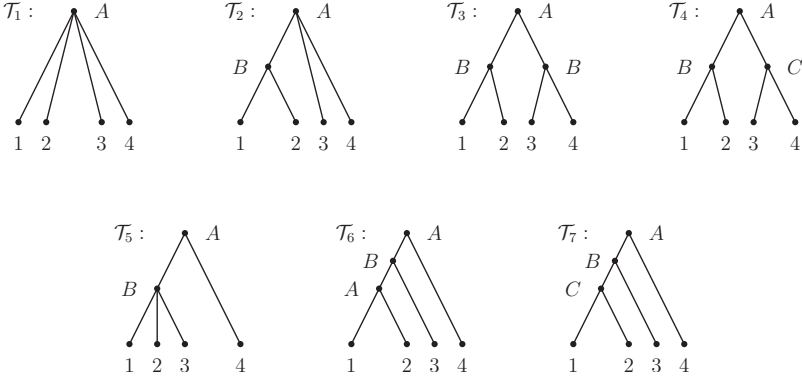


Figure 4. All possible discriminating labelled rooted trees \mathcal{T}_i , $1 \leq i \leq 7$, on $\{1, 2, 3, 4\}$, up to a relabelling of the leaves.

Table 1. For $1 \leq i \leq 7$ and $M = \{A, B, C\}$, the values of the map $\hat{\delta}_i$ represented by the labelled rooted trees \mathcal{T}_i in Figure 4. The trees \mathcal{T}_i are given in terms of their index i in the top row.

i	1	2	3	4	5	6	7
$\hat{\delta}_i(1, 2, 3)$	3A	2A+B	2A+B	2A+B	3B	A+2B	2B+C
$\hat{\delta}_i(1, 2, 4)$	3A	2A+B	2A+B	2A+B	2A+B	3A	2A+C
$\hat{\delta}_i(1, 3, 4)$	3A	3A	2A+B	2A+C	2A+B	2A+B	2A+B
$\hat{\delta}_i(2, 3, 4)$	3A	3A	2A+B	2A+C	2A+B	2A+B	2A+B

if \mathcal{T} is a representation of δ , then the subtree \mathcal{T}_Y of \mathcal{T} induced by Y is a representation of $\delta|_Y$. Furthermore, we obtain a discriminating representation of $\delta|_Y$ by collapsing all edges of \mathcal{T}_Y both of whose end vertices have the same label.

We now consider symbolic ultrametrics on a set of size four. For $M = \{A, B, C\}$, in Figure 4, we picture all possible discriminating labelled rooted trees \mathcal{T}_i , $1 \leq i \leq 7$, on $\{1, 2, 3, 4\}$ and in Table 1, we list for all $1 \leq i \leq 7$ the values of the map $\hat{\delta}_i : \binom{X}{3} \rightarrow \mathcal{M}$ that is represented by \mathcal{T}_i . As we can see from this table, all of the maps $\hat{\delta}_i$ except for $\hat{\delta}_3$ capture \mathcal{T}_i (in the sense that \mathcal{T}_i is the unique labelled rooted tree on $\{1, 2, 3, 4\}$ that represents $\hat{\delta}_i$). Now, for $Y \subseteq X$ of size four, we say that $\delta|_Y$ is of type $\hat{\delta}_i$, $i \in \{1, \dots, 7\}$ if there exists a bijection between Y and $\{1, 2, 3, 4\}$ that induces a bijection between the image of $\delta|_Y$ and the image of $\hat{\delta}_i$ such that $\delta|_Y$ and $\hat{\delta}_i$ coincide up to these bijections. Since Table 1 is exhaustive, we have:

Proposition 4.2 *Suppose that $|X| \geq 4$, that $\delta : \binom{X}{3} \rightarrow \mathcal{M}$ is a three-way symbolic map, and that $Y \subseteq X$ is a subset of size four. Then $\delta|_Y : \binom{Y}{3} \rightarrow \mathcal{M}$ is a three-way symbolic ultrametric on Y if and only if there exists some $i \in \{1, \dots, 7\}$ such that $\delta|_Y$ is of type $\hat{\delta}_i$. Moreover, if $i \neq 3$, the representation of $\delta|_Y$ is unique.*

We now turn our attention to symbolic ultrametrics on sets of size five. In the last result, we have seen that a three-way symbolic ultrametric on a set of size 4 may have more than one representation by a labelled tree. However, as we shall now show this can not happen for sets of size five.

Lemma 4.3 *If Y is a set of size five and $\delta : \binom{Y}{3} \rightarrow \mathcal{M}$ is a three-way symbolic ultrametric on Y , then δ has a unique discriminating representation.*

Proof. Suppose that δ is a three-way symbolic ultrametric on Y , and that \mathcal{T} is a discriminating representation of δ . Let $D = D_{\mathcal{T}} : \binom{Y}{2} \rightarrow M$ be the symbolic ultrametric represented by \mathcal{T} . By Theorem 2.1, it suffices to show that if δ is also represented by a labelled tree \mathcal{T}' , then $D_{\mathcal{T}'} = D_{\mathcal{T}}$.

