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## PROVABLE WELLORDERINGS OF FORMAL THEORIES FOR TRANSFINITELY ITERATED INDUCTIVE DEFINITIONS

#### W. BUCHHOLZ and W. POHLERS

## Introduction

By [12] we know that transfinite induction up to  $\Theta \varepsilon_{\Omega_{N+1}} 0$  is not provable in ID<sub>N</sub>, the theory of N-times iterated inductive definitions. In this paper we will show that conversely transfinite induction up to any ordinal less than  $\Theta \varepsilon_{\Omega_{N+1}} 0$  is provable in ID<sup>i</sup><sub>N</sub>, the intuitionistic version of ID<sub>N</sub>, and extend this result to theories for transfinitely iterated inductive definitions.

In [14] Schütte proves the wellordering of his notational systems  $\Sigma(N)$  using predicates  $\mathfrak{B}_k(a)$ : $\Leftrightarrow (a \in M_k \land \{x \in M_k : x < a\}$  is wellordered) with  $M_k := \{x \in \Sigma(N): \mathfrak{B}_0(K_0x) \land \cdots \land \mathfrak{B}_{k-1}(K_{k-1}x)\}^1$  and  $0 \le k \le N$ . Obviously the predicates  $\mathfrak{B}_0, \ldots, \mathfrak{B}_{N-1}$  are definable in  $\mathrm{ID}_N^i$  with the defining axioms:

$$(\mathfrak{B}_{k} 1) \qquad \operatorname{Prog} [M_{k}, \mathfrak{B}_{k}],$$

$$(\mathfrak{B}_{k} 2) \qquad \operatorname{Prog} [M_{k}, \mathfrak{F}] \to \forall x (\mathfrak{B}_{k} (x) \to \mathfrak{F}[x]),$$

where  $Prog[M_k, X]$  means that X is progressive with respect to  $M_k$ , i.e.

$$\operatorname{Prog}[M_k, X] : \leftrightarrow \forall x \in M_k (\forall y \in M_k (y < x \to X(y)) \to X(x)).$$

The crucial point in Schütte's wellordering proof is Lemma 19 [14, p. 130] which can be modified to

(I) 
$$\operatorname{TI}[M_{k+1}, a], Sb = k, \mathfrak{B}_k(b) \Rightarrow \mathfrak{B}_k((a, b)), \text{ for } 0 \le k \le N-1,$$

where TI[ $M_{k+1}$ , a] is the scheme of transfinite induction over  $M_{k+1}$  up to  $a^2$ . Checking the proof of (I) it turns out that besides  $(\mathfrak{B}_k 1)$  and  $(\mathfrak{B}_k 2)$  $(0 \le k \le N-1)$  only finitary methods (including mathematical induction) are used. Since the proof uses "excluded middle" only for decidable formulas it is formalizable in ID<sup>1</sup><sub>N</sub>. Following the proof of Lemma 17 in [14] one gets

(II) 
$$\mathrm{ID}_{N}^{i} \vdash \mathfrak{B}_{0}(1) \land \cdots \land \mathfrak{B}_{N-1}(\Omega_{N-1})$$
 and

(III) 
$$\mathrm{ID}_{N}^{i} \vdash \mathrm{TI}[M_{N}, \Omega_{N}].$$

From (III) one derives in the well-known way (due to Gentzen [5])

(IV) 
$$ID_N^i \vdash TI[M_N, c_n]$$
 for each  $n \in N$ ,

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<sup>&</sup>lt;sup>1</sup>  $K_n x$  is a finite set of subterms of x.  $\mathfrak{B}_n(K_n x)$  means  $\forall y \in K_n x(\mathfrak{B}_n(y))$ .

<sup>&</sup>lt;sup>2</sup> For an exact definition see notational convention (5), page 3 of the present paper.

where  $c_0 := \Omega_N$ ,  $c_{n+1} := (1, c_n)$ . By (I), (II), (IV) and the facts that  $M_0 = \Sigma(N)$ and,  $\mathfrak{B}_k(a)$  implies  $\operatorname{TI}[M_k, a]$  one gets

(V) 
$$ID_N^i \vdash TI[\Sigma(N), \Omega[c_n, 0]]$$
 for each  $n \in N$ ,

where  $\Omega[c_n, 0] := ((\cdots (c_n, \Omega_{N-1}), \dots, \Omega_1), 1)$ . Since  $\sup_{n \in N} \Omega[c_n, 0] = \Omega[(1 \# 1, \Omega_N), 0]$  and the order type of  $\{x \in \Sigma(N): x < \Omega[(1 \# 1, \Omega_N), 0]\}$  is  $\Theta \varepsilon_{\Omega_N+1} 0^3$ , transfinite induction up to any ordinal less than  $\Theta \varepsilon_{\Omega_N+1} 0$  is provable in  $\mathrm{ID}_{N}^i$ , which we will abbreviate by  $\mathrm{ID}_N^i \vdash \mathrm{TI}[<\Theta \varepsilon_{\Omega_N+1} 0]$ .