Since δ is a three-way symbolic ultrametric on Y , there exists a subset Y_0 of Y with $|Y_0| = 4$ such that $\delta|_{Y_0}$ is not of type $\hat{\delta}_3$. Thus, by Proposition 4.2, $D_{\mathcal{T}'|_{Y_0}} = D_{\mathcal{T}|_{Y_0}}$. Hence, $D_{\mathcal{T}'}(x_0, x) = D_{\mathcal{T}}(x_0, x)$ for all $x \in Y_0$ where x_0 is the unique element contained in $Y - Y_0$, since the value of $D_{\mathcal{T}'}(x_0, x)$ is given by δ and $D_{\mathcal{T}|_{Y_0}}$ as follows. Let $Y_0 = \{x, y, z, u\}$ and consider the multisets $\delta(x, y, x_0) - D_{\mathcal{T}|_{Y_0}}(x, y)$, $\delta(x, z, x_0) - D_{\mathcal{T}|_{Y_0}}(x, z)$ and $\delta(x, u, x_0) - D_{\mathcal{T}|_{Y_0}}(x, u)$ where for a multiset A with $k \geq 1$ copies of some element a , we denote by $A - a$ the multiset obtained by removing one copy of a . If there exists a unique element $c \in M$ that belongs to all three of these sets, we have $D_{\mathcal{T}'}(x_0, x) = c$. If two distinct elements of M share this property, this implies $D_{\mathcal{T}'}(x_0, y) = D_{\mathcal{T}'}(x_0, z) = D_{\mathcal{T}'}(x_0, u) \neq D_{\mathcal{T}'}(x_0, x)$. We then have $D_{\mathcal{T}'}(x_0, y) = m(\delta(y, z, x_0))$, and $D_{\mathcal{T}'}(x_0, x)$ is the single element of $\delta(x, y, x_0) - \{D_{\mathcal{T}|_{Y_0}}(x, y), D_{\mathcal{T}'}(x_0, y)\}$.

■

Note that, as the example in Table 2 shows, it is not true in general that a three-way symbolic map δ that restricts to a three-way symbolic ultrametric on all subsets Y of X of size four is a three-way symbolic ultrametric on X . However, as mentioned above, using the previous lemma we now show that considering sets of size five is enough to ensure that this is the case.

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Table 2. For $M = \{A, B\}$ and $X = \{1, 2, 3, 4, 5\}$, a three-way symbolic map $\delta : \binom{X}{3} \rightarrow \mathcal{M}$ which is not a three-way symbolic ultrametric on X but whose restriction to any subset $Y \subset X$ of size four is a three-way symbolic ultrametric on Y .

$\delta(1, 2, 3)$	2A+B	$\delta(1, 4, 5)$	2A+B
$\delta(1, 2, 4)$	2A+B	$\delta(2, 3, 4)$	3A
$\delta(1, 2, 5)$	3B	$\delta(2, 3, 5)$	2A+B
$\delta(1, 3, 4)$	3A	$\delta(2, 4, 5)$	2A+B
$\delta(1, 3, 5)$	2A+B	$\delta(3, 4, 5)$	A+2B

Theorem 4.4 *Suppose that $|X| \geq 5$ and that $\delta : \binom{X}{3} \rightarrow \mathcal{M}$ is a three-way symbolic map. Then, δ is a three-way symbolic ultrametric if and only if $\delta|_Y$ is a three-way symbolic ultrametric for all $Y \subseteq X$ of size five.*

Proof. The fact that a three-way symbolic ultrametric on X restricts to such an ultrametric on all subsets of X of size five is clear.

Conversely, assume that $\delta|_Y$ is a three-way symbolic ultrametric for all $Y \subseteq X$ of size five. For such a set Y , we denote by $\mathcal{T}_Y = (T_Y, t_Y)$ the unique (by Lemma 4.3) discriminating labelled tree that represents $\delta|_Y$, and by D_Y the symbolic ultrametric that is represented by \mathcal{T}_Y .

Clearly, if there exists a map $D : \binom{X}{2} \rightarrow M$ such that, for all subsets $Y \subseteq X$ of size five, the restriction of D to $\binom{Y}{2}$ coincides with D_Y , then D satisfies $\delta(x, y, z) = \{D(x, y), D(x, z), D(y, z)\}$, for all $x, y, z \in X$ pairwise distinct. Moreover, since D_Y is a symbolic ultrametric on any subset $Y \subseteq X$ of size five, and given that the property of being a symbolic ultrametric is based on a 4-point condition, we have that such a map D , if it exists, is also a symbolic ultrametric. Thus, if D exists, then δ is a three-way symbolic ultrametric.

To show that D exists, assume for contradiction that there exist x and y in X and two distinct subsets Y_1 and Y_2 of X of size five, both containing x and y , such that $D_{Y_1}(x, y) \neq D_{Y_2}(x, y)$. We may assume without loss of generality that $I = Y_1 \cap Y_2$ has size four. Moreover, we claim that x, y, Y_1 and Y_2 can be chosen in such a way that $\delta|_I$ is not of type $\hat{\delta}_3$, as defined in Table 1.

To prove this claim, consider the case where $\delta|_I$ is of type $\hat{\delta}_3$ (otherwise, the claim trivially holds). Assume $Y_1 = \{x, y, z, t, u_1\}$ and $Y_2 = \{x, y, z, t, u_2\}$, which implies $I = \{x, y, z, t\}$. Both the subtree of \mathcal{T}_{Y_1} induced by I and the subtree of \mathcal{T}_{Y_2} induced by I are of the form \mathcal{T}_3 in Figure 4, and their underlying phylogenetic trees are not isomorphic. We can assume that one has cherries $\{x, y\}$ and $\{t, z\}$ and the other has cherries $\{x, z\}$ and $\{t, y\}$. Then, we have not only that

$D_{Y_1}(x, y) \neq D_{Y_2}(x, y)$, but also that $D_{Y_1}(x, z) \neq D_{Y_2}(x, z)$, $D_{Y_1}(z, t) \neq D_{Y_2}(z, t)$, and $D_{Y_1}(y, t) \neq D_{Y_2}(y, t)$. Moreover, it is easy to check that there exists a subset $Y^* \subset I$ of size three such that neither $\delta|_{Y^* \cup \{u_1\}}$ nor $\delta|_{Y^* \cup \{u_2\}}$ is of type $\hat{\delta}_3$.

Since Y^* is a subset of I of size three and, in view of the four inequalities listed above, there exists two elements $x', y' \in Y^*$ such that $D_{Y_1}(x', y') \neq D_{Y_2}(x', y')$. If we denote by Y' the set $Y^* \cup \{u_1\} \cup \{u_2\}$, we have that both $Y' \cap Y_1$ and $Y' \cap Y_2$ have size four, and that at least one of $D_{Y'}(x', y') \neq D_{Y_1}(x', y')$ or $D_{Y'}(x', y') \neq D_{Y_2}(x', y')$ holds. If the first inequality holds, the claim is then satisfied for x', y', Y' and Y_1 . Otherwise, it is satisfied for x', y', Y' and Y_1 , which completes the proof of the claim.

Now, in light of the claim, the representation \mathcal{T}_I of $\delta|_I$ is unique, and so is the symbolic ultrametric D^I that is represented by \mathcal{T}_I . Moreover, D^I is precisely the restriction of D_{Y_1} to I , and the restriction of D_{Y_2} to I . In particular, we have $D(x, y) = D_{Y_1}(x, y)$ and $D(x, y) = D_{Y_2}(x, y)$, which contradicts $D_{Y_1}(x, y) \neq D_{Y_2}(x, y)$.
■

5. A Five-Point Characterization of Three-Way Symbolic Ultrametrics

We now focus on using the results in the previous two sections to derive conditions for characterizing three-way symbolic ultrametrics that are analogous to conditions (U1) and (U2) for symbolic ultrametrics.

In the following, we shall consider expressions of the form $\sum_{m \in M} \alpha_m m$, where α_m is a real number, which arise when we take linear combinations of multisets in \mathcal{M} . We shall say that such an expression $\sum_{m \in M} \alpha_m m$ is *valid for M* if the coefficient for each element in M is contained in \mathbb{N} . For example, for $M = \{a, b\}$, if $S_1 = 2a + b$, $S_2 = 2b + a$ and $S_3 = 3a$ are multisets in \mathcal{M} , then we have $\frac{1}{3}(S_1 + S_2) = a + b$, which is valid for M , but $S_3 - S_1 = a - b$ and $\frac{1}{2}(S_1 + S_3) = \frac{5}{2}a + \frac{1}{2}b$ which are not valid for M .

Now, suppose that $\delta : \binom{X}{3} \rightarrow \mathcal{M}$ is a three-way symbolic map where $|X| \geq 5$. Let $Y = \{x, y, z, u, v\}$ be a subset of X . Let $\nu_Y(\delta)$ denote the vector

$$(\delta(x, y, z), \delta(x, y, u), \dots, \delta(z, u, v)).$$

In addition, suppose that $D_Y : \binom{Y}{2} \rightarrow M$ is a map such that

$$\delta(a, b, c) = \{D_Y(a, b), D_Y(a, c), D_Y(b, c)\}$$

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for all $a, b, c \in Y$, and let $\mu_Y(\delta)$ denote the vector

$$(D_Y(x, y), D_Y(x, z), \dots, D_Y(u, v)).$$

By definition of D_Y , it is straight-forward to check that $A\mu_Y(\delta) = \nu_Y(\delta)$, where

$$A = \begin{pmatrix} 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 \end{pmatrix}.$$

Note that in Hermann et al. (2012) it was shown that the matrix A is invertible with inverse

$$A^{-1} = \frac{1}{6} \begin{pmatrix} 2 & 2 & 2 & -1 & -1 & -1 & -1 & -1 & -1 & 2 \\ 2 & -1 & -1 & 2 & 2 & -1 & -1 & -1 & 2 & -1 \\ -1 & 2 & -1 & 2 & -1 & 2 & -1 & 2 & -1 & -1 \\ -1 & -1 & 2 & -1 & 2 & 2 & 2 & -1 & -1 & -1 \\ 2 & -1 & -1 & -1 & -1 & 2 & 2 & 2 & -1 & -1 \\ -1 & 2 & -1 & -1 & 2 & -1 & 2 & -1 & 2 & -1 \\ -1 & -1 & 2 & 2 & -1 & -1 & -1 & 2 & 2 & -1 \\ -1 & -1 & 2 & 2 & -1 & -1 & 2 & -1 & -1 & 2 \\ -1 & 2 & -1 & -1 & 2 & -1 & -1 & 2 & -1 & 2 \\ 2 & -1 & -1 & -1 & -1 & 2 & -1 & -1 & 2 & 2 \end{pmatrix}.$$

Consider the product $\mu_Y(\delta) = A^{-1}\nu_Y(\delta)$. Then, as the rows of A^{-1} are indexed by pairs of distinct elements in Y , it is straight-forward to check by considering the $\{p, q\}$ th row of A^{-1} (for $p \neq q \in Y$) and putting $\{e, f, g\} = Y - \{p, q\}$ and

$$\begin{aligned} S_{p,q}^Y(\delta) &= \frac{1}{6}(2(\delta(p, q, e) + \delta(p, q, f) + \delta(p, q, g) + \delta(e, f, g)) \\ &\quad - \sum_{a,b \in Y - \{p,q\}} (\delta(p, a, b) + \delta(q, a, b))), \end{aligned}$$

that $S_{p,q}^Y = \{D_Y(p, q)\}$. Defining $S_{p,q}^Y$ as above for $Y \subseteq X$ with $|Y| = 5$ and $p \neq q \in Y$ we also have:

Proposition 5.1 *Suppose that $|X| \geq 5$, that $\delta : \binom{X}{3} \rightarrow \mathcal{M}$ is a three-way symbolic map, and that $Y \subseteq X$ has size five. There exists a map $D^Y : \binom{Y}{2} \rightarrow M$ such that $\delta(a, b, c) = \{D^Y(a, b), D^Y(a, c), D^Y(b, c)\}$ for all $a, b, c \in Y$ if and only if for all $p, q \in Y$ distinct, $S_{p,q}^Y(\delta)$ is valid for M , in which case $S_{p,q}^Y(\delta)$ is a singleton multiset.*

Proof. Suppose first that the map D^Y exists. Without loss of generality we may assume that $D^Y = D_Y$. In view of the discussion preceding the proposition, it follows that $S_{p,q}^Y(\delta)$ is valid for M for all $p, q \in Y$ distinct, as $S_{p,q}^Y(\delta) = \{D_Y(p, q)\}$.