Similar considerations apply to the wellordering proof of the system  $\overline{\Theta}(\{g\})$  given in [2]. We will prove the following *results*:

(A) 
$$ID_{\nu}^{i} \vdash TI[\langle \bar{\Theta} \varepsilon_{\Omega_{\nu}+1} 0]$$
 for any countable  $\nu \in \bar{\Theta}(\{g\})^{4}$ ,

(B) 
$$ID_{<*}^{i} \vdash TI[<\bar{\Theta}\varepsilon_{\Omega_{\Omega},+1}0]^{s},$$

where  $ID_{\nu}^{i}$  and  $ID_{<}^{i}$  are the intuitionistic versions of the theories  $ID_{\nu}$  and  $ID_{<}$  defined in [4, pp. 307–308] (see also the last paragraph of page 3 of the present paper).

 $ID_{<\nu}$  is defined to be the theory  $\bigcup_{\xi<\nu}ID_{\xi}$ . By ID we mean the theory of autonomously iterated inductive definitions (i.e. if  $\overline{ID} \vdash TI[\nu]$  then  $ID_{\nu} \subset \overline{ID}$ ). For a theory Th we take  $|Th| := \sup\{\xi \in On : Th \vdash TI[\xi]\}$ . There are the following ordinal theoretic relations:

(1)  $\overline{\Theta} \varepsilon_{\Omega_{\Omega_1}+1} 0 = \Theta \varepsilon_{\Omega_{\Omega_1}+1} 0$  and  $\overline{\Theta} \varepsilon_{\Omega_{\nu}+1} 0 = \Theta \varepsilon_{\Omega_{\nu}+1} 0$  for  $\nu < \Theta \Omega_{\Omega_1} 0$ . (So the above derived results on transfinite induction in  $ID_N^i$  ( $N < \omega$ ) are special cases of (A).)

(2)  $\Theta \Omega_{\nu} 0 = \sup_{\xi < \nu} \Theta \varepsilon_{\Omega_{\xi}+1} 0$  for limit  $\nu \le \Theta \Omega_{\Omega_1} 0$ .

(3)  $\Theta \Omega_{\nu} 0 = \nu$  for  $\nu = \Theta \Omega_{\Omega_1} 0$ .

(4)  $\Theta \Omega_{\Omega_1} 0 = \sup_{n \in \mathbb{N}} \nu_n$  with  $\nu_0 := 1$ ,  $\nu_{n+1} := \Theta \Omega_{\nu_n} 0$ .

By (A), (1)-(4) and [13] (cf. footnote 4) we get the equations:

(A1) 
$$|ID_{\nu}| = |ID_{\nu}^{i}| = \Theta \varepsilon_{\Omega_{\nu}+1} 0 \text{ for } \nu < \Theta \Omega_{\Omega_{1}} 0.$$

(A2) 
$$|ID_{<\nu}| = |ID_{<\nu}^{i}| = \Theta \Omega_{\nu} 0 \text{ for limit } \nu \leq \Theta \Omega_{\Omega_{1}} 0.$$

(A3)

 $|ID_{<\Theta\Omega_{0,0}}^{(i)}| = |\overline{ID}^{(i)}| = \Theta\Omega_{\Omega_{1}}0$  and  $\overline{ID}^{(i)}$  has the same theorems as  $ID_{<\Theta\Omega_{0,0}}^{(i)}$ .

**Preliminaries.** In the sequel we assume an arithmetization of the notational system  $\overline{\Theta}(\{g\})$ , such that all relevant ordinal sets, functions and relations of [2] (as  $\mathfrak{T}, \mathfrak{R}, K_u a, S, +, \overline{\Theta}, <$ , etc.) become primitive recursive<sup>6</sup>. We will identify ordinal notations and their arithmetizations.

Though we presume some familiarity with [2], we will give a short description of the system  $\overline{\Theta}(\{g\})$ .  $\overline{\Theta}(\{g\})$  is a set  $\mathfrak{T}$  of ordinal notations ordered by a relation <. Each element of  $\mathfrak{T}$  has the shape 0, a + b,  $\overline{\Theta}ab$  or gab with

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<sup>&</sup>lt;sup>3</sup> Cf. [3]. Note that the system  $\Sigma(N)$  in [3] is a slight modification of that in [14]. In [3] the first element of  $\Sigma(N)$  is 0 instead of 1.

<sup>&</sup>lt;sup>4</sup> Recently the second author [13] was able to show  $ID_{\nu} \not\vdash TI[\bar{\Theta}\varepsilon_{\Omega_{\nu+1}}0]$ .