To see the converse, assume that $S_{p,q}^Y(\delta)$ is valid for M for all $p \neq q \in Y$. Fix p and q . We claim that $S_{p,q}^Y(\delta)$ is a singleton multiset. To see this, put $\mathcal{A} = 2(\delta(p, q, e) + \delta(p, q, f) + \delta(p, q, g) + \delta(e, f, g))$ and $\mathcal{B} = \sum_{a,b \in Y - \{p,q\}} (\delta(p, a, b) + \delta(q, a, b))$. Then since $S_{p,q}^Y(\delta)$ is valid for M , every element in \mathcal{B} must also be an element in \mathcal{A} . Hence, $S_{p,q}^Y(\delta)$ must contain $\frac{1}{6}|\mathcal{A} - \mathcal{B}| = 1$ element as $|\mathcal{A}| = 24$ and $|\mathcal{B}| = 18$. This proves the claim.

Now, it is straight forward to see that if $S_{p,q}^Y(\delta) = \{s_{p,q}^Y\}$, for $p \neq q \in Y$, then the map $D^Y : \binom{Y}{2} \rightarrow M$ defined by putting $D^Y(p, q) = s_{p,q}^Y(\delta)$, for all $p \neq q \in Y$, satisfies the stated property. ■

We now present conditions for characterizing when a three-way symbolic map is a three-way symbolic ultrametric. For $\Sigma \in \mathcal{M}$, we define the elements $m(\Sigma)$ and $n(\Sigma)$ of M as follows:

- If Σ contains a single element $A \in M$ repeated three times, we put $m(\Sigma) = n(\Sigma) = A$.
- If Σ contains two distinct elements, we define $m(\Sigma)$ as the element of Σ appearing twice and $n(\Sigma)$ as the element appearing only once.
- If Σ contains three distinct elements, we put $m(\Sigma) = n(\Sigma) = \emptyset$.

Note that if Σ contains two or fewer distinct elements, then $\Sigma = \{m(\Sigma), m(\Sigma), n(\Sigma)\}$.

Theorem 5.2 *Suppose that $|X| \geq 5$ and that $\delta : \binom{X}{3} \rightarrow \mathcal{M}$ is a three-way symbolic map. Then δ is a three-way symbolic ultrametric if and only if the following hold:*

- (P1) *For all subsets $Y \subseteq X$ of size five and all $x, y \in Y$ distinct, $S_{x,y}^Y(\delta)$ is valid for M .*
- (P2) *For all pairwise distinct $x, y, z \in X$, $\delta(x, y, z)$ contains at most two distinct elements.*

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(P3) For all pairwise distinct $x, y, z, u \in X$ with $\delta(x, y, z) = \delta(y, z, u) \neq \delta(x, y, u) = \delta(x, z, u)$ holding, we have $m(\delta(x, y, z)) = m(\delta(x, y, u))$.

Proof. Assume first that δ is a three-way symbolic ultrametric. By Theorem 4.4 and Proposition 5.1 it follows that Properties (P1) and (P2) must hold. To see that Property (P3) holds too let $\{x, y, z, u\} \in \binom{X}{4}$ be such that $\delta(x, y, z) = \delta(y, z, u) \neq \delta(x, y, u) = \delta(x, z, u)$. Since $\delta|_{\{x, y, z, u\}}$ is a three-way symbolic ultrametric, Proposition 4.2 combined with Table 1 implies that $\delta|_{\{x, y, z, u\}}$ is either of type $\hat{\delta}_3$ and $\hat{\delta}_5$. Clearly, $m(\hat{\delta}_i(x, y, z)) = m(\hat{\delta}_i(x, y, u))$ holds for $i = 3, 5$ and, so, Property (P3) follows.

Conversely, assume that δ satisfies Properties (P1) – (P3). Consider a subset $Y \subseteq X$ of size five. By Proposition 5.1, there exists a map $D^Y : \binom{Y}{2} \rightarrow M$ such that $\delta(x, y, z) = \{D^Y(x, y), D^Y(x, z), D^Y(y, z)\}$ for all $x, y, z \in Y$. We claim that D^Y is a symbolic ultrametric. For this it suffices to show that D^Y satisfies Property (U2) as Property (U1) is a direct consequence of Property (P1).

To see that D^Y satisfies Property (U2), assume for contradiction that there exist pairwise distinct $x, y, z, u \in Y$ such that $D^Y(x, y) = D^Y(y, z) = D^Y(z, u) \neq D^Y(z, x) = D^Y(x, u) = D^Y(u, y)$. Put $A = D^Y(x, y)$ and $B = D^Y(z, x)$. Then $\delta(x, y, z) = \delta(y, z, u) = 2A + B \neq A + 2B = \delta(x, y, u) = \delta(x, z, u)$. Since, $m(\delta(x, y, z)) = A \neq B = m(\delta(x, y, u))$ also holds this is impossible in view of Property (P3). Thus, D^Y also satisfies Property (U2) and, so, is a symbolic ultrametric, as claimed.

Since D^Y is a symbolic ultrametric, there exists a labelled rooted tree \mathcal{T} that represents D^Y . Combined with the definition of D^Y it follows that \mathcal{T} also represents $\delta|_Y$. Thus, $\delta|_Y$ is a three-way symbolic ultrametric and, so, $\delta|_Y$ is a three-way symbolic ultrametric for all subsets $Y \subseteq X$ with $|Y| = 5$. By Theorem 4.4, it follows that δ is a three-way symbolic ultrametric. ■

Note that Properties (P1) – (P3) are independent of each other. Indeed, that Property (P2) is independent of Properties (P1) and (P3) and that Property (P3) is independent of Properties (P1) and (P2) is a direct consequence of the fact that Properties (U1) and (U2) are independent of each other.

To see that Property (P1) is independent of Properties (P2) and (P3), consider the three-way symbolic map $\delta : \binom{X}{3} \rightarrow \mathcal{M}_{\{A, B\}}$ defined, for all $x, y, z \in X$, by putting $\delta(x, y, z) = 2A + B$. The map δ always satisfies (P2) and (P3), but if $|X| \geq 5$, δ does not satisfy (P1).