<sup>&</sup>lt;sup>5</sup>Kino's wellordering proof for her ordinal diagrams Od(*I*) [8] is formalizable in  $ID_{<+}^{i}$ . Hence  $ID_{<+}^{i}$  + TI [ $\|Od(I), <_{\infty}\|$ ]. But as remarked in [2]  $\|Od(I), <_{\infty}\| \le \Theta\Omega_{\Omega_{1}}(\tau+1) < \Theta(\Omega_{\Omega_{1}}+1)0 < \Theta\varepsilon_{\Omega_{\Omega_{1}}+1}0$  for  $\|I\| = 1 + \tau < \Theta(\Omega_{\Omega_{1}}+1)0$ .

<sup>&</sup>lt;sup>6</sup>For a subsystem of  $\overline{\Theta}(\{g\})$  such an arithmetization will be carried out in [15].

 $a, b \in \mathfrak{T}$ . The symbols + (ordinal sum),  $\overline{\Theta}$  and g denote 2-place ordinal functions. So each term  $a \in \mathfrak{T}$  canonically represents an ordinal |a|, for example  $|\overline{\Theta}00| = 1$ ,  $|\overline{\Theta}01| = \omega$ ,  $|\overline{\Theta}10| = \varepsilon_0$ . For the order relation < we have  $a < b \leftrightarrow |a| \in |b|$ . Terms of the shape  $\overline{\Theta}ab$  or gab are called main terms; they represent ordinals closed under +.  $\Re$  is the set of all terms  $gab \in \mathfrak{T}$ ; the elements of  $\Re$  represent initial ordinals >  $\omega$ . There is a primitive recursive order isomorphism  $a \mapsto \Omega_a$  from  $\mathfrak{T}$  onto  $\Re_0 := \{0\} \cup \Re$  with  $\Omega_0 = 0$  and  $|\Omega_a| = \Omega_{|a|}$  for  $a \neq 0$ . For each  $a \in \mathfrak{T}$  there is exactly one  $x \in \mathfrak{T}$  with  $\Omega_x \le a < \Omega_{x+1}$ ; we define  $Sa := \Omega_x$  and call it the level (Stufe) of a. For all  $a, b \in \mathfrak{T}$  we have  $S\overline{\Theta}ab = Sb$ , hence  $\overline{\Theta}a0 < \Omega_1$ . We define  $M_0 := \{x \in \mathfrak{T}: Sx = 0\} = \{x \in \mathfrak{T}: x < \Omega_1\}$ .  $M_0$  represents a segment of the countable ordinals, i.e.

(1)  $\{\xi \in On : \xi < |a|\} = \{|x| : x \in M_0 \land x < a\}$  for  $a \in M_0$ .

For  $a \in \mathfrak{T}$  and  $u \in \mathfrak{R}_0$ ,  $K_u a$  is a finite set of main terms with levels  $\leq u$ . The sets  $K_u a$  have the following properties:

(2) If  $Sa \le u$ , then  $K_u a$  is the set of components of  $a^7$ .

(3)  $K_u b \subset K_u (a+b) \subset K_u a \cup K_u b$ .

(4)  $w \leq v \wedge c \in K_v a \rightarrow K_w c \subset K_w a$ .

(5)  $v < \Omega_a \rightarrow K_v \Omega_a \subset K_v a \subset K_v \Omega_a \cup \{1\}.$ 

We fix the following notational conventions:

(1) a, b, c, x, y, z denote elements of  $\mathfrak{T}$ .

(2) u, v, w denote elements of  $\Re_0$ .

(3)  $\Re, \mathfrak{X}$  serve as syntactical variables for sets  $\{x : \mathfrak{F}[x]\} \subset \mathfrak{T}$ , where  $\mathfrak{F}[x]$  is a formula of the theory considered.

(4)  $\mathcal{X} \cap a := \{x \in \mathcal{X} : x < a\}, \ \mathcal{X} \cap u^+ := \{x \in \mathcal{X} : Sx \le u\}.$ 

(5)  $\operatorname{Prog} [\mathfrak{R}, \mathfrak{X}]$  abbreviates the formula  $\forall x \in \mathfrak{R}(\mathfrak{R} \cap x \subset \mathfrak{X} \to x \in \mathfrak{X})$ . TI[ $\mathfrak{R}, \mathfrak{X}, a$ ] abbreviates the fromula  $a \in \mathfrak{R} \land (\operatorname{Prog} [\mathfrak{R}, \mathfrak{X}] \to \mathfrak{R} \cap a \subset \mathfrak{X})$ , and TI[ $\mathfrak{R}, a$ ] denotes the scheme  $\{\operatorname{TI} [\mathfrak{R}, \mathfrak{X}, a]\}_{\mathfrak{X}}$  expressing the principle of transfinite induction over  $\mathfrak{R}$  up to a.

Transfinite inductions provable in  $ID_{\nu}^{i}$  and  $ID_{<}^{i}$ . Our main tool in proving transfinite inductions will be the concept of the accessible part  $W[\mathfrak{R}]$  of a set  $\mathfrak{R}$ , usually defined by  $W[\mathfrak{R}] = \{x \in \mathfrak{R} : \mathfrak{R} \cap x \text{ is wellordered}\}$ , which is a second-order definition. This definition however can be replaced by an inductive definition, which is expressible in a first-order language by the infinite list of axioms:

(i)  $\operatorname{Prog}[\mathfrak{R}, W[\mathfrak{R}]]$  and

(ii)  $\operatorname{Prog}\left[\mathfrak{R}, \mathfrak{X}\right] \to W[\mathfrak{R}] \subset \mathfrak{X}$  for each  $\mathfrak{X}$ .