6. Reconstructing Three-Way Symbolic Ultrametric Representations Using Triplets

In this section, we are interested in determining when a three-way symbolic map δ on X is a tree-map or a symbolic ultrametric. Clearly, using the conditions given in Theorem 5.2 this can be done by examining every subset of X with size five. However, we now show how to do this using a triplet-based approach, which essentially reduces the problem to considering subsets of X of size three.

Recall that a *triplet* is a binary phylogenetic tree on three leaves. By $xy|z$ we denote the triplet with leaf-set $\{x, y, z\}$ which has x, y adjacent to the same vertex in the tree. For T a phylogenetic tree on X and $x, y, z \in X$, we say that T *displays* the triplet $xy|z$ if $\text{lca}_T(x, z) = \text{lca}_T(y, z) \neq \text{lca}_T(x, y)$. To keep notation at bay, we sometimes also say that a labelled tree $\mathcal{T} = (T, t)$ on X *displays* a triplet r if r is displayed by T .

In Semple and Steel (2003, Section 7.6), a triplet-based approach is described for deciding whether or not a two-way symbolic map δ is a symbolic ultrametric or not and, if it is, for building a labelled tree which represents δ . This approach is based on the BUILD algorithm, that was presented under that name in Aho et al. (1981). Using the results in Section 3, the BUILD algorithm also allows us to check if a three-way symbolic map is a tree-map using triplets as follows. Suppose $\delta : \binom{X}{3} \rightarrow \mathcal{M}$ is a three-way symbolic map. Pick any $r \in X$. Then, using the BUILD-based approach, we can check whether or not the map δ_r defined in Section 3. is a symbolic ultrametric by taking the set of triplets $xy|z$ with $x, y, z \in X$ distinct, for which $\delta_r(x, y) \neq \delta_r(x, z) = \delta_r(y, z)$ holds as input to BUILD. If this is not the case, then by Lemma 3.1, δ is not a three-way symbolic tree-map. Otherwise, if (T, t) is the representation of δ_r returned by BUILD, then we can simply check whether or not this leads to a representation of δ by attaching the leaf r to the root of T . If this is possible then δ is a three-way symbolic tree-map, otherwise it is not.

We now turn our attention to three-way symbolic ultrametries. We begin by presenting a key link between triplets and such maps whose proof is straight forward. Denote the underlying set of a multiset \mathcal{A} by $\underline{\mathcal{A}}$.

Proposition 6.1 *Let $\mathcal{T} = (T, t)$ be a discriminating labelled tree. For $x, y, z \in X$ distinct:*

(T1) *If $xy|z$ is a triplet displayed by T , then $t(\text{lca}_T(x, z)) = t(\text{lca}_T(y, z)) = m(\delta_{\mathcal{T}}(x, y, z))$ and $t(\text{lca}_T(x, y)) = n(\delta_{\mathcal{T}}(x, y, z))$*

(T2) *If T does not display any triplet on $\{x, y, z\}$, then $|\underline{\delta_{\mathcal{T}}(x, y, z)}| = 1$.*

Corollary 6.2 *Up to isomorphism, any labelled tree \mathcal{T} can be uniquely reconstructed from $\delta_{\mathcal{T}}$ and the set of triplets displayed by \mathcal{T} .*

Proof. Put $\mathcal{T} = (T, t)$ and $\delta = \delta_{\mathcal{T}}$. Let $t : V^o(T) \rightarrow M$ and let \mathcal{R} denote the set of triplets displayed by \mathcal{T} . Define a map $D_{\delta} : \binom{X}{2} \rightarrow M$ as follows. Suppose $x, y \in X$ distinct. If there exists some $z \in X - \{x, y\}$ such that no triplet on $\{x, y, z\}$ is contained in \mathcal{R} , then define $D_{\delta}(x, y)$ to be the element in $\delta_{\mathcal{T}}(x, y, z)$. If there exists some $z \in X - \{x, y\}$ such that $xy|z \in \mathcal{R}$ then put $D_{\delta}(x, y) = n(\delta_{\mathcal{T}}(x, y, z))$ and if $xz|y \in \mathcal{R}$ then put $D_{\delta}(x, y) = m(\delta_{\mathcal{T}}(x, y, z))$. In view of Proposition 6.1, the map D_{δ} is clearly well-defined.

The corollary now follows in view of Theorem 2.1 as D_{δ} is equal to the symbolic ultrametric $D_{\mathcal{T}}$ that is represented by \mathcal{T} (as $D_{\delta}(x, y) = t(lca(x, y)) = D_{\mathcal{T}}(x, y)$ clearly holds for all $x, y \in X$ distinct).

■

In light of Corollary 6.2, it is of interest to understand when, for a labelled tree \mathcal{T} , the set of triplets displayed by \mathcal{T} can be obtained from $\delta_{\mathcal{T}}$. The tree \mathcal{T}_3 in Figure 4, suggests that this is not always possible. In fact, as we shall show, it suffices to exclude a special type of labelled tree which we define next.

A *fixed-cherry tree* on X (with cherry $\{x_1, x_2\}$), $|X| \geq 4$, is a labelled tree $\mathcal{T} = (T, t)$ on X such that the root ρ_T of T has two children v and w with $t(v) = t(w) \neq t(\rho_T)$, v is the parent of two elements x_1 and x_2 of X , and w the parent of all elements in $X - \{x_1, x_2\}$. For example, the tree \mathcal{T}_3 in Figure 4 is a fixed-cherry tree on $X = \{1, 2, 3, 4\}$ with cherry $\{1, 2\}$. Note that if $\mathcal{T} = (T, t)$ is a fixed-cherry tree with cherry $\{x_1, x_2\}$ and $x, y, z \in X$ distinct, then $\delta_{\mathcal{T}}(x, y, z) = \{t(w), t(w), t(w)\}$ if neither x_1 nor x_2 belong to $\{x, y, z\}$ and $\delta(x, y, z) = \{t(\rho_T), t(\rho_T), t(w)\}$ else. We call a three-way symbolic map $\delta : \binom{X}{3} \rightarrow \mathcal{M}$ that satisfies these conditions for some $x_1 \neq x_2 \in X$ a *fixed cherry map* (with cherry $\{x_1, x_2\}$). The following observation is straight-forward to check.