The theories  $ID_{\nu}^{i}$  (with  $\nu \in M_{0}$ ) and  $ID_{<}^{i}$  are formal theories for iterations of such inductive definitions. They are first-order extensions of Heyting's arithmetic, where  $ID_{\nu}^{i}$  allows iteration of monotone inductive definitions along the segment  $M_{0} \cap \nu$ , while  $ID_{<}^{i}$  allows iteration along the accessible part  $W_{0} := W[M_{0}]$  of  $M_{0}$ . Besides the axioms for iteration of inductive definitions (cf. [4, p. 307, (i), (ii)]) there are the axioms:

(TI<sub> $\nu$ </sub>) Prog[ $M_0, \mathcal{X}$ ]  $\rightarrow M_0 \cap \nu \subset \mathcal{X}$  for each  $\mathcal{X}$ , in ID<sup>*i*</sup><sub> $\nu$ </sub>,

<sup>&</sup>lt;sup>7</sup>For each  $a \neq 0$  there are uniquely determined main terms  $a_1 \ge \cdots \ge a_n$   $(n \ge 1)$  such that  $a = a_1 + \cdots + a_n$ . We call  $a_1, \ldots, a_n$  the components of a. 0 is defined to have no components.

asserting the wellordering of  $M_0$ ,

$$(W_01) \qquad \qquad \operatorname{Prog}[M_0, W_0] \qquad \text{and} \qquad \qquad$$

$$(W_02) \qquad \operatorname{Prog}[M_0, \mathcal{X}] \to W_0 \subset \mathcal{X} \qquad \text{for each } \mathcal{X}, \text{ in } \mathrm{ID}_{<^*}^i.$$

defining the accessible part  $W_0$  of  $M_0$ .

In order to treat  $ID_{*}^{i}$  and  $ID_{*}^{i}$  simultaneously as far as possible, we refer to both as  $ID^{i}$  and define A to be the set  $\{\Omega_{x} : x < \nu\}$  in the case of  $ID_{*}^{i}$  and the set  $\{\Omega_{x} : x \in W_{0}\}$  in the case of  $ID_{*}^{i}$ . Then A is a segment of  $\Re_{0} \cap \Omega_{\Omega_{1}}$  with  $0 \in A$ .

In the sequel u, v are reserved to denote elements of A!

Define  $\mathfrak{A}[X, Y, x, y]$  to be the formula  $\mathfrak{F}[x] \land \forall x_0 < x(\mathfrak{F}[x_0] \to x_0 \in X)$ , where  $\mathfrak{F}[x]$  stands for  $Sx \leq \Omega_y \land \forall z_1 < y(\{(z_0, z_1): z_0 \in K_{\Omega_{z_1}}x\} \subset Y))$ . Then  $\mathfrak{A}[X, Y, x, y]$  is an arithmetic formula such that each occurrence of X is positive. To  $\mathfrak{A}$  corresponds a set constant  $P^{\mathfrak{A}}$  (cf. [4, p. 307]). We define

$$W_{\Omega_{y}} := \{ x : \langle x, y \rangle \in P^{\mathfrak{A}} \} \text{ and}$$
$$M_{u} := \{ x : Sx \leq u \land \forall v < u (K_{v}x \subset W_{v}) \}.$$

Then the axioms (i) and (ii) of [4, p. 307] become

(W1) 
$$\forall u \in A (\operatorname{Prog}[M_u, W_u])$$
 and

(W2)  $\forall u \in A (\operatorname{Prog}[M_u, \mathcal{X}] \to W_u \subset \mathcal{X})$  for each  $\mathcal{X}$ .

(W1) and (W2) assert that  $W_u$  is the accessible part of  $M_u$ . Clearly for u = 0,  $M_u$  coincides with the previously defined set  $M_0 = \mathfrak{T} \cap \Omega_1$ , and in the case of  $ID'_{<}$ . the set  $W_0$  defined by (W1), (W2) coincides with the set  $W_0$  defined by (W01), (W2) coincides with the set  $W_0$  defined by (W01), (W2) coincides with the set  $W_0$  defined by (W01), (W2) coincides with the set  $W_0$  defined by (W01), (W2) coincides with the set  $W_0$  defined by (W01), (W2) coincides with the set  $W_0$  defined by (W01), (W2) coincides with the set  $W_0$  defined by (W01), (W2) coincides with the set  $W_0$  defined by (W01), (W2) coincides with the set  $W_0$  defined by (W01), (W2) coincides with the set  $W_0$  defined by (W01), (W2) coincides with the set  $W_0$  defined by (W01), (W2) coincides with the set  $W_0$  defined by (W01), (W2) coincides with the set  $W_0$  defined by (W01), (W2) coincides with the set  $W_0$  defined by (W01), (W2) coincides with the set  $W_0$  defined by (W01), (W2) coincides with the set  $W_0$  defined by (W01), (W2) coincides with the set  $W_0$  defined by (W01), (W2) coincides with the set  $W_0$  defined by (W01), (W02).

(6)  $\forall x \in W_u (x \in M_u \land M_u \cap x = W_u \cap x)$ , i.e.  $W_u$  is a segment of  $M_u$ .

(7)  $a \in W_u \to \mathrm{TI}[M_u, \mathfrak{X}, a].$ 

By (3) and the definition of  $M_u$  we get

(8)  $a, b \in M_u \rightarrow a + b \in M_u$  and  $a + b \in M_u \rightarrow b \in M_u$ .

The following lemmata 1-3 are straightforward modifications of corresponding lemmata in [9], [10], [11] and [14].

LEMMA 1. (a)  $a, b \in W_u \rightarrow a + b \in W_u$  and

(b)  $Sa \le u \land K_u a \subset W_u \rightarrow a \in W_u$  are provable in  $ID^i$ .

PROOF. By (6) and (8).  $a, b \in W_u \land \forall x \in M_u \cap b(a + x \in W_u) \rightarrow a + b \in M_u \land M_u \cap (a + b) \subset W_u$ . Hence by (W1),  $a \in W_u \rightarrow \operatorname{Prog}[M_u, \{x : x + a \in W_u\}]$  and thence by (W2),  $a, b \in W_u \rightarrow a + b \in W_u$ . Part (b) is an immediate consequence of (a) and (2).

LEMMA 2. (a)  $a \in W_u \to K_v a \subset W_v$  and

(b)  $v < u \rightarrow W_v = W_u \cap v^+$  are provable in  $ID^i$ .

PROOF. Suppose  $a \in W_u \wedge v = \Omega_x$ . Using  $(\text{TL}_v)$  or  $(W_02)$  resp. we prove  $K_v a \subset W_v$  by transfinite induction on x. For v < u we have  $K_v a \subset W_v$  by  $a \in W_u \subset M_u$ . From  $u \leq v$  we get  $M_v \cap u^+ \subset M_u$ , hence by (W1),  $\text{Prog}[M_u, \{x : x \in M_v \to x \in W_v\}]$  and thence by  $(W2), W_u \subset \{x : x \in M_v \to x \in W_v\}$ , i.e.  $W_u \cap M_v \subset W_v$ . By the induction hypothesis we have  $K_w a \subset W_w$  for all

w < v and hence  $a \in W_u \cap M_v \subset W_v$ . By (2), (4) and (6) we then get  $K_v a \subset M_v \cap (a+1) \subset W_v$ . Part (b) follows from (a) by Lemma 1(b).

LEMMA 3.  $u \in W_u$  is provable in  $ID^i$ .

PROOF. We have  $\{x: \Omega_x \in A\} \subset W_0$ , which is trivial for  $ID_{\leq^*}^i$  and proved by  $(TI_{\nu})$  and (W1) in  $ID_{\nu}^i$ . So for  $u = \Omega_x$  we get  $x \in W_0$  and thence by (5) and Lemma 2(a),  $K_{\nu}u \subset K_{\nu}x \subset W_{\nu}$  for all v < u, which implies  $u \in M_u$ . Now suppose  $a \in M_u \cap u$ , then  $Sa \in A \cap u$ ,  $K_{Sa}a \subset W_u$  and by the lemmata 1(b), 2(b),  $a \in W_u$ . Hence  $M_u \cap u \subset W_u$  and by (W1),  $u \in W_u$ .

DEFINITION.  $Q := \{x : \exists u (x \in W_u)\} = \bigcup_{u \in A} W_u; M := \{x : \forall v (K_v x \subset W_v)\}.$ 

Consequences. 1. Obviously  $Sa \in A$  for all  $a \in Q$ . By Lemma 3 we then have  $\forall x (x \in Q \rightarrow Sx \in Q)$  and  $Q \cap \Re_0 = A$ . Hence by Lemma 2(b) for all  $u \in Q$ 

 $(9) \quad Q \cap u^+ = W_u$ 

and  $M_u = \{x : Sx \le u \land \forall w (w \in Q \cap u \to K_w x \subset Q)\}$ . That means the set Q is "ausgezeichnet" in the sense of [2, p. 18] with  $M_u^O \cap u^+ = M_u$  and  $W_u^O = W[M_u^O \cap u^+] = W_u$ .

2. By (6) and (9)  $\operatorname{Prog}[Q, \mathcal{X}] \to \operatorname{Prog}[M_u, \{x : x \in W_u \to x \in \mathcal{X}\}]$  and thence by (W2)

(10)  $\operatorname{Prog}[Q, \mathfrak{X}] \to Q \subset \mathfrak{X}$  for each  $\mathfrak{X}$ ,

which is the first-order formulation of the fact that Q is wellordered.