Lemma 6.3 *Suppose that $|X| \geq 5$ and that δ is a three-way symbolic map on X . Then δ can be represented by a fixed-cherry tree on X with cherry $\{x_1, x_2\}$ if and only if δ is a fixed-cherry map with cherry $\{x_1, x_2\}$.*

Note that a triplet $xy|z$ with $x, y, z \in X$ is displayed by a fixed-cherry tree on X with cherry $\{x_1, x_2\}$ if and only if either $\{x, y\} = \{x_1, x_2\}$ or $z \in \{x_1, x_2\}$, and $x, y \in X - \{x_1, x_2\}$ hold. In particular, if $|X| > 4$ and δ is a fixed-cherry map, then the cherry can be easily identified from δ , and therefore also all of the triplets displayed by \mathcal{T} .

We now consider how to obtain the triplets displayed by a labelled tree \mathcal{T} in case \mathcal{T} is not a fixed-cherry tree. We start with a useful lemma. Suppose $\mathcal{T} = (T, t)$ is a labelled tree and $Y \subseteq X$, $|Y| \geq 4$, is such that $\delta_{\mathcal{T}|_Y}$ has a unique discriminating representation. Then we denote that representation by $\mathcal{T}_Y = (T_Y, t_Y)$.

Lemma 6.4 *Let \mathcal{T} be a discriminating labelled tree on X and assume that $Y \subseteq X$ is such that $\delta_{\mathcal{T}|_Y}$ has a unique discriminating representation. If t is a triplet displayed by \mathcal{T}_Y , then t is displayed by \mathcal{T} .*

Proof. Put $\delta = \delta_{\mathcal{T}}$ and $\mathcal{T} = (T, t)$. It suffices to note that \mathcal{T}_Y is obtained from \mathcal{T} by first taking the subtree T' of T induced by Y , and then collapsing edges of T' both of whose end vertices have the same label under the restriction t' of t to $V(T')$. Clearly, \mathcal{T}_Y is a discriminating representation of $\delta|_Y$. By assumption, it follows that \mathcal{T}_Y is the unique discriminating representation of $\delta|_Y$.

It is well-known (Semple and Steel, 2003, Theorem 6.4.1) that the set \mathcal{R} of triplets displayed by T' is contained in the set of triplets displayed by T . Since the process of collapsing edges of T' removes triplets from \mathcal{R} , but does not add any, it follows that a triplet displayed by T_Y is also displayed by T .

■

We now present the main result of this section.

Theorem 6.5 *Suppose that $|X| \geq 4$ and that \mathcal{T} is a labelled tree on X that is not a fixed-cherry tree. Then, for all $x, y, z \in X$ distinct, \mathcal{T} displays the triplet $xy|z$ if and only if one of the following two properties holds:*

- (P1) *There exists some $u \in X$ such that $\delta_{\mathcal{T}}(x, u, z) = \delta_{\mathcal{T}}(y, u, z) \neq \delta_{\mathcal{T}}(x, y, u)$ and if $|\underline{\delta_{\mathcal{T}}(x, y, u)}| = 1$ then $\delta_{\mathcal{T}}(x, y, u) \neq \delta_{\mathcal{T}}(x, y, z)$.*
- (P2) *There exists some $u \in X$ such that $|\{\delta_{\mathcal{T}}(x, u, z), \delta_{\mathcal{T}}(y, u, z), \delta_{\mathcal{T}}(x, y, u)\}| = 3$ and $m(\delta_{\mathcal{T}}(x, u, z)) = m(\delta_{\mathcal{T}}(y, u, z)) \neq m(\delta_{\mathcal{T}}(x, y, u))$.*

Proof. Put $\mathcal{T} = (T, t)$ and $\delta = \delta_{\mathcal{T}}$. Assume first that $x, y, z \in X$ distinct are such that T displays the triplet $xy|z$. Put $v = \text{lca}(x, z)$ and $w = \text{lca}(x, y)$. We proceed using a case-analysis on the structure of T . Since \mathcal{T} is not a fixed-cherry tree we need to consider the following (not necessarily disjoint) cases: (a): w is not a child of v , (b): v is not the root of T or has outdegree three or more, (c): w has a child that is neither x nor y , and (d): there exists a vertex v_0 on the path from v to z with $t(v_0) \neq t(w)$.

Case (a): Consider the parent v_0 of w , and an element u in X that is below v_0 but not below w (see Figure 5(a)). Since \mathcal{T} is a discrim-

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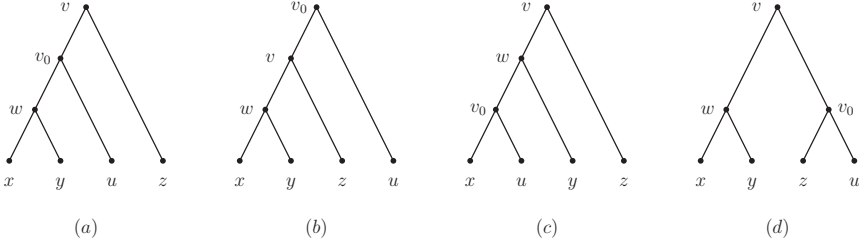


Figure 5. Cases (a)-(d) for the case-analysis carried out in the proof of Theorem 6.5. See text for details.

inating representation for $\delta_{\mathcal{T}}$, we have $t(v_0) \neq t(w)$. Hence, $\delta(x, u, z) = \delta(y, u, z) = \{t(v_0), t(v), t(v)\}$ and $\delta(x, y, u) = \{t(w), t(v_0), t(v_0)\}$. Consequently, $\delta(x, u, z) = \delta(y, u, z) \neq \delta(x, y, u)$. Note that if $t(v) = t(w)$, then $\delta(x, y, z) = \{t(w), t(v), t(v)\}$ and, so, $|\underline{\delta(x, y, z)}| = 1$. But then $\delta(x, y, u) \neq \delta(x, y, z)$ as $|\underline{\delta(x, y, u)}| = 2$. Hence, the second condition in Property (P1) holds, too. So assume that $t(w) \neq t(v)$. Then $|\underline{\delta(x, y, u)}| = 2$ and so the second condition in Property (P1) does not apply.

Case (b): Consider an element of $u \in X$ such that $v_0 := \text{lca}(u, z) = \text{lca}(u, x)$ (see Figure 5(b)). If w is not a child of v then Property (P1) follows by Case (a). So assume that w is a child of v . Then $t(v) \neq t(w)$ as \mathcal{T} is a discriminating representation for $\delta_{\mathcal{T}}$. Since $\delta(x, u, z) = \delta(y, u, z) = \{t(v), t(v_0), t(v_0)\}$ and $\delta(x, y, u) = \{t(w), t(v_0), t(v_0)\}$ we have $\delta(x, u, z) = \delta(y, u, z) \neq \delta(x, y, u)$. Since the choice of v_0 implies that $|\underline{\delta(x, y, u)}| \neq 1$ the second condition in Property (P1) does not apply. Hence, Property (P1) is also satisfied in this case.

Case (c): Then there is some $u \in X$ below w that is neither x nor y . We may assume without loss of generality that $w = \text{lca}(y, u)$. Put $v_0 = \text{lca}(x, u)$ (see Figure 5(c)). Note that $v_0 = w$ may hold. Clearly, $\delta(x, u, z) = \{t(v_0), t(v), t(v)\}$, $\delta(y, u, z) = \{t(w), t(v), t(v)\}$ and $\delta(x, y, u) = \{t(v_0), t(w), t(w)\}$. If $v_0 \neq w$ then $\delta(y, u, z) \neq \delta(x, u, z) \neq \delta(x, u, y)$. Hence, $|\{\delta(y, u, z), \delta(x, u, z), \delta(x, u, y)\}| = 3$. Since $m(\delta(x, u, z)) = t(v) = m(\delta(y, u, z))$ and $m(\delta(x, y, u)) = t(w)$, Property (P2) follows. So assume that $v_0 = w$. Then $\delta(y, u, z) = \delta(x, u, z) = \{t(w), t(v), t(v)\}$ and $\delta(y, u, x) = \{t(w), t(w), t(w)\}$. In view of Property (P1) holding if Case (a) applies, we may assume without loss of generality that w is a child of v . Since \mathcal{T} is a discriminating representation of $\delta_{\mathcal{T}}$ we have $t(v) \neq t(w)$. Hence, $\delta(y, u, z) = \delta(x, u, z) \neq \delta(x, y, u)$. Since $|\underline{\delta(x, y, z)}| = 1$ and $|\underline{\delta(x, y, z)}| \neq 1$, Property (P1) follows in this case, too.

Case (d): Let $u \in X$ such that $v_0 = \text{lca}(z, u)$ (see Figure 5(d)). Then $\delta(x, u, z) = \delta(y, u, z) = \{t(v_0), t(v), t(v)\}$ and $\delta(x, y, u) = \{t(w), t(v), t(v)\}$. If $t(w) = t(v_0)$ we have $\delta(x, u, z) = \delta(y, u, z) = \delta(x, y, u)$. In view of Property (P1) holding if Case (a) applies, we may assume without loss of generality that w is a child of v . Hence, $t(w) \neq t(v)$ because \mathcal{T} is a discriminating representation for δ . But then $|\underline{\delta(x, y, u)}| \neq 1$ and, so, the second condition in Property (P1) does not apply.

Conversely, let $x, y, z \in X$ distinct. Assume first that there exists some $u \in X - \{x, y, z\}$ such that Property (P1) is satisfied for the namesakes of u, x, y , and z . Consider the restriction δ' of δ to $\{x, y, u, z\}$. Let $\mathcal{T}' = (T', t')$ denote a discriminating representation of δ' . To see that $xy|z$ is displayed by T' we claim first that \mathcal{T}' is the unique discriminating representation of δ' . To see the claim, we show that $xy|z$ is displayed by T' . Assume for contradiction that the triplet $xy|z$ is not displayed by T' . In view of the first condition in Property (P1), the outdegree of the root $\rho_{T'}$ cannot be four. Hence, one of the triplets $x|yz$ and $y|xz$ must be displayed by T' and T' is either resolved or unresolved. Assume first that T' is resolved. Then a straight forward case analysis concerned with adding u to the triplet $x|yz$ implies that that triplet cannot be displayed by T' . Swapping the roles of x and y in that argument also implies that the triplet $y|xz$ cannot be displayed by T' either. Thus, T' must be unresolved and, so, either $\rho_{T'}$ has outdegree three or one of the children of $\rho_{T'}$ has outdegree three.

If T' displays the triplet $x|yz$ and the outdegree of $\rho_{T'}$ is three then $|\underline{\delta(y, u, z)}| = 1$. Hence, $\delta(x, y, u) = \delta(x, y, z)$ in view of the second condition in Property (P1) which is impossible. Thus, one of the children of $\rho_{T'}$ has outdegree three. But this is impossible in view of the first condition in Property (P1). Similar arguments imply that the triplet displayed by T' cannot be $y|xz$ either which is impossible. Thus, T' must display the triplet $xy|z$. Consequently, either $\rho_{T'}$ is the parent of u and z or x, y , and u have the same parent. In either case it follows that δ' cannot be of type $\hat{\delta}_3$. Thus, \mathcal{T}' is the unique discriminating representation of δ' , as claimed. By Lemma 6.4, it follows that $xy|z$ must be displayed by \mathcal{T} .

Assume next that there exists some $u \in X - \{x, y, z\}$ such that Property (P2) is satisfied for the namesakes of u, x, y , and z . Consider again the restriction δ' of δ to $\{x, y, u, z\}$. Then δ' must have a representation $\mathcal{T}' = (T', t')$. In view of the first condition of Property (P2), \mathcal{T}' must be discriminating. A straight forward case analysis implies that δ' cannot be of type $\hat{\delta}_3$. Thus, \mathcal{T}' is the unique discriminating representation of δ' .