3. Since Q is "ausgezeichnet" (provable in  $ID^i$ ) we may follow the proof of Theorem 15(b) in [2, p. 19] and get the formula

$$a \in M \land \forall x \in M \cap a(Q \subset R_x) \rightarrow \operatorname{Prog}[Q, R_a]^{\mathrm{s}},$$

where  $R_a := \{y : \overline{\Theta} ay \in \mathcal{X} \to \overline{\Theta} ay \in Q\}$ . Here besides the premise "Q ausgezeichnet" only methods formalizable in Heyting's arithmetic are used. By (10) it follows

(11)  $\operatorname{Prog}[M, \{x : \forall y \in Q (\overline{\Theta} xy \in \mathfrak{T} \to \overline{\Theta} xy \in Q)\}].$ 

4. From outside we know that  $\overline{\Theta}(\{g\}) = (\mathfrak{T}, <)$  is wellordered and hence  $W_u = M_u = \{x : Sx \le u\}$  and  $M = \mathfrak{T}$  which implies  $Q = M \cap \Omega_{\sigma}$ ,  $\sigma$  defined by:

DEFINITION.

$$\sigma := \begin{cases} \nu, & \text{in the case of } \mathrm{ID}_{\nu}^{i}, \\ \Omega_{1}, & \text{in the case of } \mathrm{ID}_{<^{\bullet}}^{i}. \end{cases}$$

Of course  $W_u = M_u = \{x : Sx \le u\}$  is not provable in ID<sup>*i*</sup>, but the weaker assertion  $Q = M \cap \Omega_{\sigma}$  is provable as the following theorem shows.

THEOREM 1.  $\Omega_{\sigma} \in M$  and  $Q = M \cap \Omega_{\sigma}$  are provable in  $ID^{i}$ .

**PROOF.** By Lemma 2(a) we have  $Q \subset M \cap \Omega_{\sigma}$ . If  $a \in M$  and  $Sa \in A$  we get  $K_{Sa} \ a \subset W_{Sa}$  and by Lemma 1(b),  $a \in W_{Sa} \subset Q$ . So we just have to prove  $a \in M \cap \Omega_{\sigma} \rightarrow Sa \in A$  and  $\Omega_{\sigma} \in M$ . The proofs differ for  $\mathrm{ID}_{\nu}^{i}$ ,  $\mathrm{ID}_{<^{\bullet}}^{i}$ .

1. ID<sup>*i*</sup><sub> $\nu$ </sub>. Then  $A = \{w : w < \Omega_{\sigma}\}$  and trivially  $a \in M \cap \Omega_{\sigma} \to Sa \in A$  holds. By (TI<sub> $\nu$ </sub>) and (W1) we get  $\sigma = \nu \in W_0$ . Hence by Lemma 2(a) and (5)  $\forall v (K_{\nu}\Omega_{\sigma} \subset K_{\nu}\sigma \subset W_{\nu})$  which means  $\Omega_{\sigma} \in M$ .

<sup>&</sup>lt;sup>8</sup> In [2] this formula is proved with Q in place of M, but an analysis of the proof shows that it is enough to have the premise  $a \in M \land \forall x \in M \cap a(Q \subset R_x)$ .

2.  $\mathrm{ID}_{<}^{i}$ . Suppose  $a \in M \cap \Omega_{\sigma}$ . Then  $K_0 a \subset W_0$  and  $Sa = \Omega_x$  for some  $x < \Omega_1 = \sigma$ . By (5)  $K_0 x \subset K_0 \Omega_x \cup \{1\}$ . Further  $K_0 Sa \subset K_0 a$  and  $1 \in W_0$ . We therefore get by Lemma 1(b)  $x \in W_0$  and thence  $Sa \in A$ . For 0 < u,  $K_0 \Omega_{\Omega_1} = \emptyset$  and  $K_u \Omega_{\Omega_1} = \{\Omega_1\}$ . Obviously  $\Omega_1 \in A$  and hence  $\Omega_1 \in W_{\Omega_1}$  by Lemma 3. By Lemma 2(b) it follows  $\forall u (K_u \Omega_{\Omega_1} \subset W_u)$ , which means  $\Omega_{\sigma} = \Omega_{\Omega_1} \in M$ .

Now by (10) and Theorem 1 we get

(12) TI[ $M, \Omega_{\sigma}$ ].

Hence by (11)  $\forall y \in Q(\bar{\Theta}\Omega_{\sigma}y \in \mathfrak{T} \to \bar{\Theta}\Omega_{\sigma}y \in Q)$ . Since  $0 \in Q \land \bar{\Theta}\Omega_{\sigma}0 \in \mathfrak{T} \cap \Omega_1$  we obtain  $\bar{\Theta}\Omega_{\sigma}0 \in Q \cap \Omega_1$ . Hence by (9) and (7) TI[ $M_0, \bar{\Theta}\Omega_{\sigma}0$ ]. This means we are able to *collapse* the wellordering  $M \cap \Omega_{\sigma}$  to the provable ordinal  $\bar{\Theta}\Omega_{\sigma}0^\circ$ . This is a special case of the following theorem, which is proved by the above considerations with c in place of  $\Omega_{\sigma}$ .