In view of Table 1, there must exist at least two subsets Y and Y' of $\{x, y, z, u\}$ of size three satisfying $\delta(Y) = \delta(Y')$. Since $|\{\delta(x, u, z), \delta(y, u, z), \delta(x, y, u)\}| = 3$, it follows that $\{x, y, z\}$ must be one of these subsets. If $\delta(x, y, z) = \delta(x, y, u)$, then $D_{\mathcal{T}}(x, u) = m(\delta(x, u, z))$ and $D_{\mathcal{T}}(y, u) = m(\delta(y, u, z))$ must hold where $D_{\mathcal{T}}$ is the symbolic ultrametric represented by \mathcal{T} . Indeed, since $\delta(x, y, u) = \delta(x, y, z)$, one of the following two cases must hold: (α) $D_{\mathcal{T}}(x, z) = D_{\mathcal{T}}(x, u)$ and $D_{\mathcal{T}}(y, z) = D_{\mathcal{T}}(y, u)$ and (β) $D_{\mathcal{T}}(x, z) = D_{\mathcal{T}}(y, u)$ and $D_{\mathcal{T}}(y, z) = D_{\mathcal{T}}(x, u)$. However Case (β) implies $\delta(x, z, u) = \delta(y, z, u)$, which is impossible in view of the assumption that $\{\delta(x, y, u), \delta(y, z, u), \delta(x, z, u)\}$ has size three. Thus, Case (α) must hold. But then $D_{\mathcal{T}}(x, u) = m(\delta(x, u, z))$ and $D_{\mathcal{T}}(y, u) = m(\delta(y, u, z))$, as required.

Since, by assumption, we also have $m(\delta(x, u, z)) = m(\delta(y, u, z))$ we obtain $D_{\mathcal{T}}(x, u) = D_{\mathcal{T}}(y, u)$. Since $D_{\mathcal{T}}(x, u)$ and $D_{\mathcal{T}}(y, u)$ are both elements in the multiset $\delta(x, y, u)$, we obtain $m(\delta(x, u, z)) = D_{\mathcal{T}}(x, u) = m(\delta(x, y, u))$, which is impossible in view of (P2). Thus, we either have $\delta(x, y, z) = \delta(x, u, z)$ or $\delta(x, y, z) = \delta(y, u, z)$. Note that the roles of x and y are interchangeable here, so we may assume without loss of generality that $\delta(x, y, z) = \delta(y, u, z)$.

Using similar arguments as before, we have $D_{\mathcal{T}}(x, z) = m(\delta(x, u, z))$, $D_{\mathcal{T}}(x, y) = m(\delta(x, u, y))$, and $D_{\mathcal{T}}(y, z) = m(\delta(y, u, z))$ in this case. By Property (P2), it follows that $D_{\mathcal{T}}(x, z) = D_{\mathcal{T}}(y, z) \neq D_{\mathcal{T}}(x, y)$. Thus, $xy|z$ is displayed by T' and, by Lemma 6.4, $xy|z$ is also displayed by T .

■

We now explain how, as a direct consequence of Lemma 6.3 and Theorem 6.5, it is possible to decide whether or not a three-way symbolic map δ on a set X with $|X| \geq 5$ is a three-way symbolic ultrametric by considering triplets and, if so, construct the labelled tree \mathcal{T} which represents δ .

First, check if δ is a fixed-cherry map. If this is the case, then δ is a three-way symbolic ultrametric and \mathcal{T} can be easily constructed. If not, then compute the set $\text{Tr}(\delta)$ of triplets of X satisfying Properties (P1) or (P2), and use it as input to the BUILD algorithm. If there is no tree that display all the triplets in $\text{Tr}(\delta)$, then δ is not a three-way symbolic ultrametric. Otherwise, using the tree T that is constructed from the BUILD algorithm and the map δ , it is straight-forward to decide if there is a labelling map t for T such that (T, t) represents δ . If this is the case, then δ is a three-way symbolic ultrametric which has the computed labelled tree (T, t) as a representation, otherwise it is not.

Note that BUILD may return a tree T from $\text{Tr}(\delta)$ even if the map δ is not a three-way symbolic ultrametric. For example, let $M = \{A, B\}$

and consider the map $\delta : \binom{X}{3} \rightarrow \mathcal{M}$ where $X = \{1, 2, 3, 4, 5\}$, and $\delta(x, y, z) = 3A$ if $\{x, y, z\} = \{3, 4, 5\}$, and $\delta(x, y, z) = 2A + B$ otherwise. Although the map is clearly not representable by a labelled tree, we have $\text{Tr}(\delta) = \{34|1, 34|2, 35|1, 35|2, 45|1, 45|2\}$, and it is easy to check that there exists a phylogenetic tree on X whose set of displayed triplets is $\text{Tr}(\delta)$.

7. Conclusion

We conclude by presenting some possible future directions:

- In Chepoi and Fichet (2007), some relationships are derived between three-way and two-way dissimilarities in general. It would be interesting to see which of these relationships might be extendable to symbolic three-way maps.
- We define three-way tree-maps in terms of medians in leaf-labelled trees. Can any of our results be extended to median networks (Bandelt et al., 1995)? Also, can our results concerning three-way symbolic ultrametrics be extended to rooted phylogenetic networks? (cf. e.g. Huber and Scholz, 2018)
- We can clearly consider generalizations of three-way symbolic maps to k -way symbolic maps, $k \geq 2$ (see e.g. Chepoi and Fichet, 2007; Deza and Rosenberg, 2000; Warrens, 2010), and therefore generalize the notion of a three-way symbolic ultrametric in the natural way. If δ is a k -way symbolic map and its restriction to every $k + 2$ subset is a k -way symbolic ultrametric, then is δ a symbolic ultrametric?
- In Hellmuth et al. (2013), an application of symbolic ultrametrics to constructing genome-based phylogenies is presented. It would be interesting to see if this application could be extended to three-way maps. Note that results presented in Levy et al. (2006) might be relevant in this context. Also, it would be interesting to develop associated algorithms such as those in Lafond and El-Mabrouk (2015), for three-way symbolic maps.

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