THEOREM 2 (COLLAPSING PROPERTY). If TI[M, c] is provable in  $ID^i$  and  $\overline{\Theta}c0 \in \mathfrak{T}$ , then  $TI[M_0, \overline{\Theta}c0]$  is provable in  $ID^i$ .

Starting from (12) we now prove TI[M, c] for each  $c \in M \cap \overline{\Theta} 1\Omega_{\sigma}$  using Gentzen's [5] method for proving  $TI[<\varepsilon_0]$  in number theory.

DEFINITION.  $\bar{\mathcal{X}} := \{x : \bar{\Theta} \cap \Omega_{\sigma} \leq x \lor \forall y (M \cap y \subset \mathcal{X} \to M \cap (y + \bar{\Theta} \cap x) \subset \mathcal{X})\}.$ 

LEMMA 4.  $\forall y (M \cap y \subset \tilde{X} \to M \cap (y + \Omega_{\sigma}) \subset \tilde{X}) \to \operatorname{Prog}[M, \overline{\tilde{X}}]$  is provable in ID<sup>i</sup>.

PROOF. We have to prove  $M \cap (b + \overline{\Theta}0a) \subset \mathcal{X}$  under the assumptions (1)  $\forall y (M \cap y \subset \mathcal{X} \to M \cap (y + \Omega_{\sigma}) \subset \mathcal{X}),$  (2)  $M \cap a \subset \overline{\mathcal{X}},$  (3)  $a < \overline{\Theta}1\Omega_{\sigma},$ (4)  $M \cap b \subset \mathcal{X}.$ 

By (1) and (4) we get  $M \cap (b + \Omega_{\sigma} \cdot n) \subset \mathcal{X}$  for all  $n \in \mathbb{N}$  using mathematical induction. Hence  $M \cap (b + \overline{\Theta}0\Omega_{\sigma}) \subset \mathcal{X}$  because of  $\sup_{n \in \mathbb{N}} \Omega_{\sigma} \cdot n = \overline{\Theta}0\Omega_{\sigma}$ . Suppose  $z \in M \cap (b + \overline{\Theta}0a)$ . We may assume  $b + \overline{\Theta}0\Omega_{\sigma} \leq z$ . By (3)  $z < b + \overline{\Theta}0a < b + \overline{\Theta}1\Omega_{\sigma}$ . Hence  $z = b + \overline{\Theta}0a_1 \cdot n + z_1$  with  $1 \leq n \in \mathbb{N}$ ,  $\Omega_{\sigma} \leq a_1 < a, z_1 < \overline{\Theta}0a_1$ . By  $\forall v(v < \Omega_{\sigma} = Sa_1)$  and the definition of  $K_v$ —it is the case that  $K_va_1 = K_v\overline{\Theta}0a_1 \subset K_vz$ . So we get  $a_1 \in M \cap a$  since  $z \in M$  is assumed. By (2), (3), (4) we get  $M \cap (b + \overline{\Theta}0a_1 \cdot (n+1)) \subset \mathcal{X}$  using mathematical induction. Hence  $z \in \mathcal{X}$ .

DEFINITION.  $c_0 := \Omega_{\sigma}, c_{n+1} := \overline{\Theta} 0 c_n$ .

One easily proves  $c_n \in M$ ,  $c_n < \overline{\Theta} 1 \Omega_{\sigma} = \sup_{k \in \mathbb{N}} c_k$  and  $\overline{\Theta} c_n 0 \in \mathfrak{T}$ .

THEOREM 3. TI[M,  $c_n$ ] is provable in ID<sup>i</sup> for each  $n \in N$ .

PROOF. We prove the theorem by 'metainduction' on *n*. By (3) it follows that  $a + b \in M \rightarrow b \in M$ . Hence  $\operatorname{Prog}[M, \mathcal{X}] \land M \cap a \subset \mathcal{X} \rightarrow \operatorname{Prog}[M, \{x : a + x \in M \rightarrow a + x \in \mathcal{X}\}]$  and thence by (12)

(\*) 
$$\operatorname{Prog}[M, \mathcal{X}] \to \forall y (M \cap y \subset \mathcal{X} \to M \cap (y + \Omega_{\sigma}) \subset \mathcal{X}).$$

By (\*) and  $c_0 = \Omega_{\sigma} \in M$  we have  $\operatorname{TI}[M, c_0]$ . For n > 0 we have the induction hypothesis  $\operatorname{TI}[M, c_{n-1}]$ . Hence  $\operatorname{Prog}[M, \overline{X}] \to c_{n-1} \in \overline{X}$  which implies  $\operatorname{Prog}[M, \overline{X}] \to M \cap c_n \subset \overline{X}$ . By (\*) and Lemma 4 we get  $\operatorname{Prog}[M, \overline{X}] \to \operatorname{Prog}[M, \overline{X}]$  and therefore  $\operatorname{Prog}[M, \overline{X}] \to M \cap c_n \subset \overline{X}$ . Hence  $\operatorname{TI}[M, c_n]$ .

THEOREM 4. In ID<sup>i</sup>, TI[ $M_0$ , a] is provable for each  $a < \overline{\Theta}(\overline{\Theta} 1 \Omega_{\sigma}) 0$ .

<sup>&</sup>quot;Remember that by (1)  $M_0 \cap \overline{\Theta} \Omega_\sigma 0$  represents the ordinal  $\overline{\Theta} \Omega_\sigma 0$ .

**PROOF.** For  $a < \overline{\Theta}(\overline{\Theta} 1\Omega_{\sigma})0$  there is an  $n \in \mathbb{N}$  with  $a < \overline{\Theta}c_n 0$ . By Theorem 3 we have  $TI[M, c_n]$ , which can be collapsed to  $TI[M_0, \overline{\Theta}c_n 0]$  by Theorem 2. So  $TI[M_0, a]$  holds.

Since  $M_0 \cap \overline{\Theta}(\overline{\Theta} 1\Omega_{\sigma})0$  represents the segment  $\{\xi \in On : \xi < \overline{\Theta} \varepsilon_{\Omega_{\sigma+1}}0\}$ , the results (A) and (B) stated in the introduction follow from Theorem 4.

FINAL REMARKS. The above proof of transfinite induction admits the following generalization. Let Th be a theory containing Heyting's arithmetic and axioms for iterations of inductive definitions along a provably wellordered subset Α of Ra. Then the sets  $W_{\mu} := W[M_{\mu}], \quad M_{\mu} := \{x : Sx \leq x\}$  $u \wedge \forall v \in A \cap u(K_v x \subset W_v) \}$  $Q := \{x : \exists u \in A \, (x \in W_u)\}$  $(u \in A),$ and  $M := \{x : \forall u \in A \ (K_u x \subset W_u)\}$  are definable in Th, and if the formula  $\forall u \in W_u$  $A(u \in W_u \land \forall x \in W_u(Sx \in A))$  is provable in Th, one gets:

I. Q is wellordered, i.e.

$$\mathrm{Th} \vdash \mathrm{Prog}[Q, \mathcal{X}] \to Q \subset \mathcal{X}.$$

II. Collapsing property.

$$\mathrm{Th} \vdash \mathrm{TI}[M, c] \Rightarrow \mathrm{Th} \vdash \mathrm{TI}[M_0, \bar{\Theta}c0] \qquad (\text{for } \bar{\Theta}c0 \in \mathfrak{T}).$$

III. Extension to the next  $\varepsilon$ -number.

$$\operatorname{Th} \vdash \operatorname{TI}[M, \Omega_a] \Rightarrow \operatorname{Th} \vdash \operatorname{TI}[M, c] \quad \text{for each } c \in M \cap \overline{\Theta} 1\Omega_a.$$

From I, II, III it follows:

IV.

$$\mathrm{Th} \vdash \Omega_a \in M \land M \cap \Omega_a \subset Q \Rightarrow \mathrm{Th} \vdash \mathrm{TI} [<\bar{\Theta}(\bar{\Theta} \mid \Omega_a) 0]$$

for  $\overline{\Theta}(\overline{\Theta} 1 \Omega_a) 0 \in \mathfrak{T}$ .

As an example we regard the following definition by transfinite recursion on  $\nu \in \mathfrak{T} \cap \Omega_1$ :

 $\lambda(0) := 0, \qquad A_0 := \emptyset.$   $\lambda(\nu+1) := \Omega_{\lambda(\nu)+1}, \qquad A_{\nu+1} := A_{\nu} \cup \{\Omega_x : \lambda(\nu) \le x \in W[M^{\nu}]\},$ with  $M^{\nu} := \{x : x < \lambda(\nu+1) \land \forall u \in A_{\nu}(K_u x \subset W_u)\}$ and  $W_u$  defined as above by iteration of
inductive definitions along  $A_{\nu}$ .

 $\lambda(\nu) := \sup_{\xi < \nu} \lambda(\xi)^{10}, \quad A\nu := \bigcup_{\xi < \nu} A_{\xi} \text{ for limit ordinals } \nu.$ 

Let  $ID_{\nu}^{*}$  be the theory, which allows to define  $A_{\nu}$  and to iterate inductive definitions along  $A_{\nu}$  (ID<sup>\*</sup> for example is ID<sub><</sub>). Then by the above considerations we get:

$$\mathrm{ID}_{\nu}^{*} \vdash \mathrm{TI}[<\bar{\Theta}(\bar{\Theta}1\Omega_{\lambda(\nu)})0].$$

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